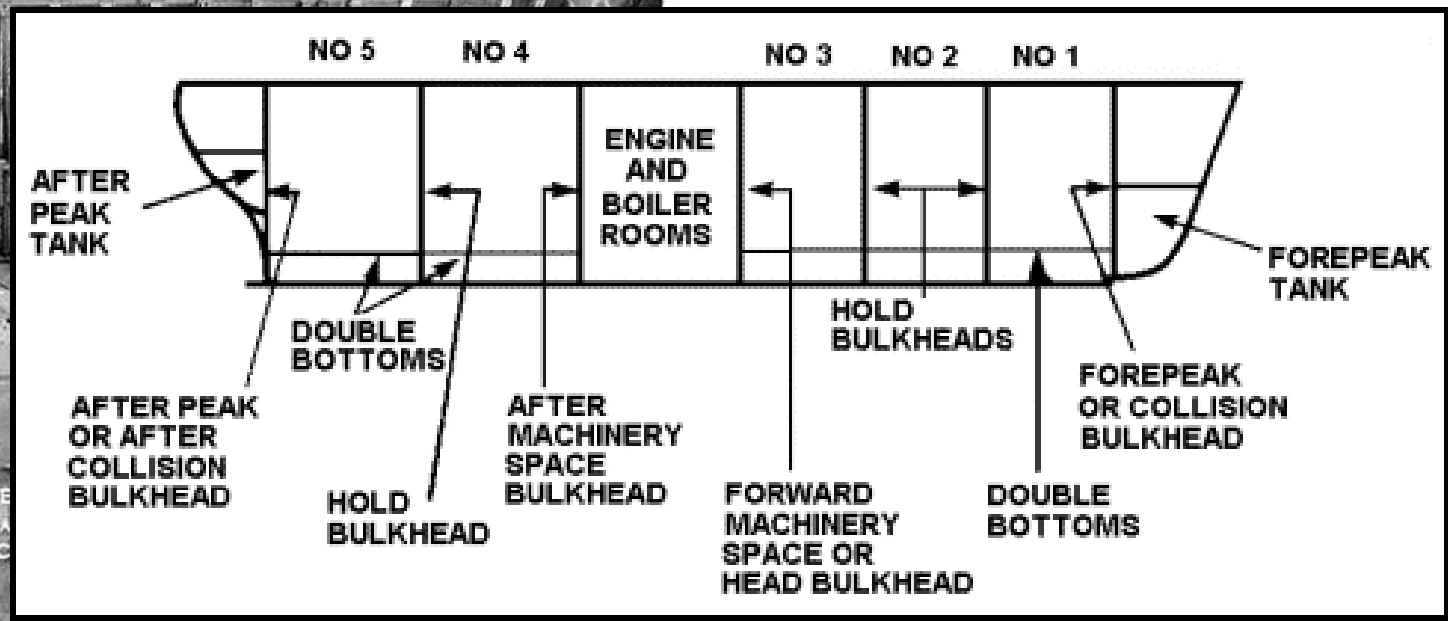
on

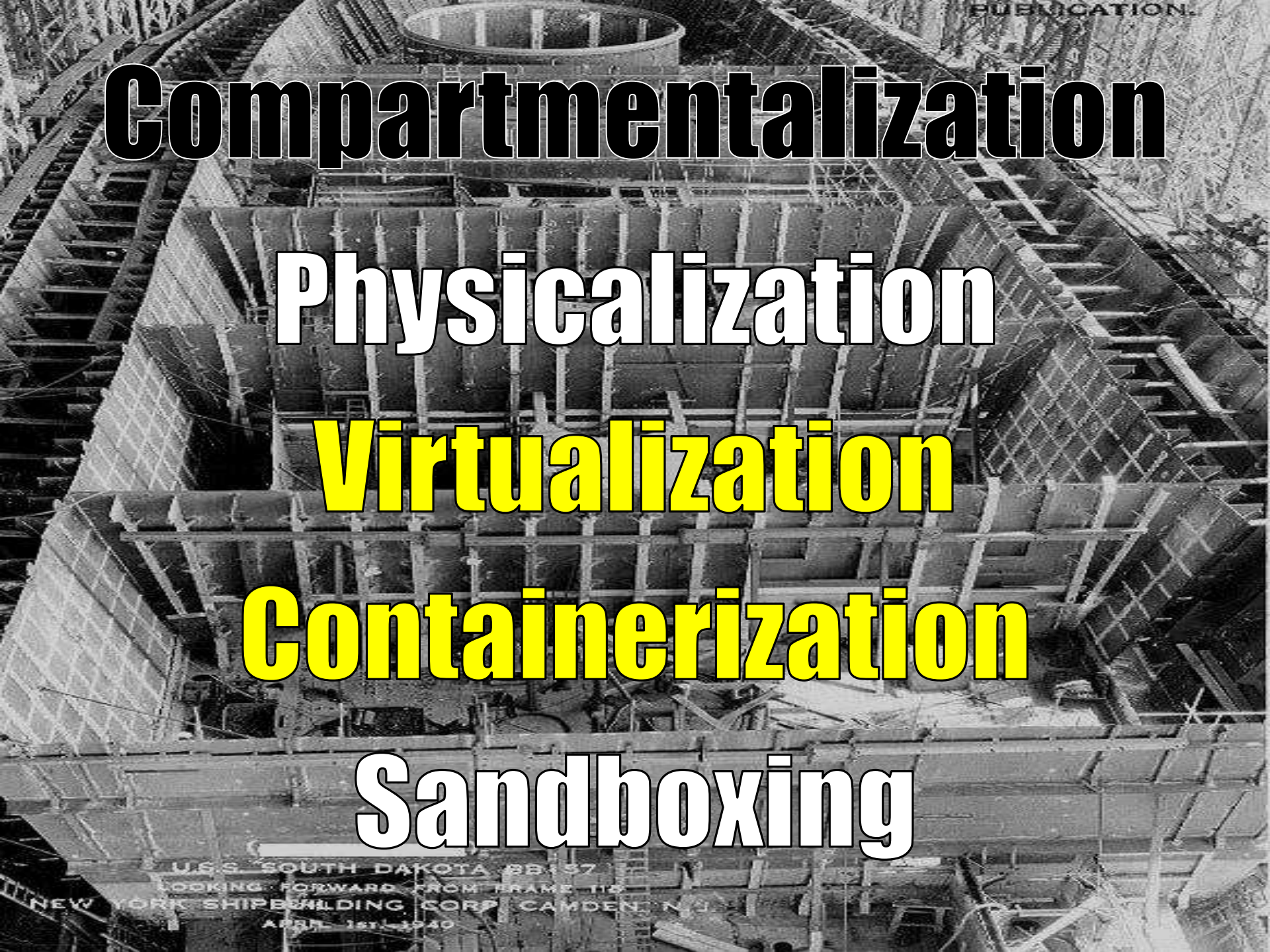
```
$ whoami

Name:     Zvi Avraham
E-mail:   zvi@nivertech.com
```

Compartmentalization

/ˈkɒm.
pɑː(ɹ)t.
mɛnt̬l.
aɪ̯ˌzeɪ.
ʃən/




Compartmentalization
Physicalization
Virtualization
Containerization
Sandboxing
U.S.S. SOUTH DAKOTA BB 57
LOOKING FORWARD FROM FRAME 115
NEW YORK SHIPBUILDING CORP. CAMDEN, N.J.
APRIL 1st 1940

Physicalization

- The opposite of Virtualization
- dedicated machines
- no virtualization overhead
- no noisy neighbors
  - nobody stealing your CPU cycles, IOPS or bandwidth
  - your EC2 instance may have a Netflix “roommate” ;)
- Mostly used by ARM-based public clouds
- also called *Bare Metal* or HPC clouds

Sandbox –
a virtual container in which untrusted code can be safely run

Sandbox examples: ZeroVM & AWS Lambda

based on Google Native Client: A Sandbox for Portable, Untrusted x86 Native Code

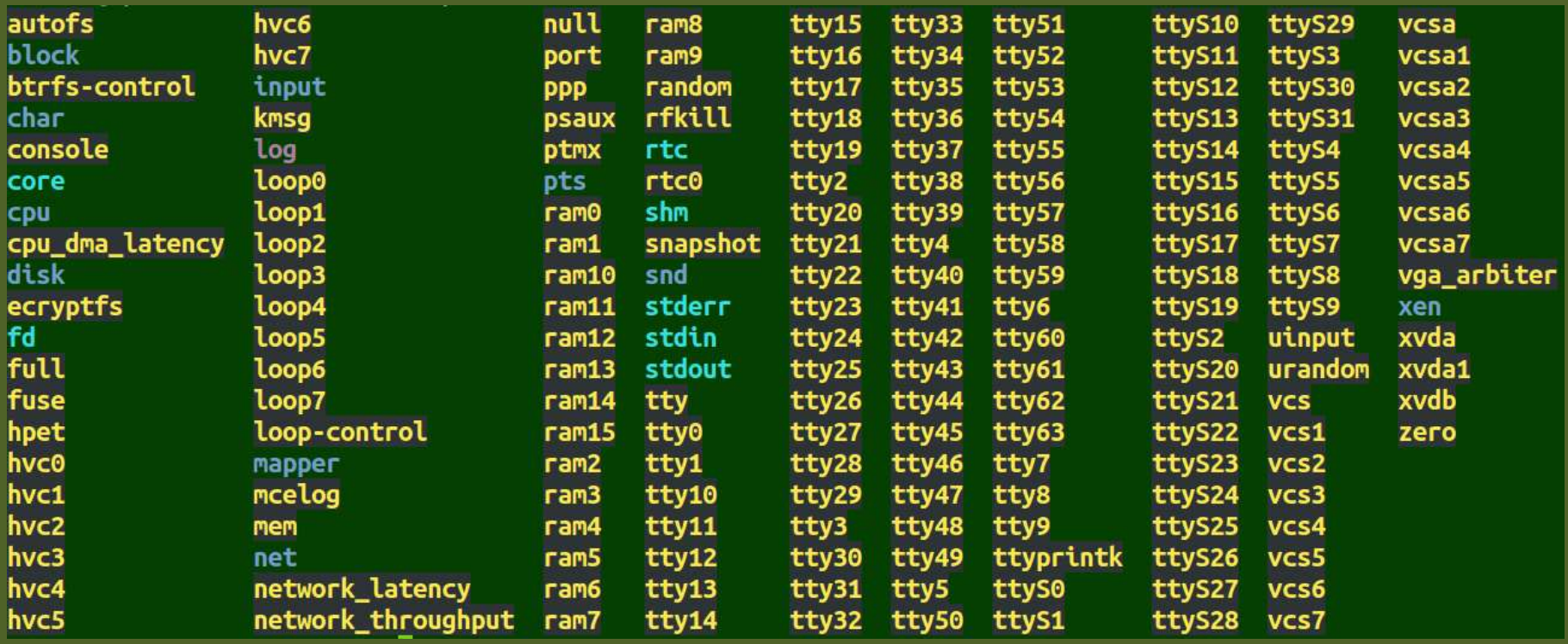
Compartmentalization in terms of Virtualization

Physicalization	No Virtualization
Virtualization	HW-level Virtualization
Containerization	OS-level Virtualization
Sandboxing	Userspace-level Virtualization*

Cloud runs on virtual HW

Does the OS on your Cloud instance still supports floppy drive?

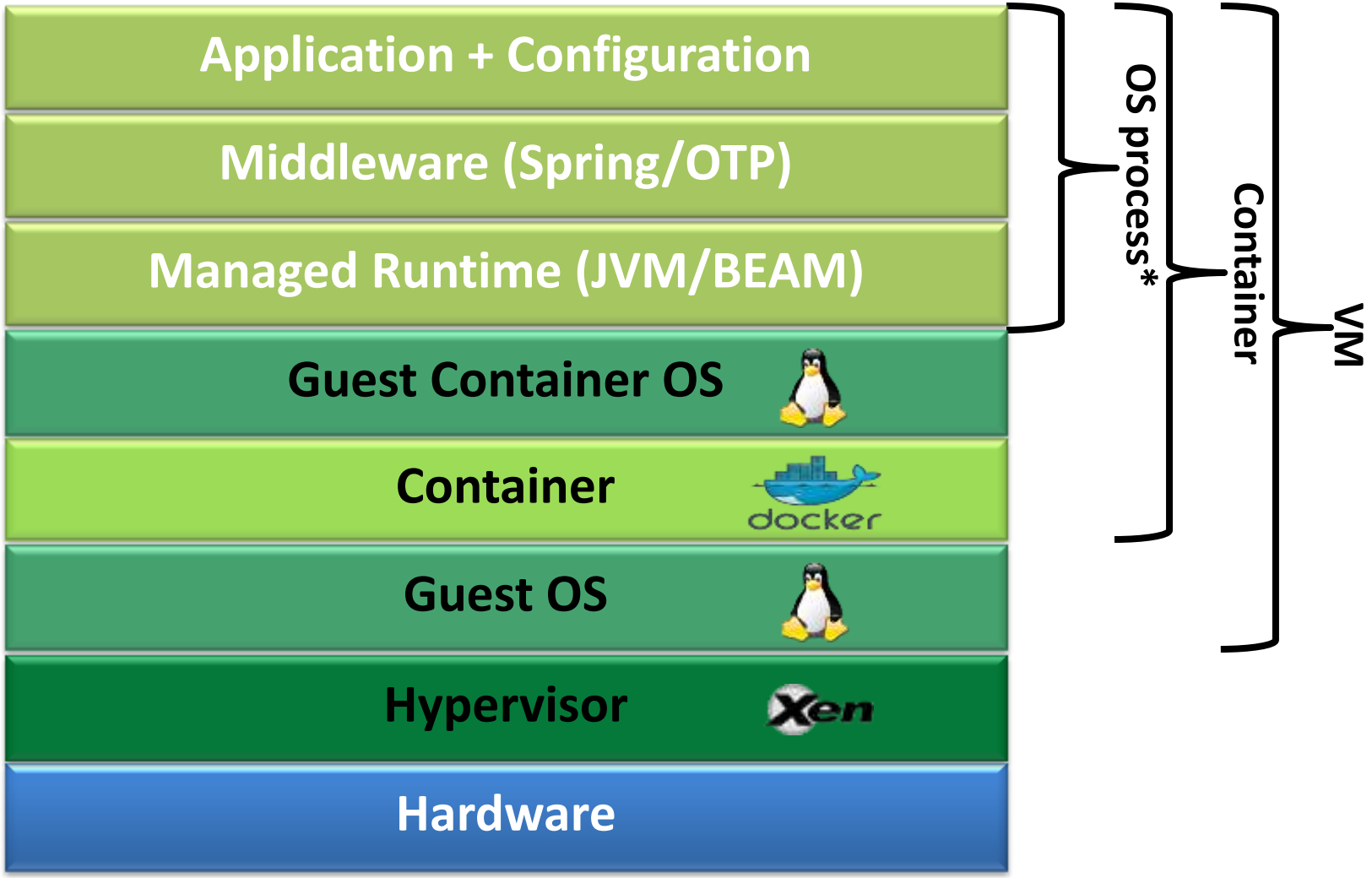
$ ls /dev
on Ubuntu 14.04 AWS EC2 instance



```
autofs          hvc6                null    ram8    tty15   tty33   tty51    ttyS10   ttyS29   vcsa
block           hvc7                port    ram9    tty16   tty34   tty52    ttyS11   ttyS3    vcsa1
btrfs-control   input               ppp     random  tty17   tty35   tty53    ttyS12   ttyS30   vcsa2
char            kmsg                psaux   rfkill  tty18   tty36   tty54    ttyS13   ttyS31   vcsa3
console         log                 ptmx    rtc     tty19   tty37   tty55    ttyS14   ttyS4    vcsa4
core            loop0               pts     rtc0    tty2    tty38   tty56    ttyS15   ttyS5    vcsa5
cpu             loop1               ram0    shm     tty20   tty39   tty57    ttyS16   ttyS6    vcsa6
cpu_dma_latency loop2               ram1    snapshot tty21  tty4    tty58    ttyS17   ttyS7    vcsa7
disk            loop3               ram10   snd     tty22   tty40   tty59    ttyS18   ttyS8    vga_arbiter
ecryptfs        loop4               ram11   stderr  tty23   tty41   tty6     ttyS19   ttyS9    xen
fd              loop5               ram12   stdin   tty24   tty42   tty60    ttyS2    uinput   xvda
full            loop6               ram13   stdout  tty25   tty43   tty61    ttyS20   urandom  xvda1
fuse            loop7               ram14   tty     tty26   tty44   tty62    ttyS21   vcs      xvdb
hpet            loop-control        ram15   tty0    tty27   tty45   tty63    ttyS22   vcs1     zero
hvc0            mapper              ram2    tty1    tty28   tty46   tty7     ttyS23   vcs2
hvc1            mcelog              ram3    tty10   tty29   tty47   tty8     ttyS24   vcs3
hvc2            mem                 ram4    tty11   tty3    tty48   tty9     ttyS25   vcs4
hvc3            net                 ram5    tty12   tty30   tty49   ttyprintk ttyS26  vcs5
hvc4            network_latency     ram6    tty13   tty31   tty5    ttyS0    ttyS27   vcs6
hvc5            network_throughput  ram7    tty14   tty32   tty50   ttyS1    ttyS28   vcs7
```

- 64 teletype devices?
- 32 serial ports?
- Sound?
- VGA?

