

A CutEr Tool

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Overview

- Testing
 - demo: unit, property-based, and concolic testing in Erlang
- Concolic execution for Erlang
 - demo
- Support for type specifications
 - short demo
- **CutEr**: A **C**oncolic **U**nit **T**esting tool for **E**rlang
- A “real” experience from using CutEr
 - short demo
- Concluding remarks & future work

Testing

- Testing is important
- **Unit testing** is the most widely used method

Tools: xUnit, EUnit

- Functional languages have mainly explored variants of **property-based random testing**

Tools: Haskell QuickCheck, EQC, PropEr, Triq, ...

First demo!

A first example

- A program unit:

```
classify(L) ->
  case lists:sum(L) of
    S when S < 0 -> negative;
    S when S < 4711 -> small;
    S when S > 4711 -> big;
    _ -> erlang:error(badmatch)
  end.
```

A function that classifies a list of numbers

- In general, pattern matching in Erlang provides a powerful mechanism for program assertions

```
[42,x,x|_] = f(...)
```

Testing

- In imperative languages, researchers have argued for the benefits of **concolic testing**
 - Fully automatic testing approach
 - **Concolic = Concrete + Symbolic**
 - Aims to achieve high path coverage

Tools: DART, CUTE, Symbolic Java PathFinder, jCUTE, SAGE, ...

Concolic execution

- Also known as **dynamic symbolic execution**
- **Main idea:**
 - during concrete execution, collect symbolic constraints on program inputs that cause the program to follow a specific execution path and
 - use these constraints to force execution of other paths
- **Properties/advantages:**
 - concrete execution makes available accurate information about program state which may not be easily accessible when using e.g. random testing or static analysis techniques

Implementation of concolic execution

- Symbolic execution is enabled by instrumenting the program with code that collects *path constraints* without disrupting its concrete execution
- Each variable that depends on input has both a concrete and a symbolic value associated to it
- Path constraints are expressed in an appropriate logic
- Off-the-shelf constraint solvers, often SMT ones, are used to solve these constraints and generate new inputs that will steer the future test runs to explore unexplored paths
- The execution paths can be expressed as a *symbolic execution tree*
 - each leaf node has a path constraint describing the input values that force the program to follow that specific path

A second example

```
-module(sf2).  
-export([foo/1]).  
  
foo(L) ->  
    lists:foreach(fun fcmp/1, L).  
  
fcmp(X) ->  
    case cmp(X) of  