“It’s DUPLICATED on so many LAYERS”



We run Single App per VM

We run in Single User mode

So, why do we still use Operating Systems designed for the desktops in the Cloud?

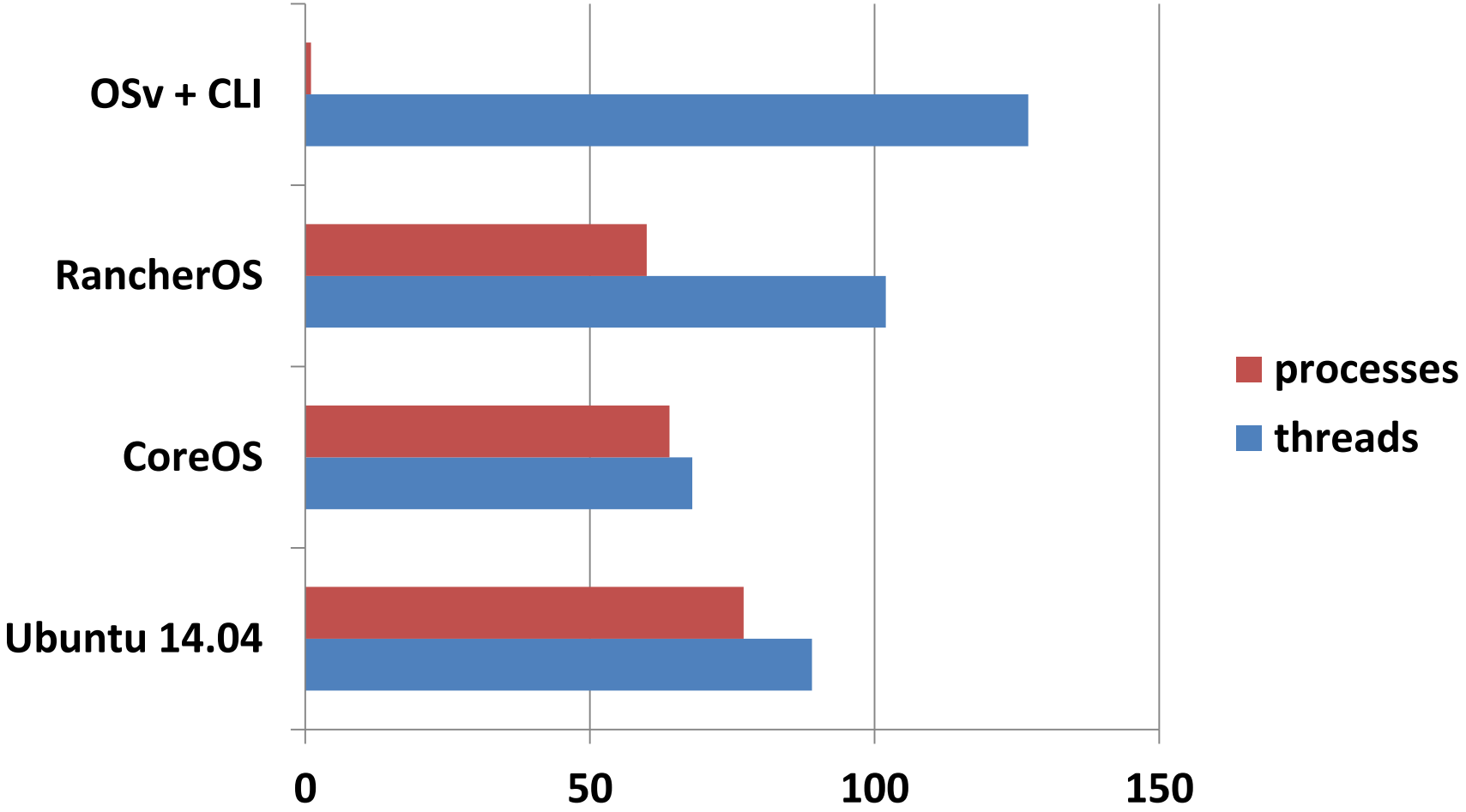
Minimalistic Linux OSes

- Embedded Linux versions
- DamnSmall Linux
- Linux with BusyBox

Min. Linux OSes for Containers
JeOS – "Just Enough OS"

- CoreOS
- RancherOS
- RedHat Project Atomic
- VMware Photon
- Intel Clear Linux
- Hyper

of Processes and Threads per OS



What's Unikernel?

- A "library" operating system
- kernel library linked with your application
- A kernel that can only support one process
- Single Address Space
- No kernel / userspace separation
- *"Unikernels: Library Operating Systems for the Cloud"*
  - **http://anil.recoil.org/papers/2013-asplos-mirage.pdf**

What's Anykernel?

- Programming discipline for kernel code reuse
- Capabilities
  - NetBSD filesystems as Linux processes
  - User-space TCP/IP stack
- *"The Design and Implementation of the Anykernel and Rump Kernels", Antti Kantee*
  - **http://book.rumpkernel.org/**

Unikernel + Anykernel

- Unikernels originally were designed to run on top of Hypervisor
- Now combined with Anykernel / Rump kernel ideas, some unikernels, like ***MirageOS*** or ***LING VM*** can run on ***Bare Metal***

A new Erlang platform – LING – runs directly on Xen.

You get less administrative headaches,
better security and performance.

LING is highly-compatible with Erlang/OTP.

LING understands .beam files.

Develop on Erlang/OTP –
deploy using LING.

http://erlangonxen.org/

LING: NOT JUST ERLANG ON XEN PORTING LING TO ARM/MIPS MICROCONTROLLERS

Viktor Sovietov

Embedded Erlang is Real

LING: NOT JUST ERLANG ON XEN PORTING LING TO ARM/MIPS MICROCONTROLLERS

LING is an Erlang platform with minimal requirements with respect to its software environment. Until now LING h
virtualised x86. We managed to port LING to ARM (Raspberry Pi). The port to PIC32 (MIPS) microcontrollers is in the
discusses the challenges of running Erlang on bare metal and potential benefits of LING/Erlang as a basis for embedde

Talk objectives:

- Describes of how Erlang is applicable to develop IoT and embedded applications;

Target audience:

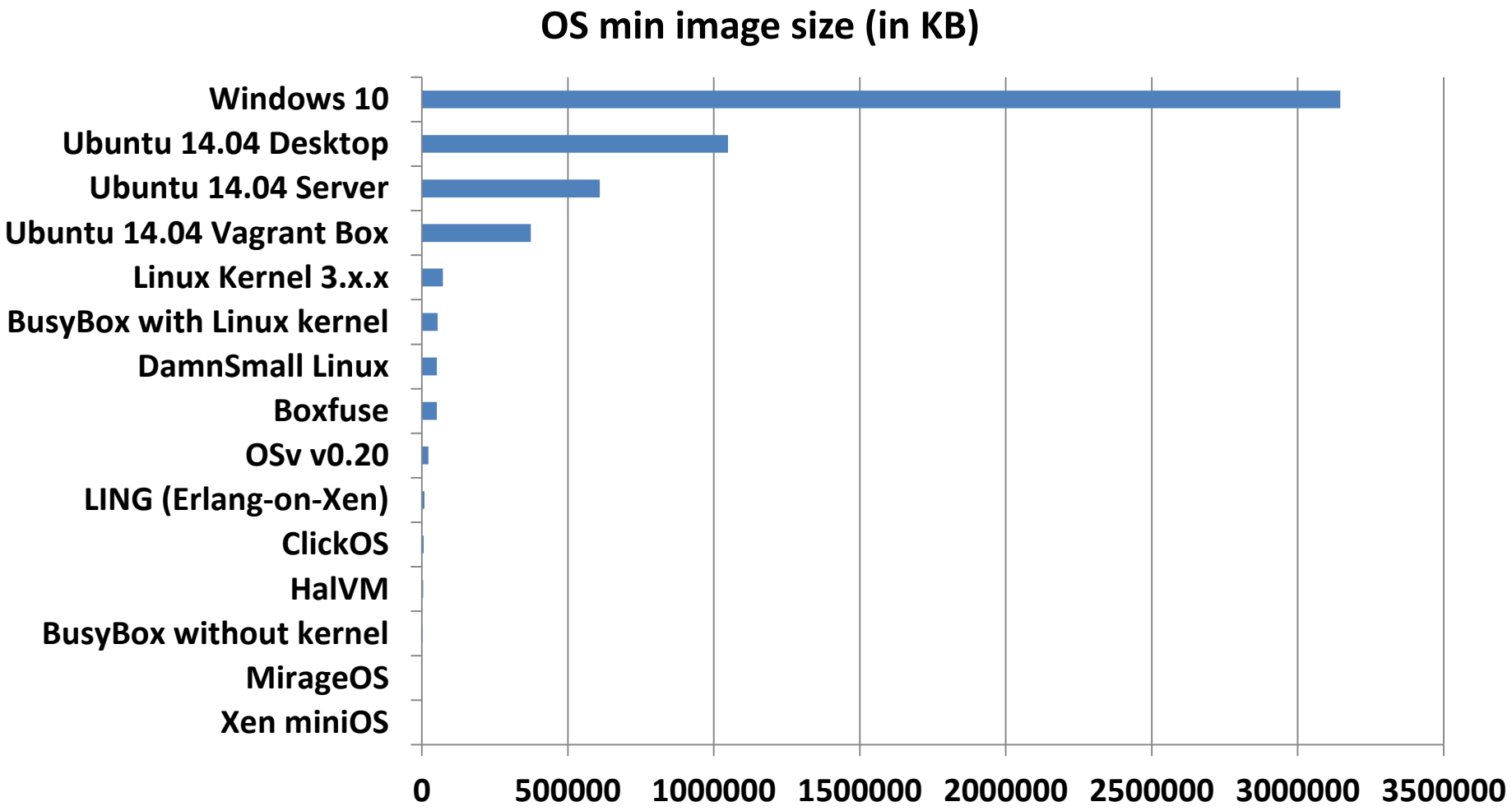
- Erlang developers, technology entrepreneurs interested in IoT;

http://osv.io

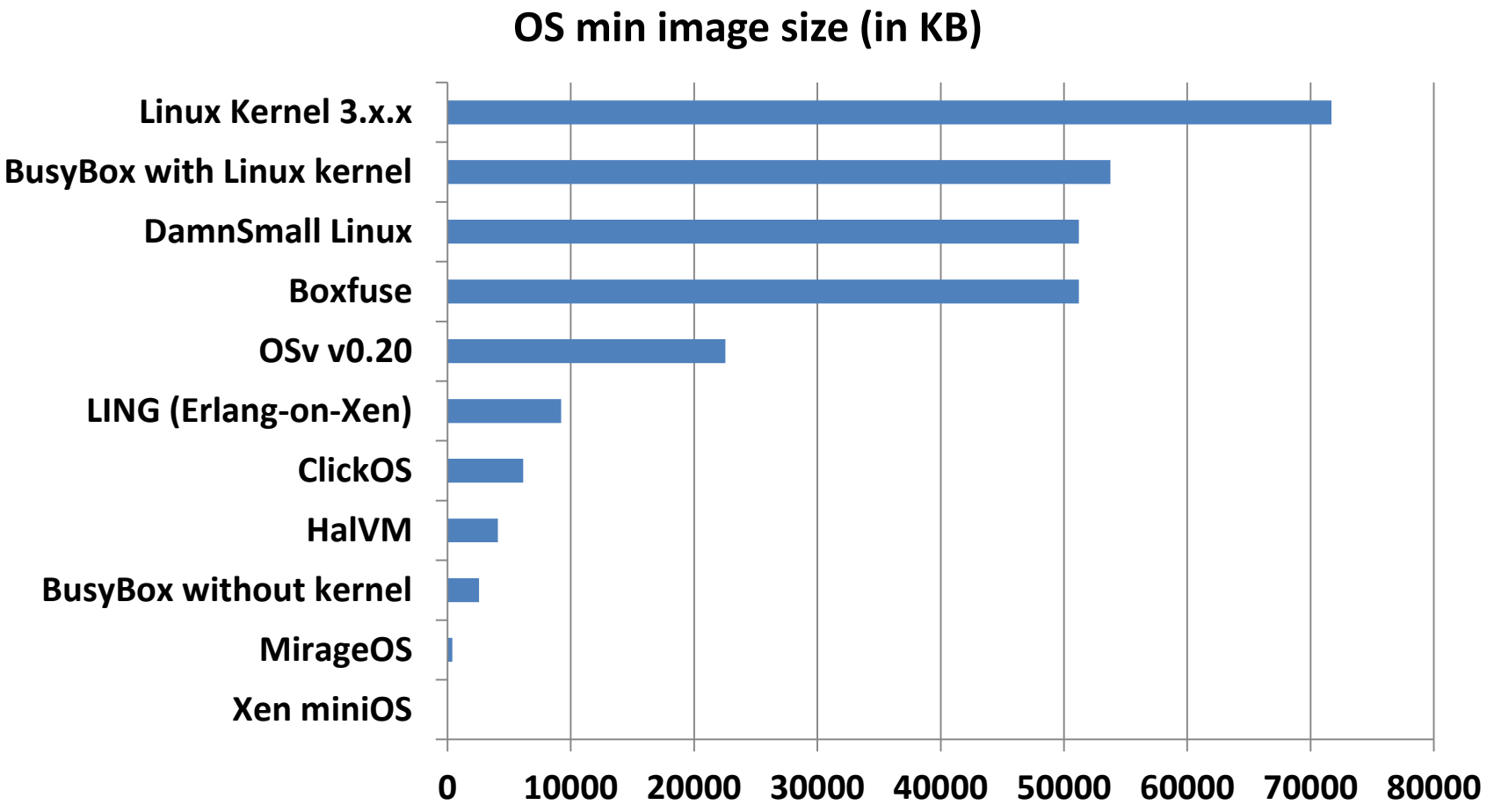
Unikernel Projects

Name	Target
MirageOS	OCaml
HalVM	Haskell
ClickOS	Click DSL
Clive	Go (inpired by Plan 9 & CSP)
Boxfuse	JVM
LING VM (Erlang-on-Xen)	Erlang
OSv	POSIX*, JVM, High Performance non-POSIX API

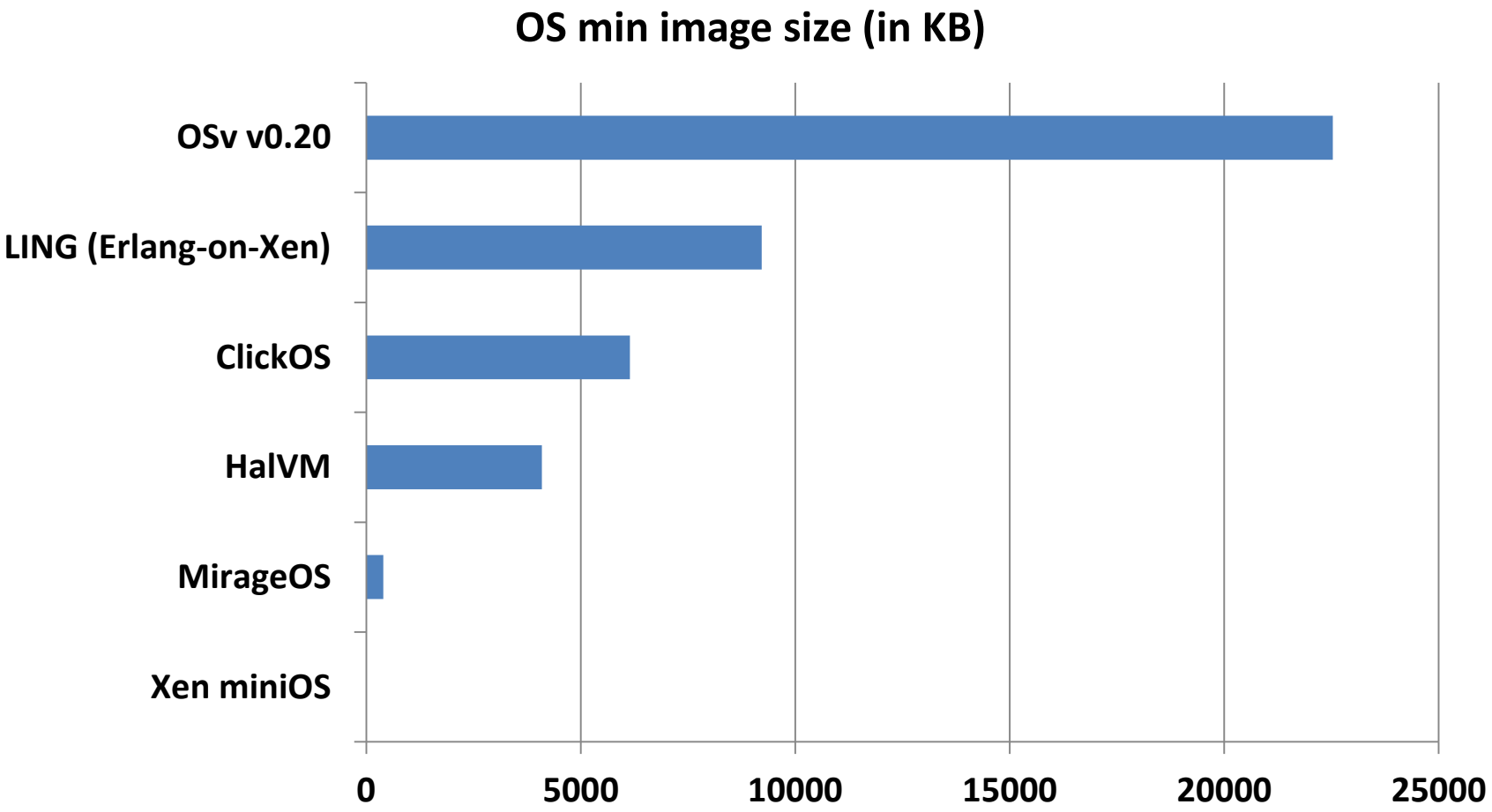
Windows 10 & Ubuntu – are outliers



OSv is half-way between Unikernels & Linux Kernel



OSv is “fat” Unikernel or Anykernel-ish



Specialized Unikernels

- Compile to very small images
- Boot very fast
- requires writing all code in Higher Level statically typed language, like OCaml or Haskell
- Hard/impossible to port existing code
- Very secure: tiny attack surface

Anykernel-ish or “fat unikernels”

- Larger images (still much smaller that Linux)
- Longer boot times (much faster that Linux)
- Larger attack surface (much smaller that Linux)
- Easier to port existing code
- More pragmatic solution overall


YO UNIKERNEL SO FAT
IT SEMI-LINUX ABI COMPATIBLE
imgflip.com

OSv is "Fat" Unikernel / Anykernel-ish

- OSv sometimes called ***"fat" unikernel***
- since OSv images are a little bit larger than for other unikernels
- It also called ***anykernel***-ish, since it
  - run on top of multiple hypervisors
  - provide TCP/IP stack and filesystem
- Small price to pay for LINUX ABI compatibility


CAPSTAN

Capstan – Docker-like CLI for OSv

```
$ capstan
NAME:   capstan - pack, ship, and run applications in light-weight VMs
USAGE: capstan [global options] command [command options] [args...]
VERSION: v0.1.8
COMMANDS:
info            show disk image information
import          import an image to the local repository
pull            pull an image from a repository
rmi             delete an image from a repository
run             launch a VM. You may pass the image name as the first arg
build           build an image
images, i       list images
search          search a remote images
instances, I    list instances
stop            stop an instance
delete          delete an instance
help, h         Shows a list of commands or help for one command
```

Download & Run OSv+CLI image

```
$ capstan run cloudius/osv
Downloading cloudius/osv/index.yaml...
170 B / 170 B [====================] 100.00 %
Downloading cloudius/osv/osv.qemu.gz...
21.78 MB / 21.78 MB [==============] 100.00 %
Created instance: cloudius-osv
OSv v0.20
eth0: 192.168.122.15
pthread_setcancelstate() stubbed
/#
```


ERLANG
Match
Made
in
Heaven

designed for the cloud

on

Containers add new layers
vs
Unikernels remove layers

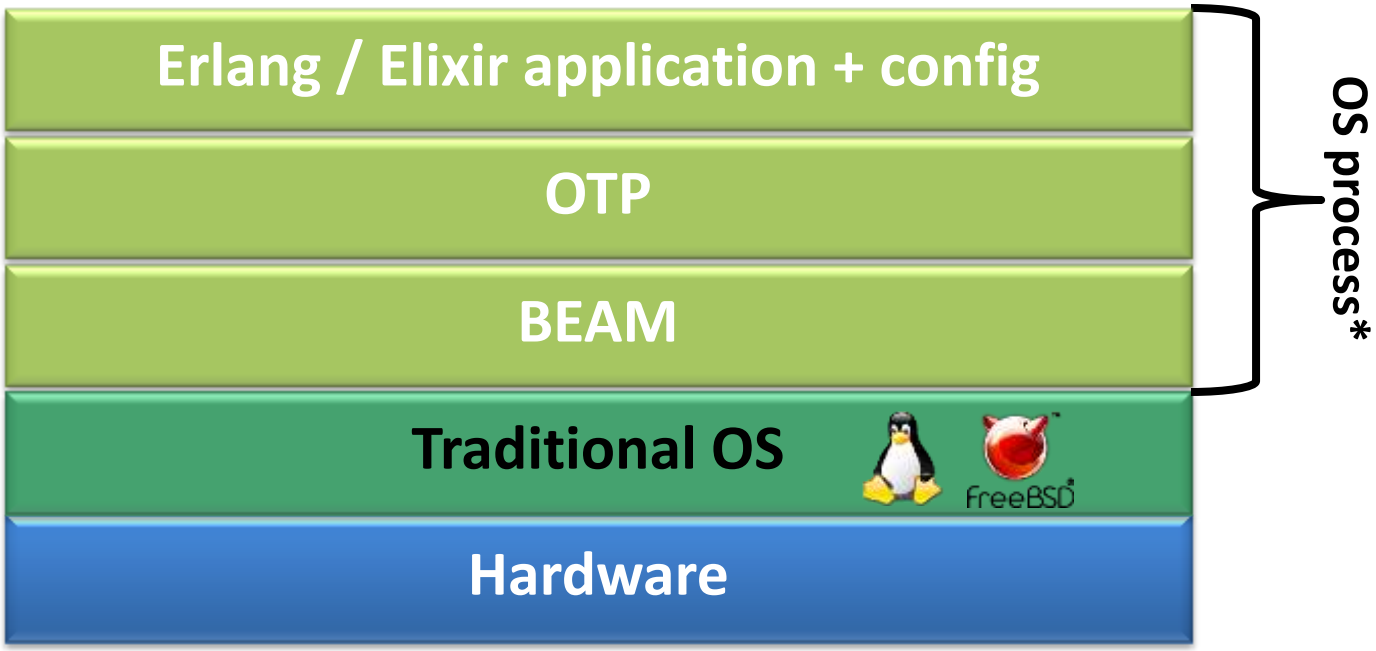
"All problems in computer science can be solved by another level of indirection“

– David Wheeler

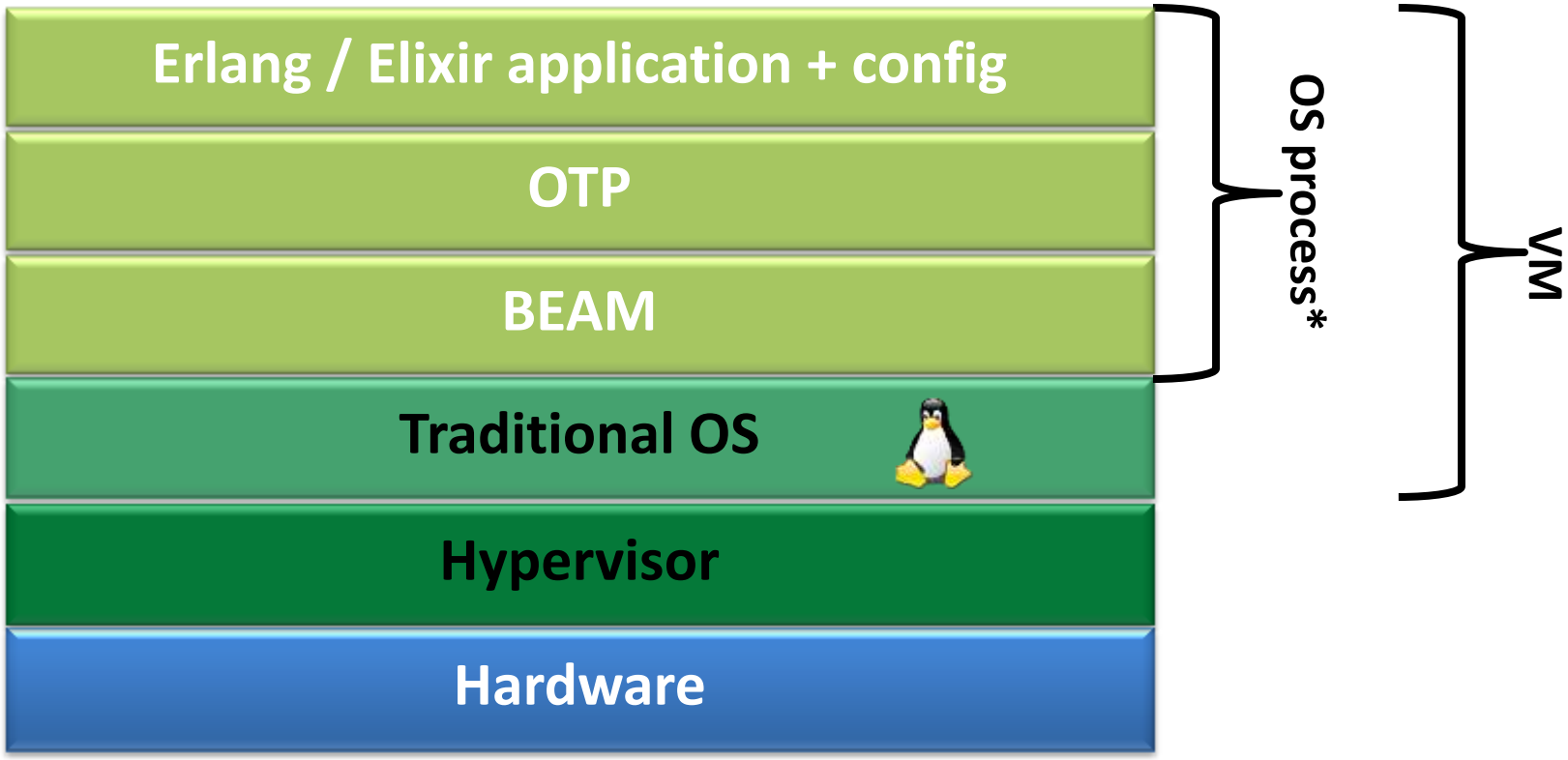
"...except for the problem of too many layers of indirection.“

– Kevlin Henney

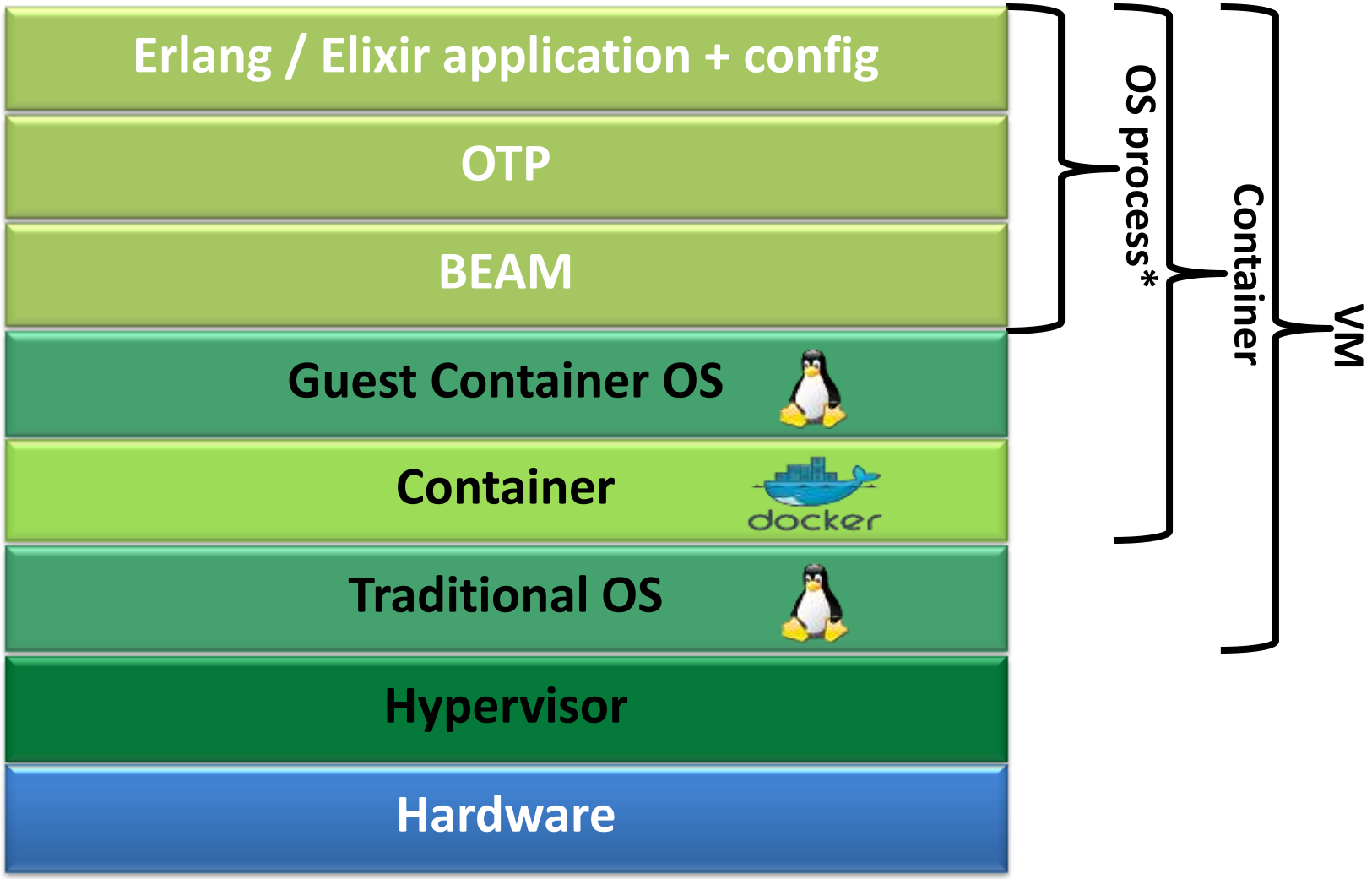
Erlang on Physical Machine



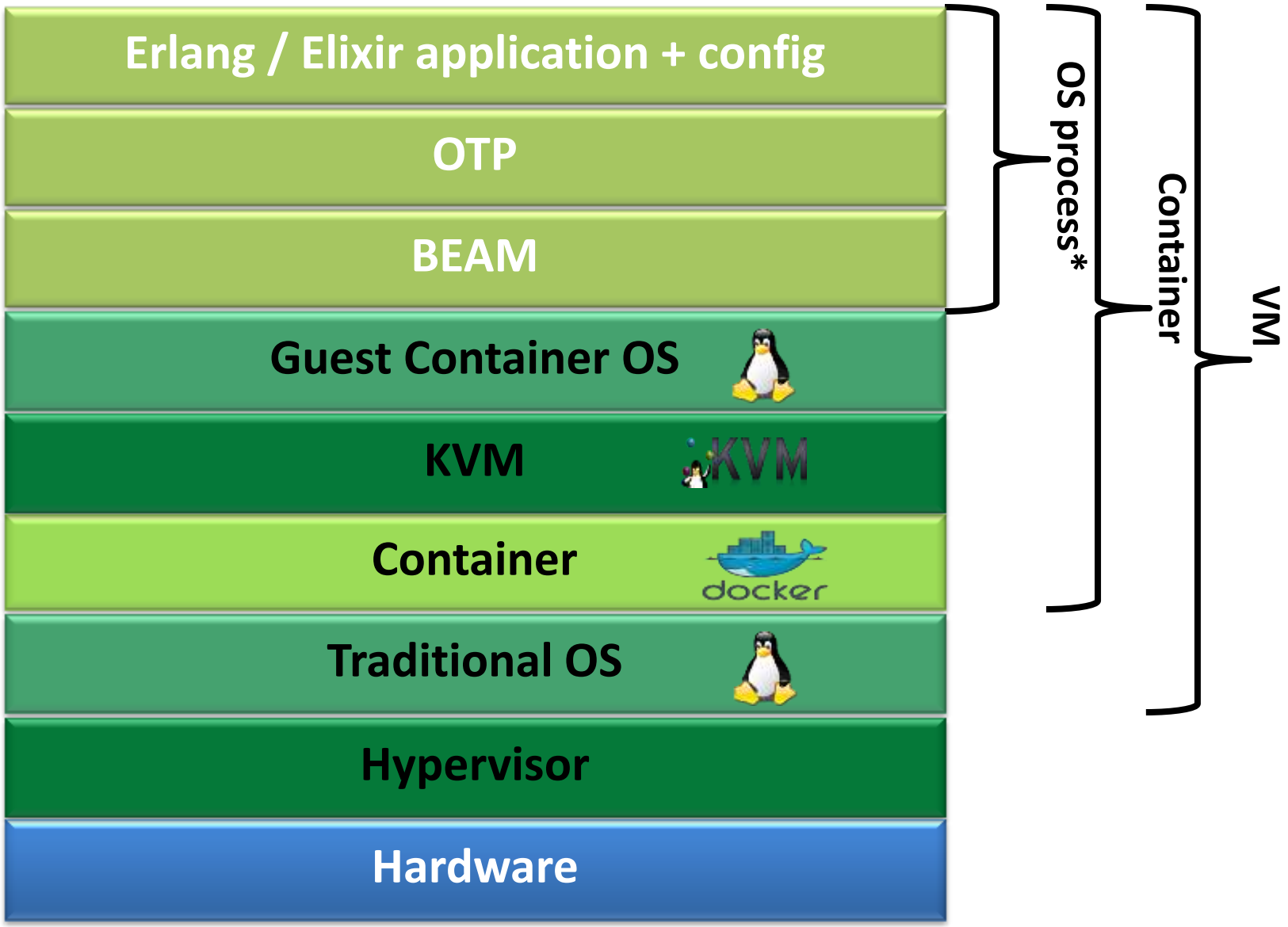
Erlang in VM


Erlang / Elixir application + config
OTP
BEAM
Traditional OS
Hypervisor
Hardware
OS process*
VM

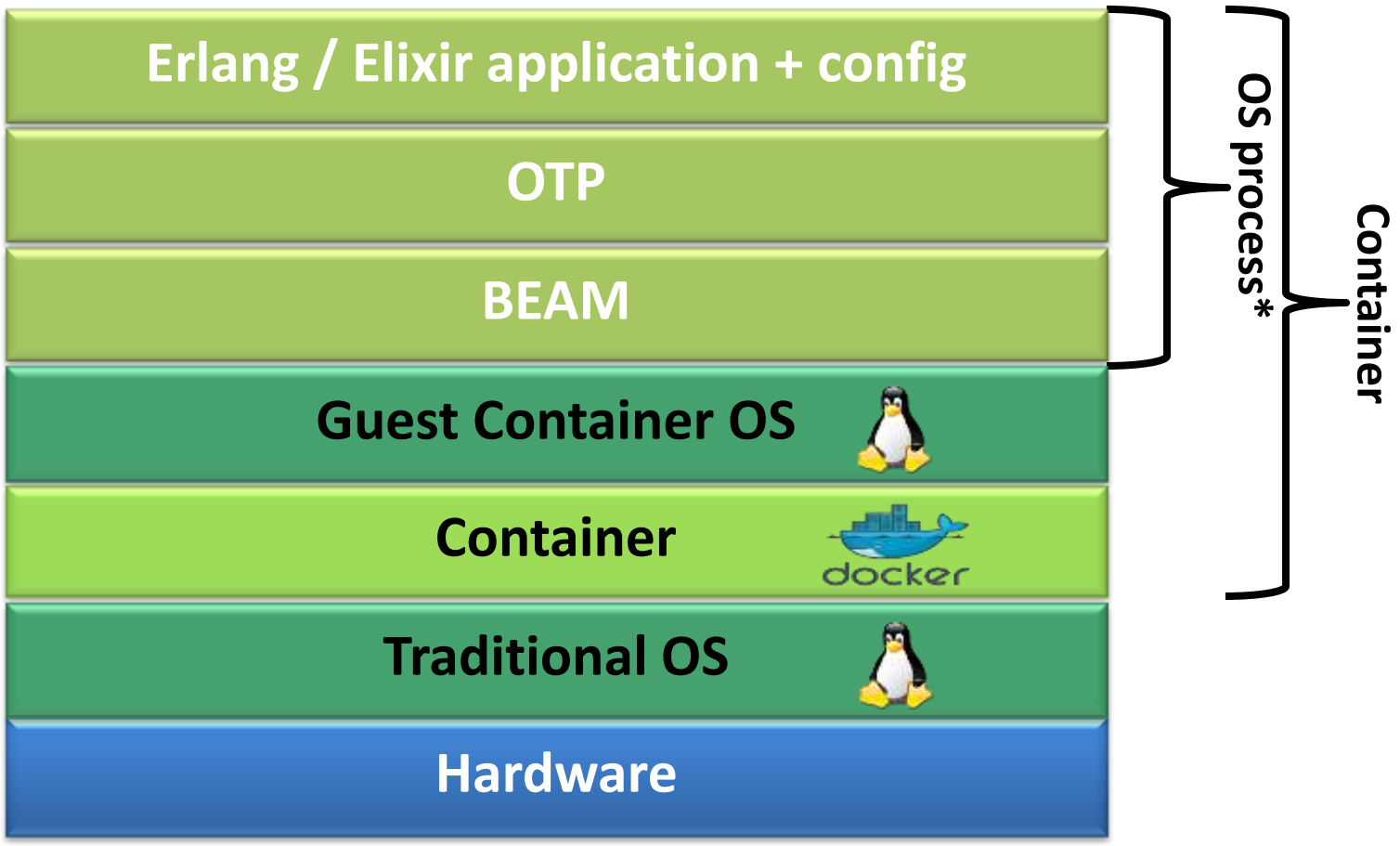
Container in VM



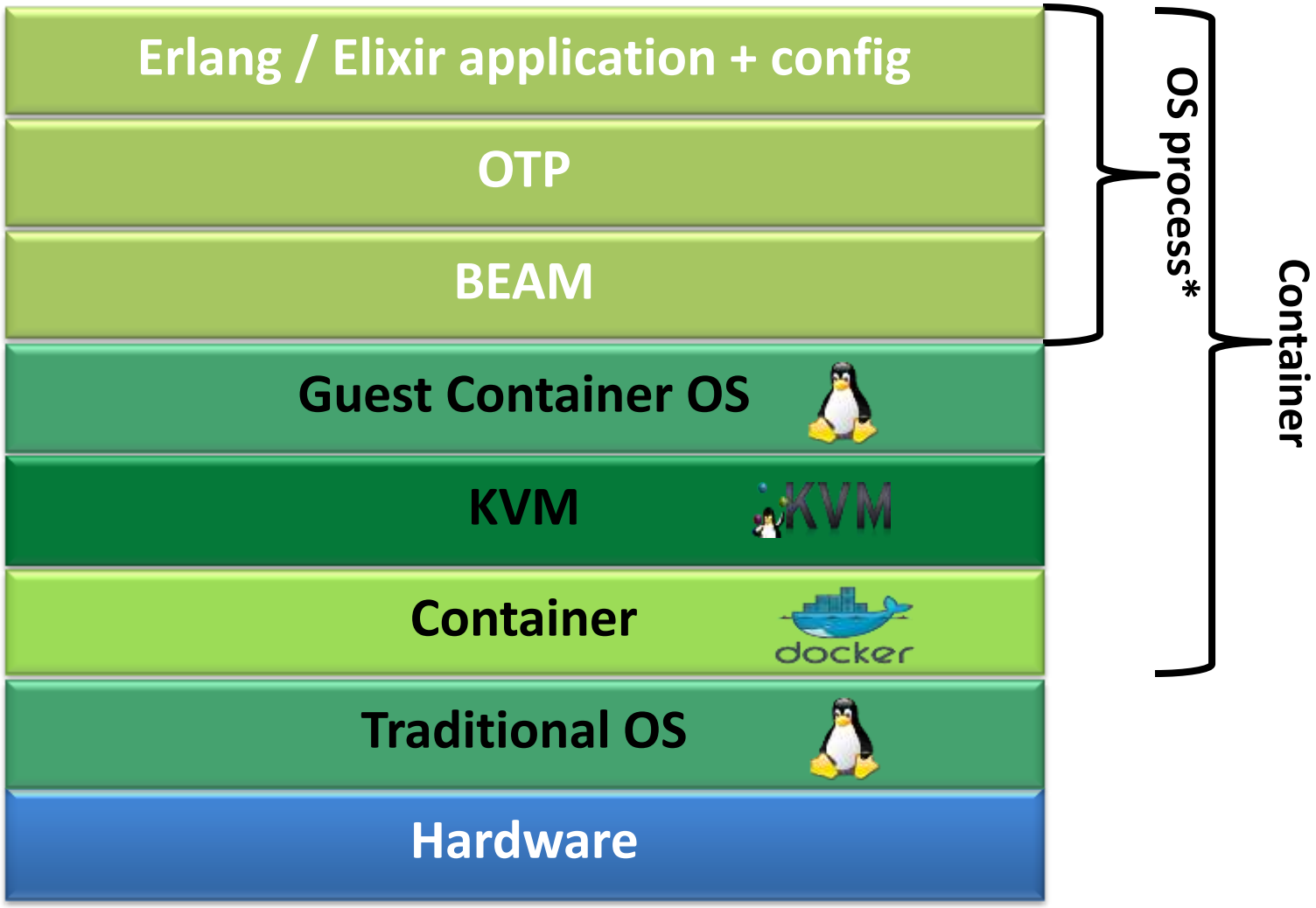
Multitenancy - VM in Container in VM



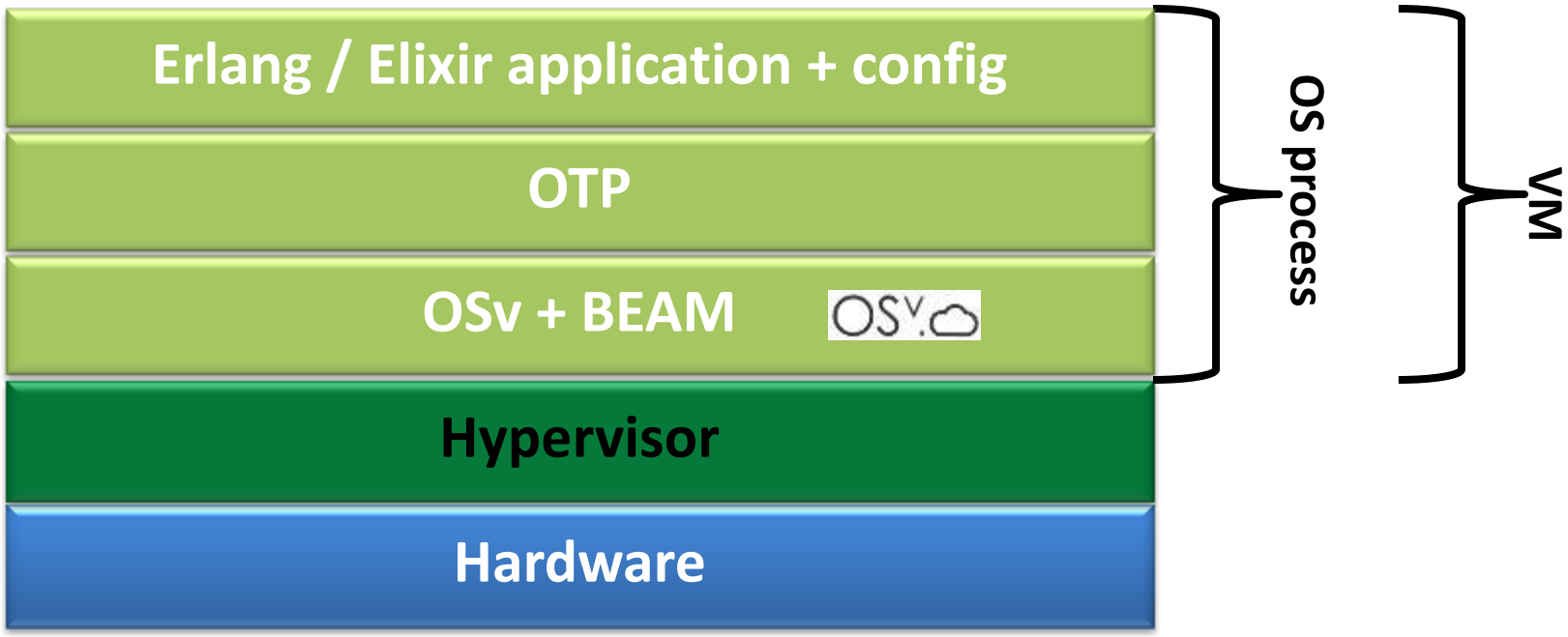
Container on Physical Machine



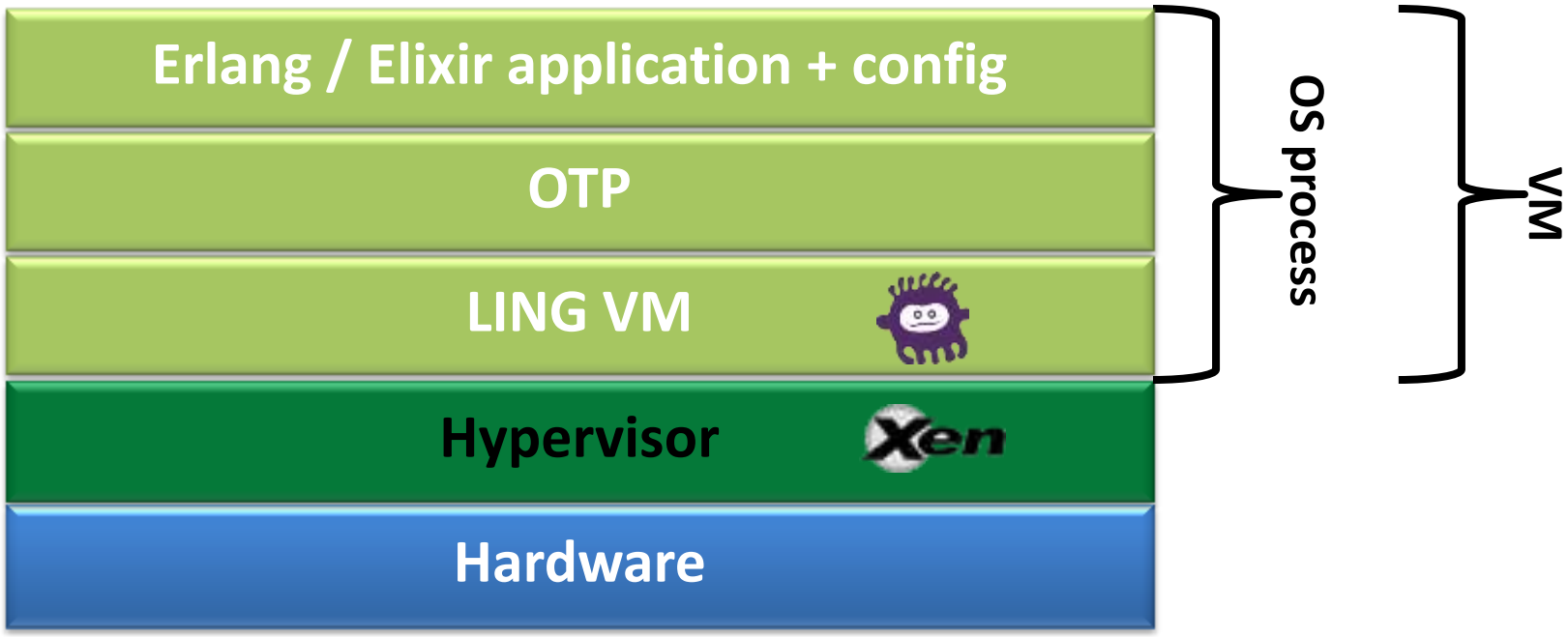
VM in Container on Physical Machine for Multitenancy



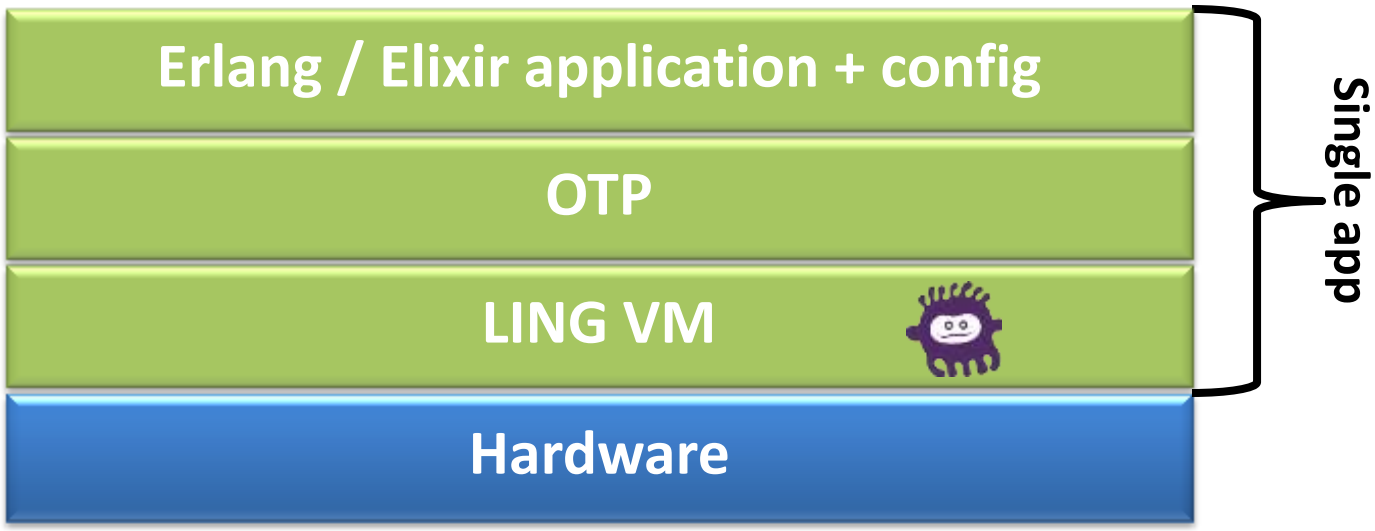
Erlang on OSv



LING VM (Erlang-on-Xen)



LING VM on Bare Metal



Erjang – Erlang for JVM

OSv has built-in JVM support, so Erjang runs w/o porting on OSv

on

OSv

```
$ git clone https://github.com/cloudius-systems/osv-apps
$ cd osv-apps/erjang/
$ cat Capstanfile
base:
    cloudius/osv-openjdk
cmdline: >
    /java.so -jar erjang.jar
build:
    ./GET
files:
    /erjang.jar: erjang.jar
```

Build OSv+JVM+Erjang image

```
$ cd osv-apps/erjang
$ capstan build
Building erjang...
Downloading cloudius/osv-
openjdk/index.yaml...
169 B / 169 B [==================] 100.00 %
Downloading cloudius/osv-openjdk/osv-
openjdk.qemu.gz...
74.26 MB / 74.26 MB [============] 100.00 %
Uploading files...
1 / 1 [=========================] 100.00 %
```

```
$ capstan run
Created instance: erjang
OSv v0.20
eth0: 192.168.122.15
** Erjang R16B01 **  [root:/~resource]
[erts:5.10.2] [smp S:1 A:10]
[java:1.7.0_55] [unicode]
WARNING: fcntl(F_SETLK) stubbed
Eshell V5.10.2  (abort with ^G)
1> lists:reverse("Hello, World!").
"!dlroW ,olleH“
2> q().
ok
```

Porting a C/C++ application to OSv

- A single-process application
  - may not ***fork()*** or ***exec()***
- ***Position Independent Code***
  - recompile with ***–fPIC*** flag
- Need to rebuild as a shared object (***.so***)
  - link with ***–shared*** flag
- or as ***Position Independent Exec.*** (***– fpie***)
  - can run the same executable in Linux and OSv

Build OSv image with native code

hello.cc

```
1 #include <iostream>
2
3 int main()
4 {
5     std::cout << "Hello, world!" << std::endl;
6 }
```

Makefile

Makefile

```
1  CXXFLAGS  = -g -Wall -std=c++11 -fPIC $(INCLUDES)
2
3  TARGET = hello
4
5  OBJ_FILES = hello.o
6
7  quiet = $(if $V, $1, @echo " $2"; $1)
8  very-quiet = $(if $V, $1, @$1)
9
10 all: $(TARGET).so
11
12 %.o: %.cc
13     $(call quiet, $(CXX) $(CXXFLAGS) -c -o $@ $<, CXX $@)
14
15 $(TARGET).so: $(OBJ_FILES)
16     $(call quiet, $(CXX) $(CXXFLAGS) -shared -o $(TARGET).so $^, LINK $@)
17
18 clean:
19     $(call quiet, rm -f $(TARGET).so $(OBJ_FILES), CLEAN)
```

```
$ git clone https://github.com/cloudius-systems/capstan-example
$ cd capstan-example
$ cat Capstanfile
base:
    cloudius/osv-base
cmdline:
    /tools/hello.so
build:
    make
files:
    /tools/hello.so: hello.so
```

Build OSv + Hello World image

```
$ capstan build
Building capstan-example...
Downloading cloudius/osv-base/index.yaml...
154 B / 154 B [==================] 100.00 %
Downloading cloudius/osv-base/osv-base.qemu.gz...
20.13 MB / 20.13 MB [============] 100.00 %
Uploading files...
1 / 1 [========================] 100.00 %
```

Run OSv + Hello World image

```
$ capstan run
Created instance: capstan-example
OSv v0.20
eth0: 192.168.122.15
Hello, world!

$
```

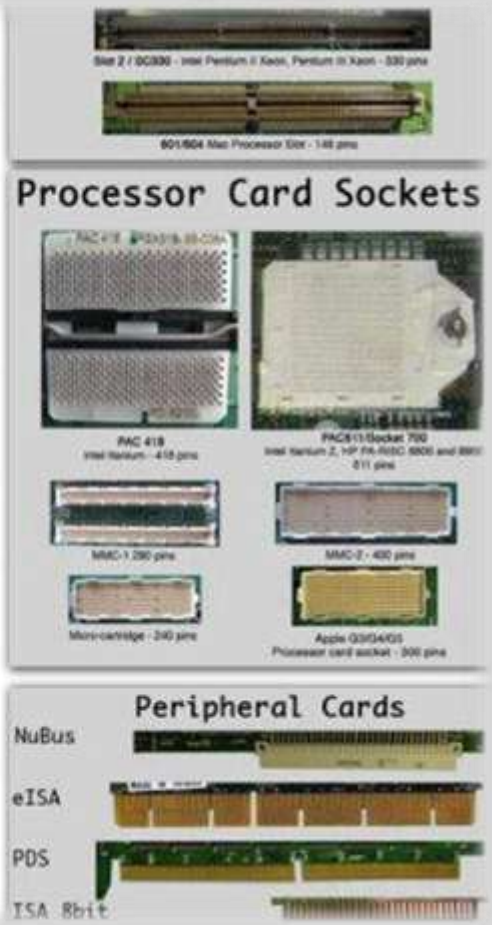
PORTING ERLANG/OTP TO OSV

No fork() or exec() in OSv

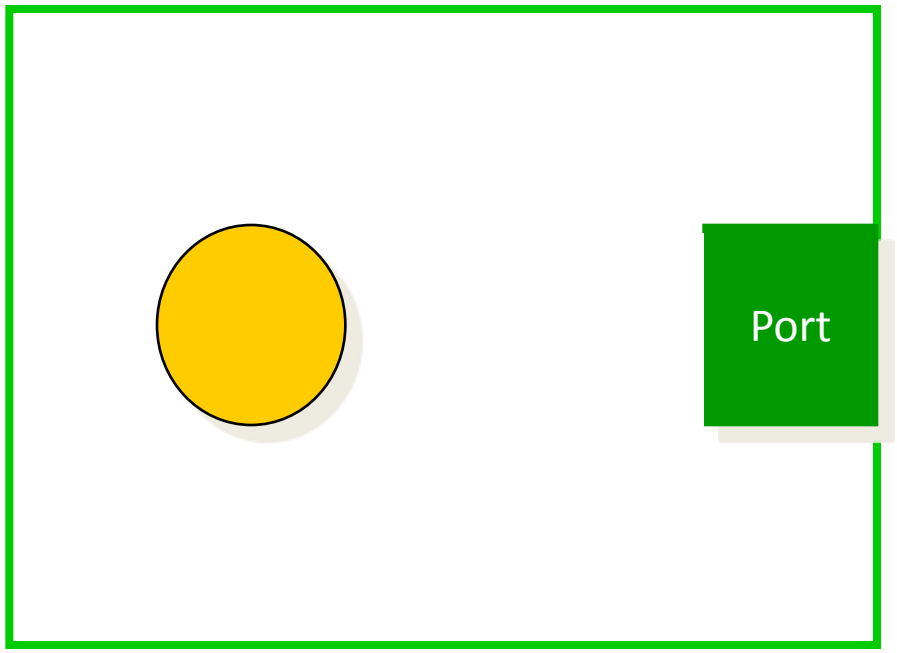
BUT ERLANG ♥ FORKS

ERLANG ♥ PORTS

Interfaces to the outside world

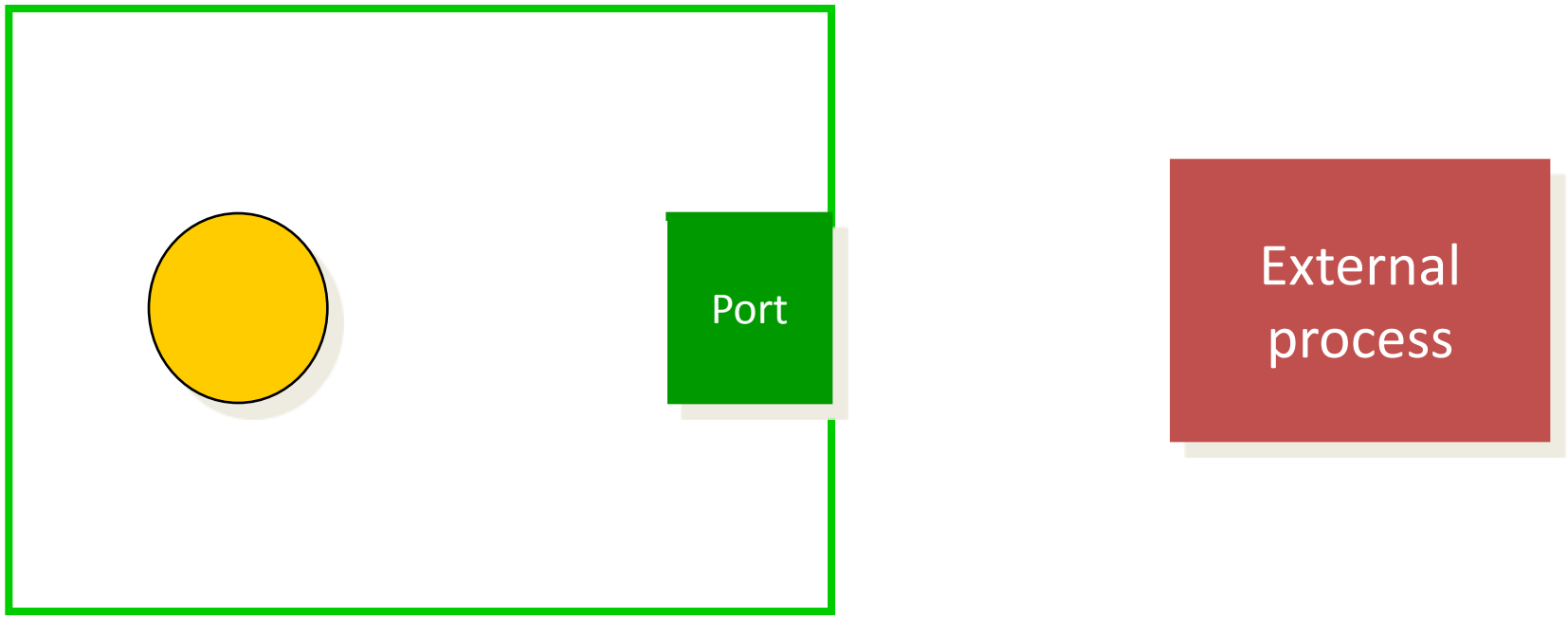


How Erlang Port Works



Source: Ulf Wiger, Erlang Open Telecom Platform

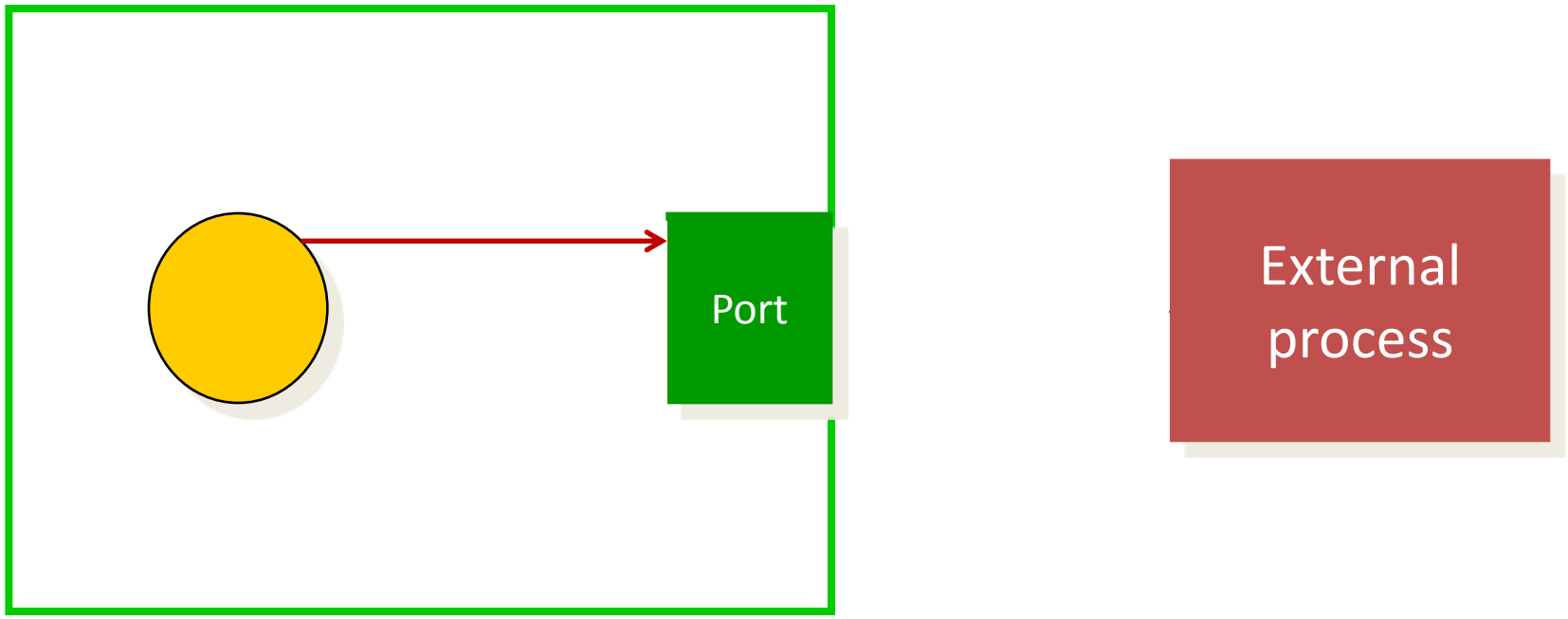
fork/exec the Port executable



- **Erlang Port is a middleman/proxy**
- **allow External Process to pretend to be another Erlang process**
- **allow communication using regular Erlang message passing**

Source: Ulf Wiger, Erlang Open Telecom Platform

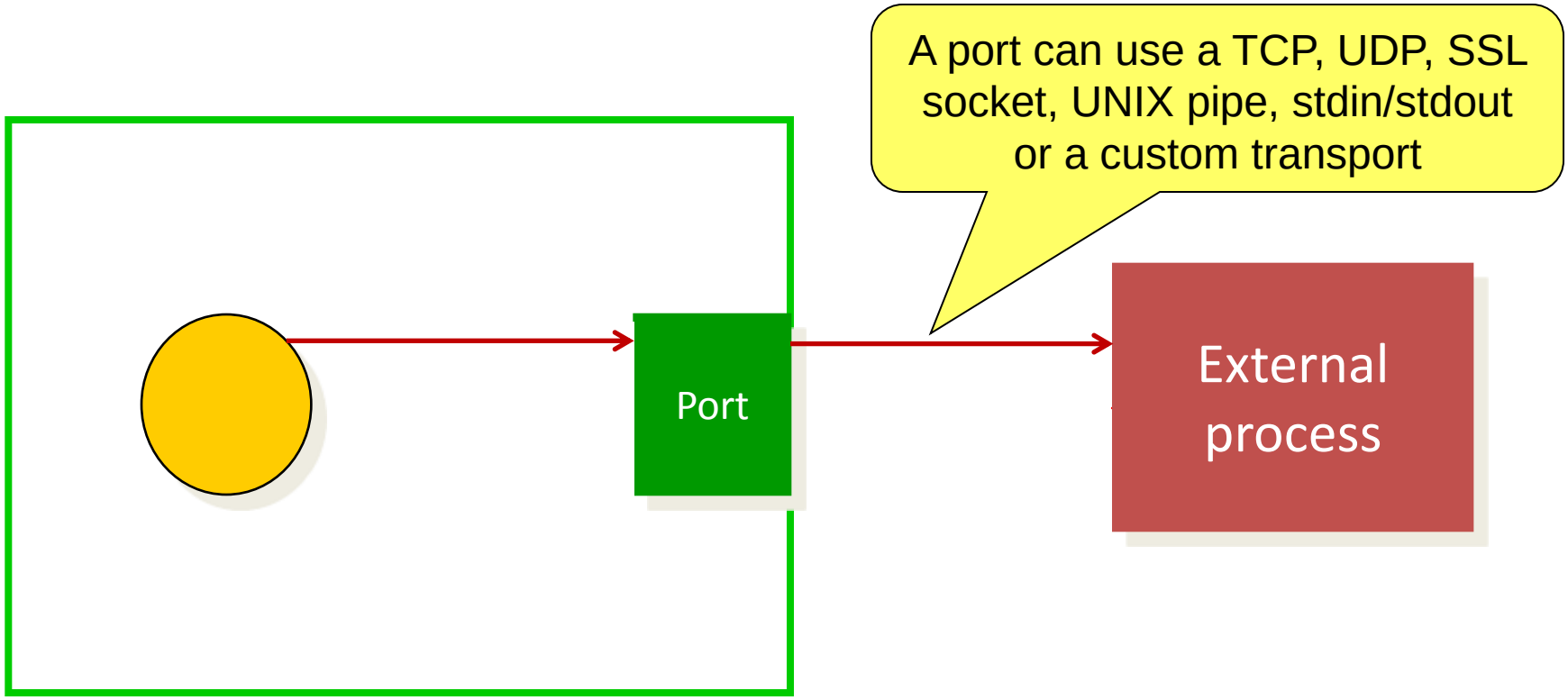
Erlang process sends msg to Port




```
Port ! {self(), {command, [1,2,3]}}
```

Source: Ulf Wiger, Erlang Open Telecom Platform

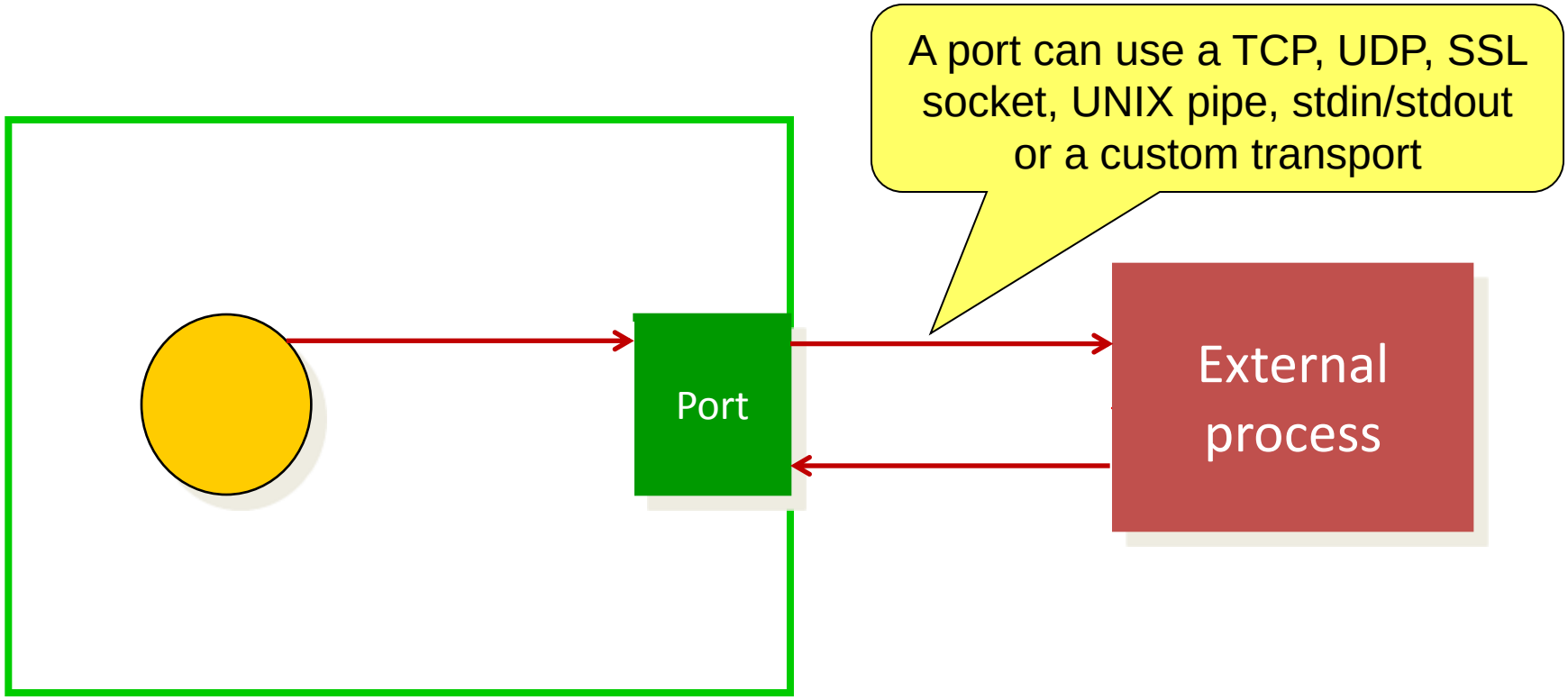
Port send msg to External Process




```
Port ! {self(), {command, [1,2,3]}}
```

Source: Ulf Wiger, Erlang Open Telecom Platform

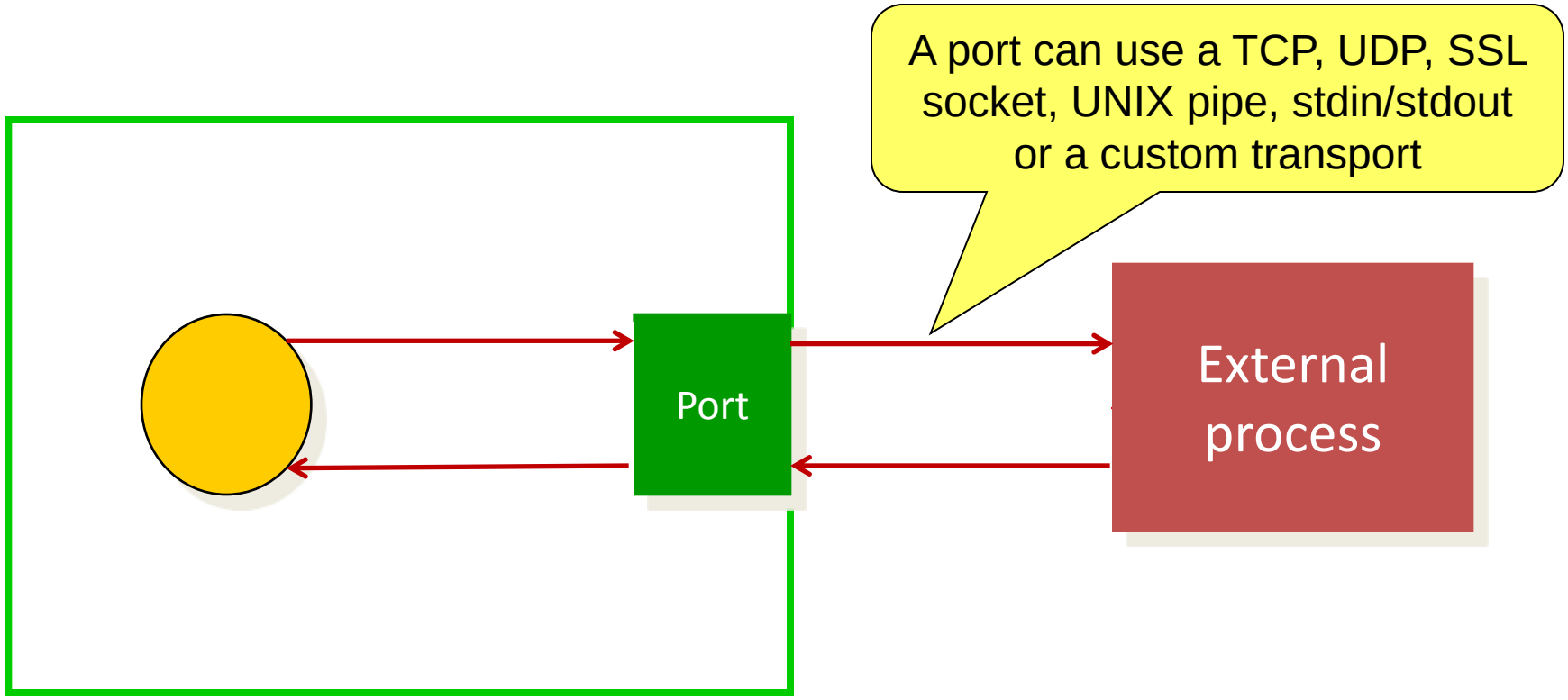
Port rcv msg from External Process




```
receive
    {Port, {data, Info}} -> ...
end
```

Source: Ulf Wiger, Erlang Open Telecom Platform

Port send msg to Erlang Process




```
receive
    {Port, {data, Info}} -> ...
end
```

Source: Ulf Wiger, Erlang Open Telecom Platform

Communicating with External world

	In-proc	Out-of-proc
Custom Protocol	Linked-in Drivers (.so)	Ports (executable)
Generic Protocol	NIF (.so)	C-Node

How to adapt existing Ports to OSv?

- the solution was to rewrite ***ports using OS processes*** as ***linked-in drivers using POSIX threads***
- With some exceptions

Community Port of Erlang/OTP to OSv

- I “ported” Erjang a year ago
- I tried to port official Erlang/OTP – but quickly lost myself in the Erlang build system
- Also reliance on lots of ports and shell scripts made me less optimistic
- Cloudius is focused on POSIX & JVM, Erlang is not a priority
- OSv mailing list discussed porting Erlang, but there was no progress

Community Port of Erlang/OTP to OSv

- People who helped to port Erlang to OSv:
  - ***Yao Zheng*** (https://github.com/bhuztez)
  - ***Bas Wegh***
  - ***Zika L'Etrange***

- ***Yao Zheng*** also ported following projects:
  - Elixir
  - LFE
  - yaws

The most important executables

- **erl** – wrapper to ***erlexec***
- ***epmd*** – Erlang Port Mapping Daemon
- ***heart*** – watchdog for BEAM (Erlang VM)
- ***inet_gethost*** – native DNS resolver
- ***osmon*** – mem_sup, disk_sup, cpu_sup
- ***odbserver*** – ODBC connectivity

Not everything ported yet.

erl

- ***erlang.so*** – OSv shared object executable wrapper to ***erlexec***
- pass ***vm.args*** and other CLI parameters to ***erlexec***

```
base:
  cloudius/osv-base
cmdline: >
   /usr/lib64/erlang.so -env HOME / \
   /etc/erlang/vm.args \
   /etc/default/erlang/vm.args
build:
  make
files:
  /usr/lib64/erlang/: ROOTFS/usr/lib64/erlang/
  /usr/lib64/erlang.so: erlang.so
  /etc/default/erlang/: default/
  /etc/erlang/vm.args: default/vm.args
```

```
$ git clone git@github.com:cloudius-systems/osv-apps.git
$ cd osv-apps/erlang
$ capstan run
Created instance: erlang
OSv v0.20
eth0: 192.168.122.15
sched_getaffinity() stubbed
Eshell V6.2  (abort with ^G)
1> lists:reverse("Hello, World!").
"!dlroW ,olleH“
2>
```

epmd – Erlang Port Mapper Daemon

- Option #1: **use pure Erlang implementation of epmd**:
  - https://github.com/bhuztez/noepmd
  - https://github.com/lemenkov/erlpmd
- Option #2: **run unmodified epmd as a POSIX thread, instead of a separate OS process**
- Option #2 was selected
- Q: How epmd implemented in Erlang on RTOS like VxWorks or OSE?

inet_gethost

```
1> httpc:request("http://google.com").

=ERROR REPORT==== 6-May-2015::13:43:58 ===
Error in process <0.79.0> with exit value:
{unknown,[{erlang,open_port,[
  {spawn,"inet_gethost 4"}, ...
```

inet_gethost

- Currently DNS resolving doesn't work yet
- Force to use pure-Erlang inet DNS resolver
- It's a default in Erlang on VxWorks & ENEA OSE RTOS

erl -sname mynode -kernel inetrc "'./erl_inetrc'"

{edns, 0}

http://www.erlang.org/doc/apps/erts/inet_cfg.html

osmon

- ***osmon*** ports where modified from standalone OS processes to linked-in drivers on POSIX thread
  - ***disk_sup***
  - ***mem_sup***
  - ***cpu_sup*** is disabled since OSv lacks required API

OTP apps + BEAM + Unikernel OS

- Erlang releases already may include Erlang VM
- Adding a Unikernel OS to releases will make them truly self-contained!

Development process

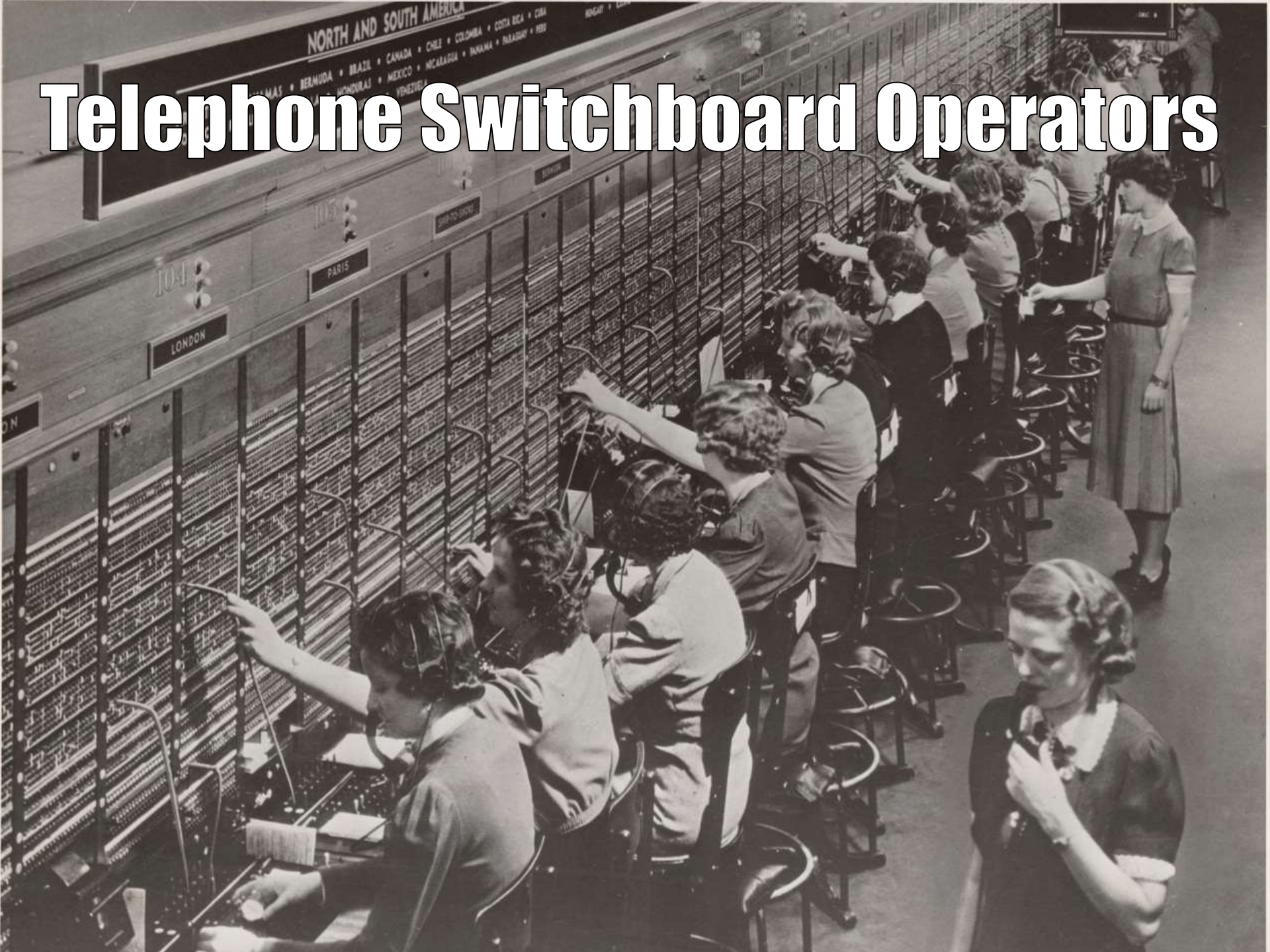
- Currently Erlang OSv port is developed as a set of patch files in Osv apps repo:
  - https://github.com/cloudius-systems/osv-apps
- Better way would be, to make it part of Erlang/OTP official repository (just like other OSes):
  - https://github.com/erlang/otp

Why Use Unikernels?

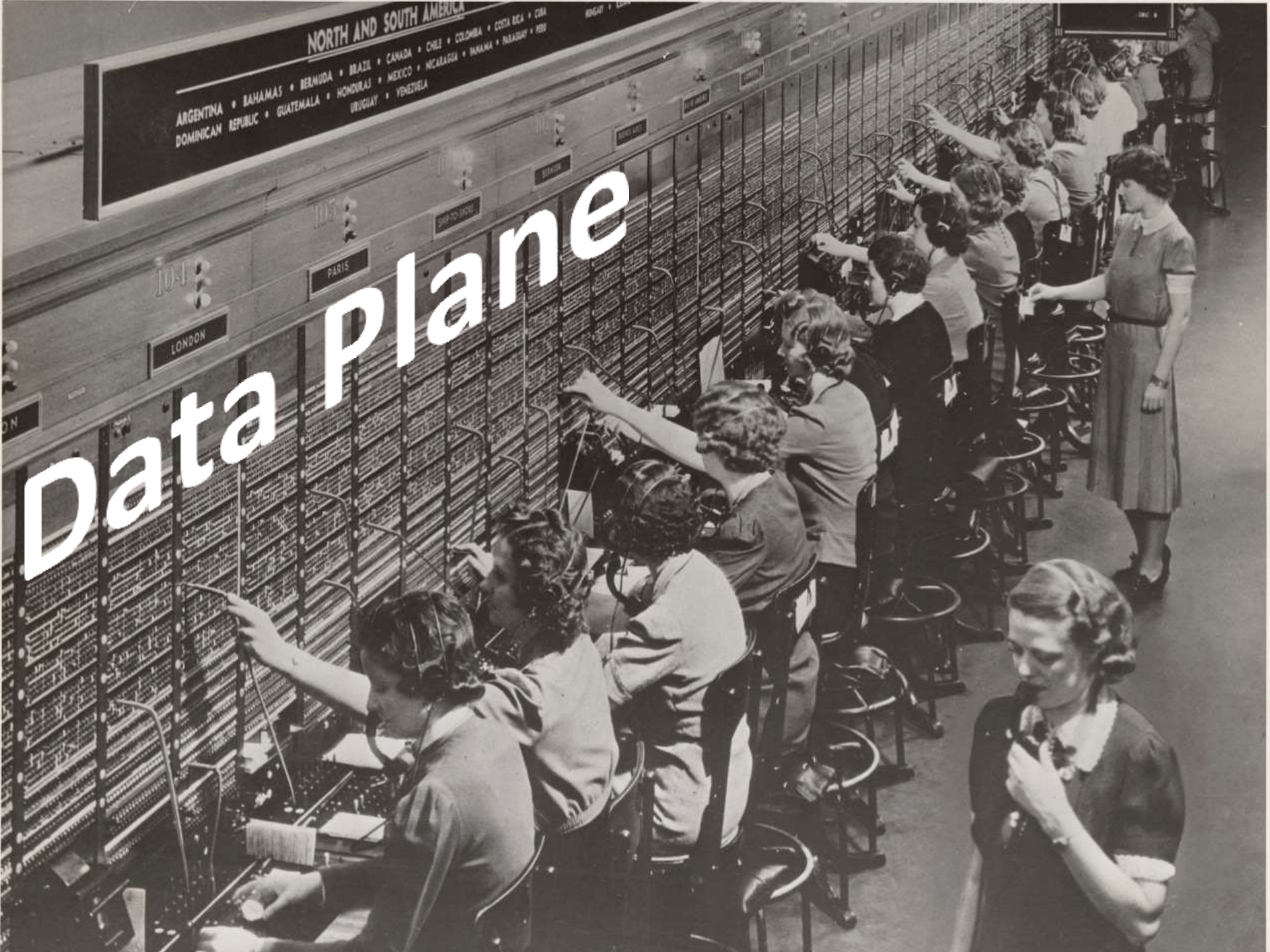
HIGH PERFORMANCE NETWORKING & I/O

Long time before “Erlang The Movie”. Telephone Switches worked like …

Source: Erlang The Movie



Telephone Switchboard Operators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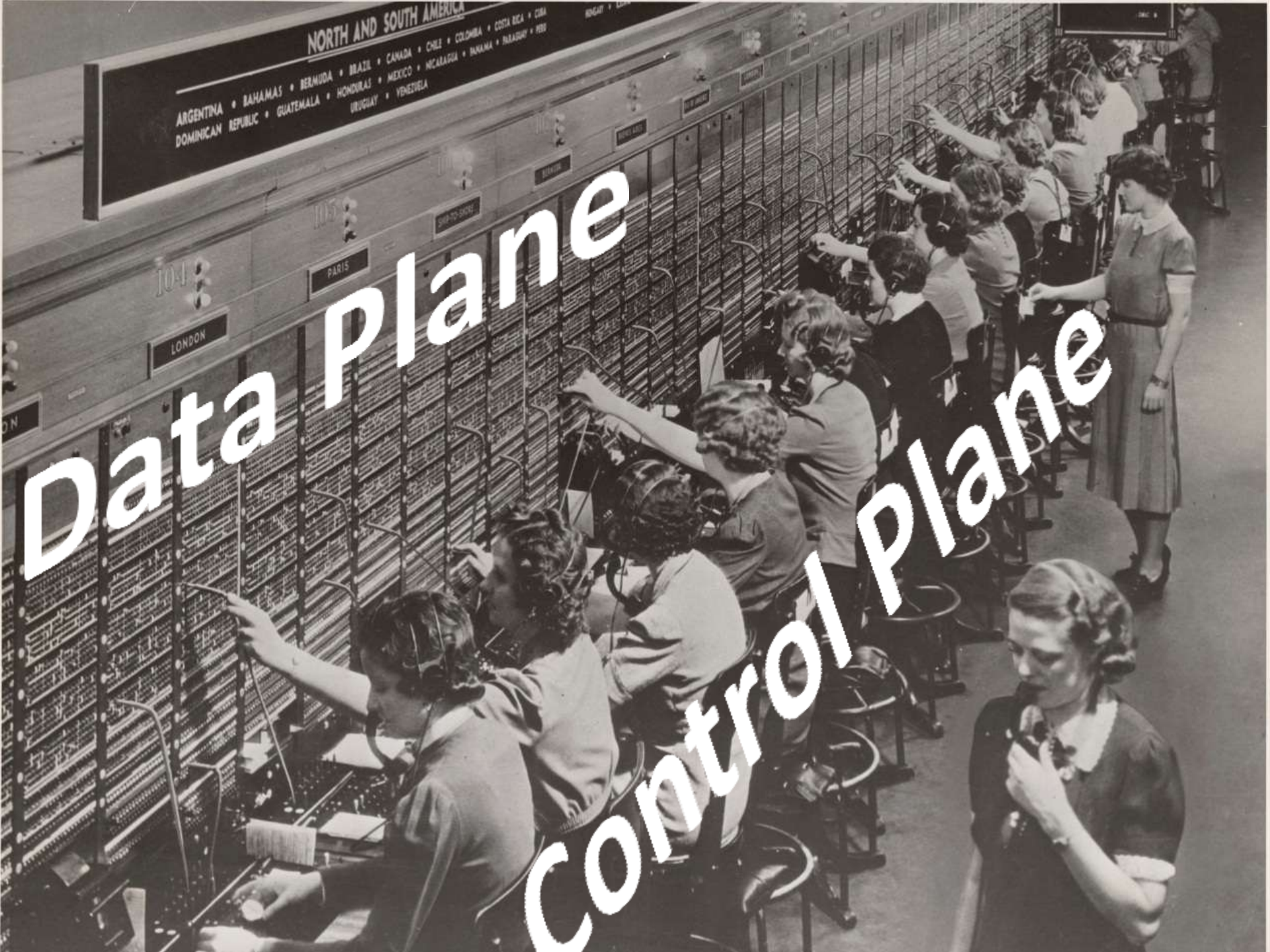
URUGUAY • VENEZUELA
PARIS
LONDON
Data Plane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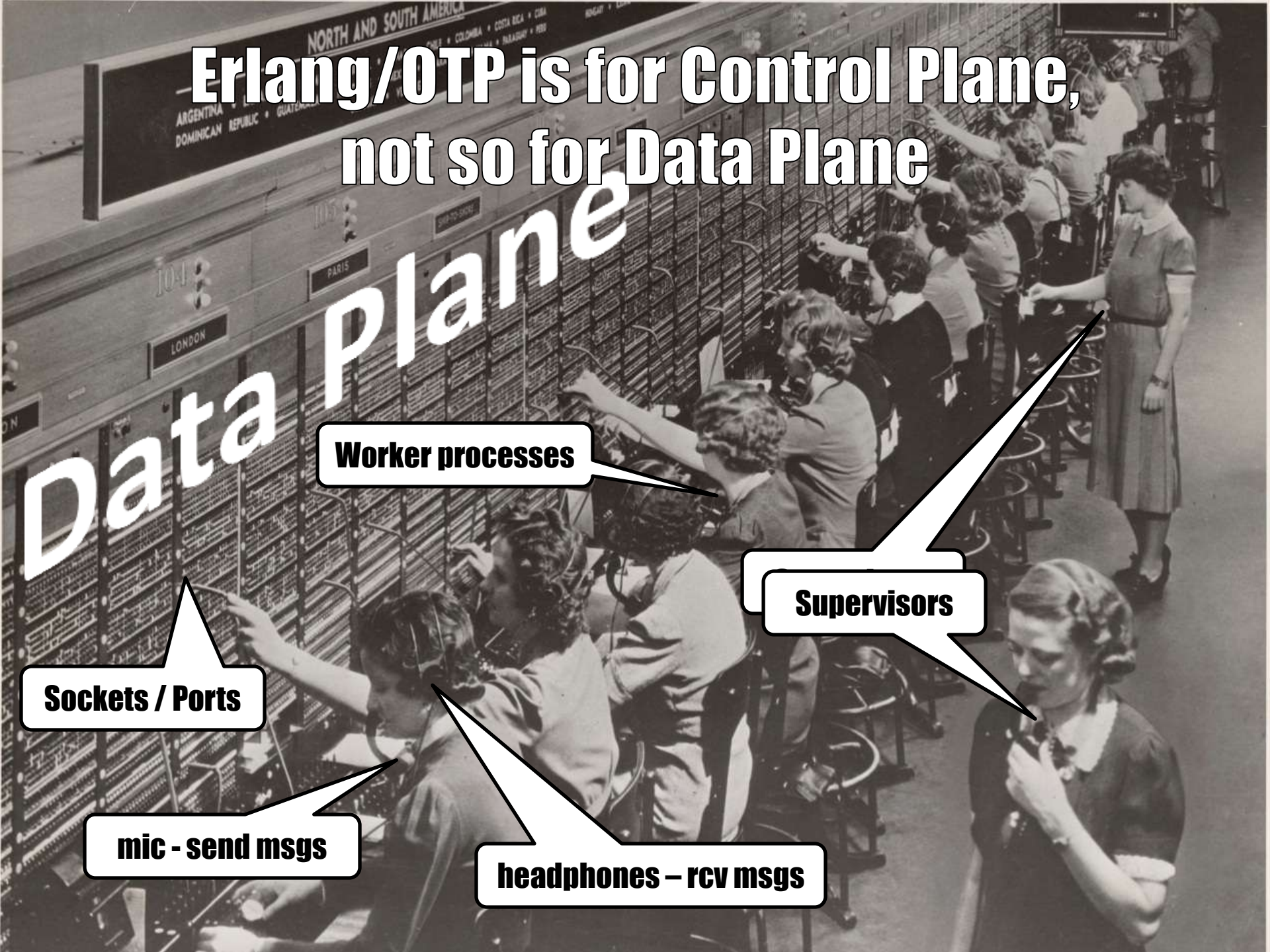
Data Plane
Control Plane


Erlang/OTP is for Control Plane,
not so for Data Plane
Data Plane
Worker processes
Supervisors
Sockets / Ports
mic - send msgs
headphones – rcv msgs

Data Plane was handled by specialized HW or RTOS

Source: Erlang The Movie

Erlang for both Control Plane + Data Plane

- In traditional OSes like Linux, IO & TCP/IP implemented via system calls, which has kernel overhead
- Many unikernels have user-space TCP/IP stack
- OSv implements Van Jacobsen's Net Channels optimization
- Optimized for NFV & SDN applications
- It's already possible with LING VM
- And with non-POSIX APIs in OSv

Comparing LING VM vs Erlang on OSv

LING VM (ERLANG-ON-XEN) & ERLANG/OTP ON OSV

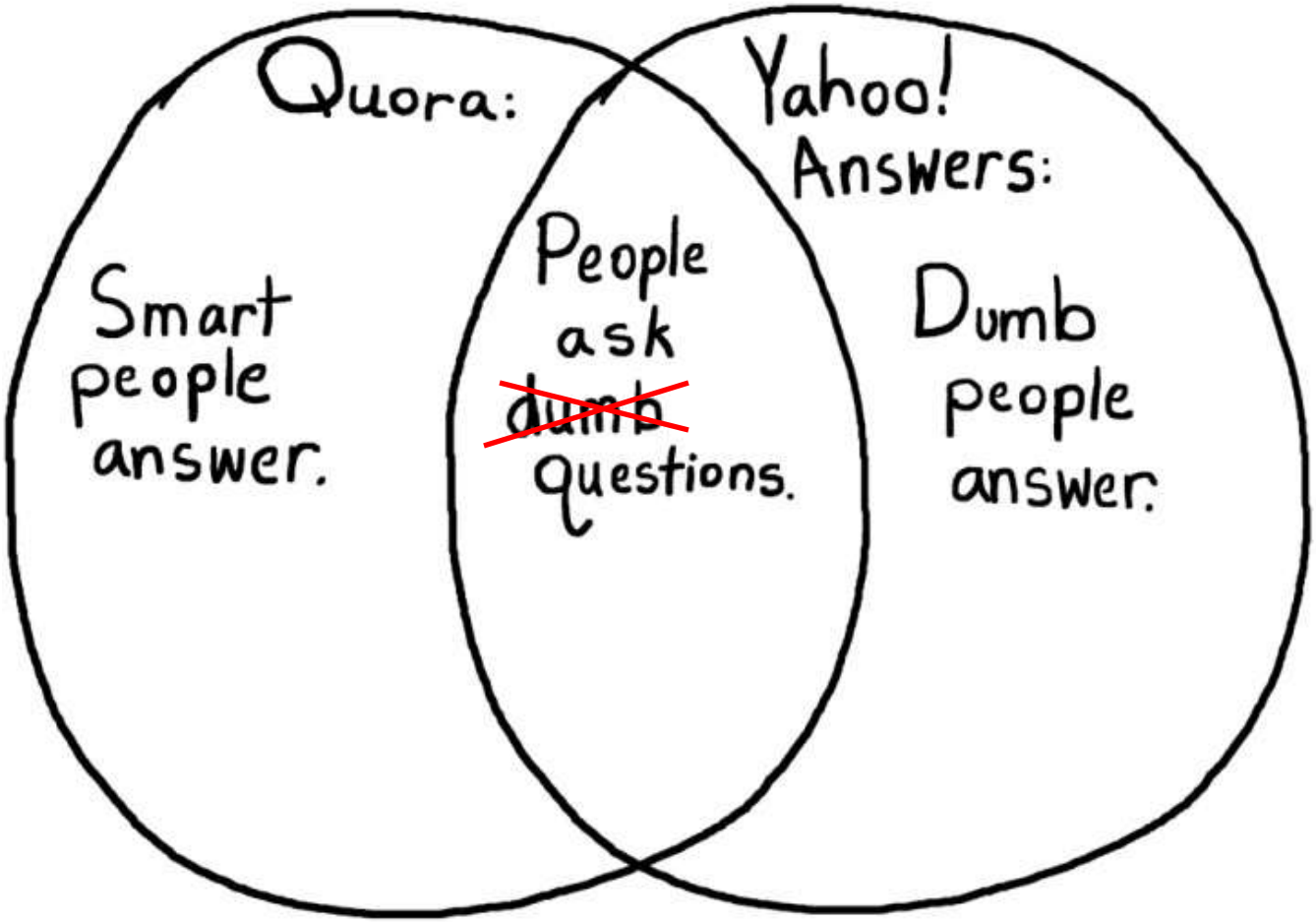
	LING VM	Erlang/OTP on OSv
Erlang VM	Custom VM (latest update from OTP – Dec 2014)	Open-source Ericsson's Erlang/OTP 17.3
State	Production-grade (3.5 years)	Experimental (6 months)
License	Proprietary (Sleepycat)	BSD 3-clause
SMP (threads)	-	+
Ports	-	-
NIFs	-	+ (possible)
HiPE	-	- (possible)
Disable GC	+	-

	LING VM	Erlang/OTP on OSv
Hypervisors	Xen only	KVM, Xen, VMware, VirtualBox
Bare Metal	Bare Metal RPi	-
Cloud support	AWS	AWS, Google Cloud
Machine arch	X86-64,ARM-32,MIPS	X86-64, ARM-64
Min img size	9 MB	21 MB
Boot time	< 100 ms	< 1 sec
“Dockerfiles”	railing	Capstan
Filesystem	goofs (Plan9)	ZFS
Data Plane, NFV/SDN	+	possible via non-POSIX API?

Takeaways

- It seems that after Docker, ***Unikernels*** are going to be ***"The Next Big Thing"*** in the Cloud
- Erlang community now has 2 options to choose from - try both
- Try to port your Erlang app
- Try to port your favorite open-source Erlang app
- But only when it makes sense!
- Download OSv and help porting standard Ericsson's Erlang/OTP to Osv
- **https://github.com/erlang-on-osv**

Thank You! Now Q&A



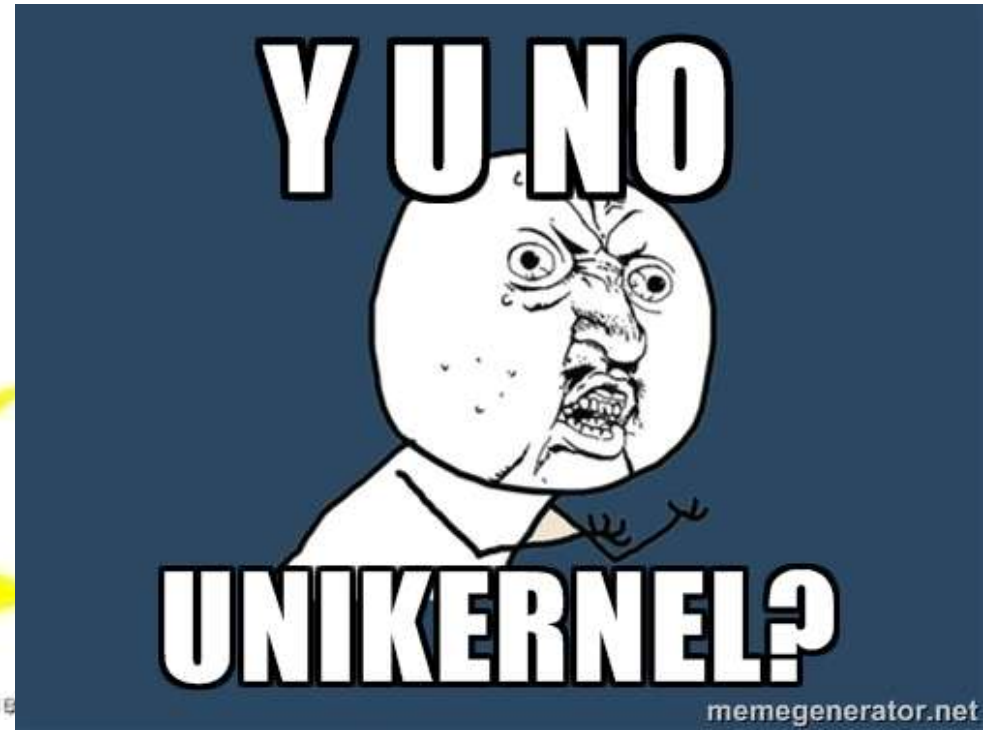
All images are taken from Google Image search and various other places on the Internet

Thanks!

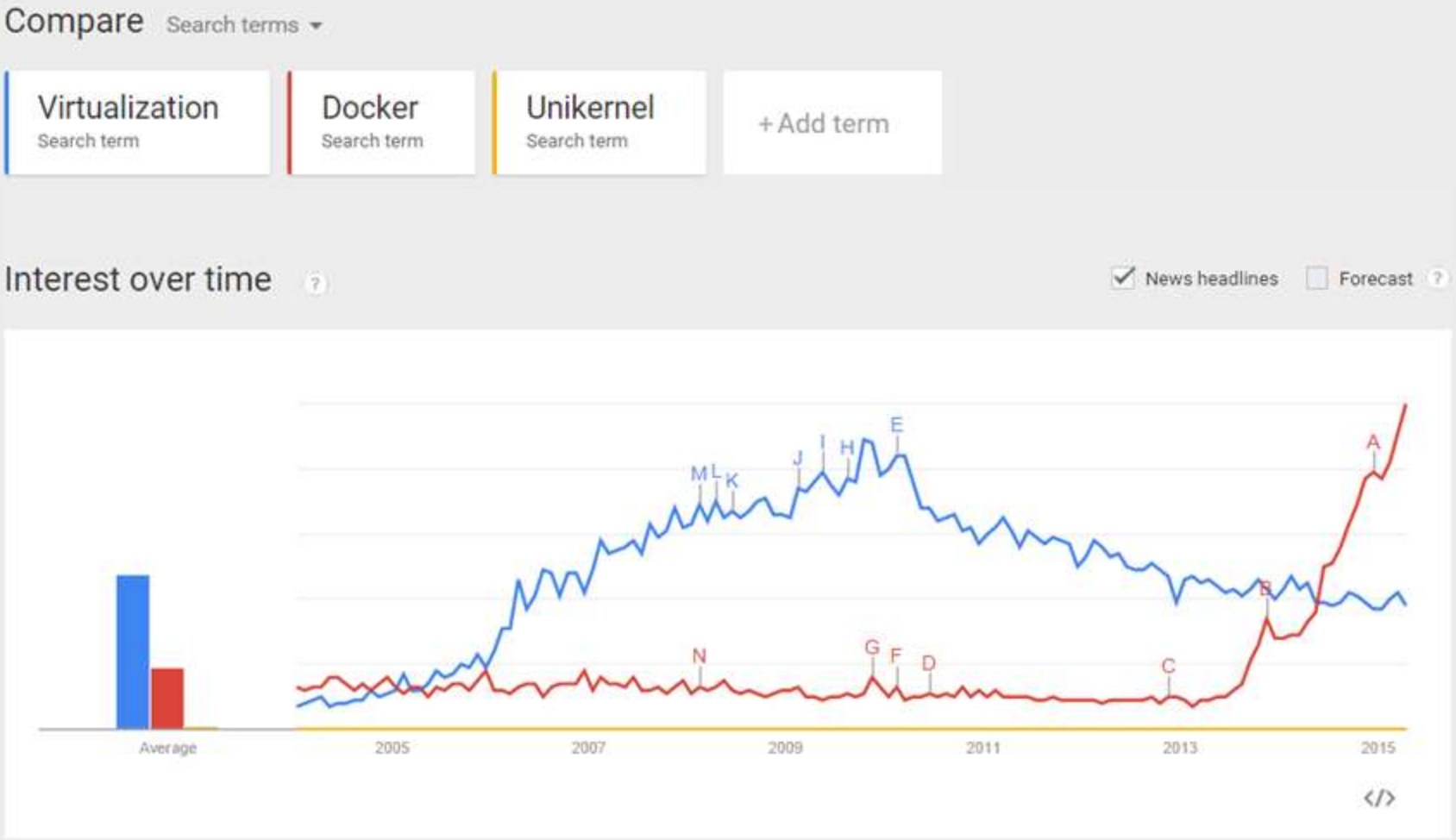
- Thanks to everybody who gave me the feedback on this presentation:
  - Dor Laor
  - Tzach Livyatan
  - Yao Zheng
  - Dave Cottlehuber

BACKUP SLIDES

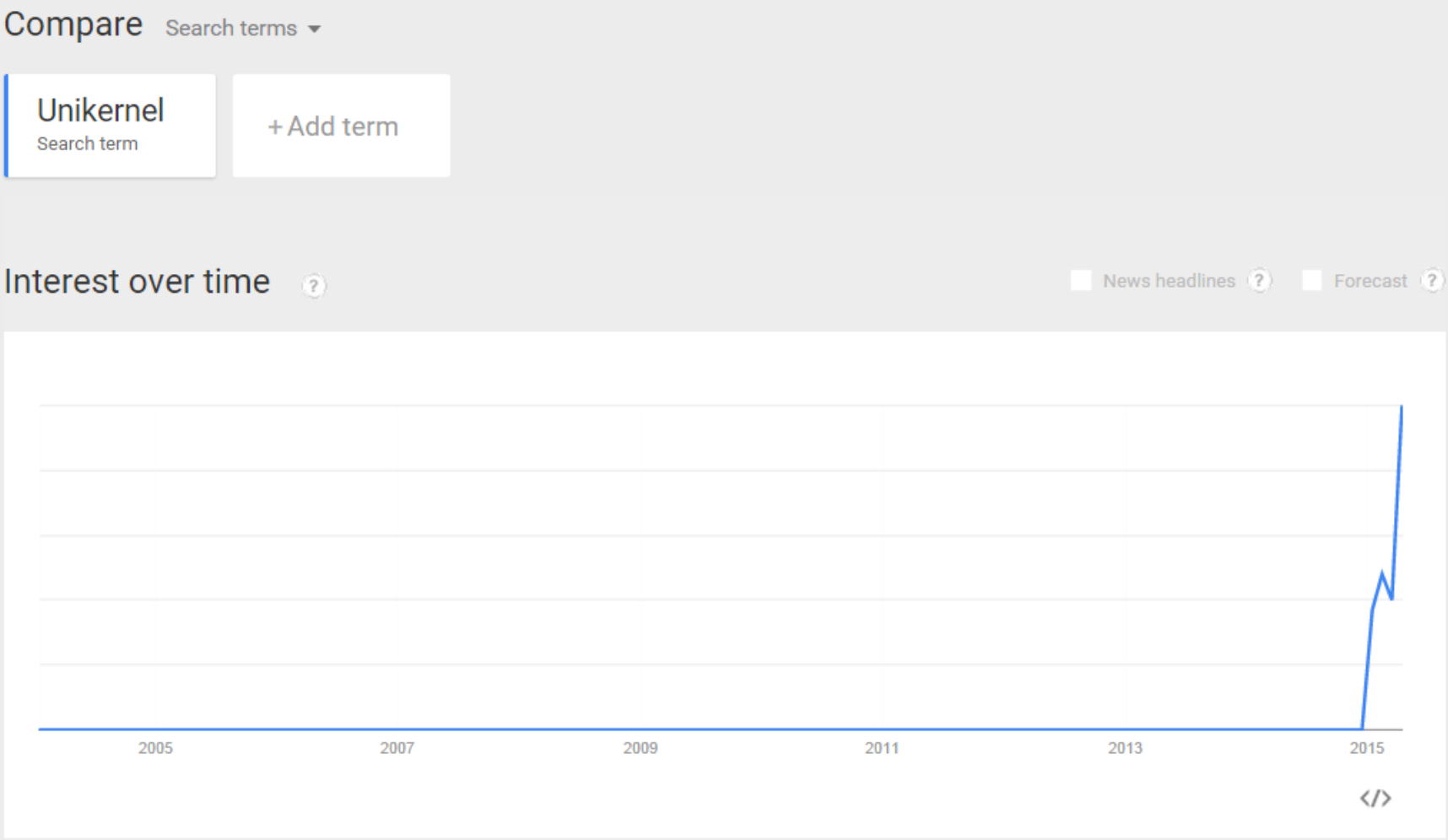
Mandatory Meme slide ;)



Old-style VMs peaked in 2010
Containers on the rise since 2013



Unikernels spiking in 2015



Ask HN: Is Docker just a step on the pathway to Unikernel deployments?

2 points by andrewstuart 326 days ago | 2 comments

Seems to me Docker is aiming for a similar thing to Unikernels - i.e. to abstract away the host operating system.

But surely Unikernels are a much more simple way to go?

Is Docker's success really just a stepping stone towards doing away with the host OS entirely via Unikernel deployments?

wmf 326 days ago

Containers and libOSes/unikernels are two opposite approaches to remove redundant "yo dawg" layers of virtualization. It's not clear to me that a hypervisor+libOS is simpler than a containerized kernel, and Linux seems better maintained than Xen.

jesusmichael 326 days ago

Containers seem to be heading more in the direction of VMs. Which would compartmentalize the services of a given application, but then you might need a separate container for a given app. I don't know if you can completely get away from the problem of shipping code to new environments without configuration/customization.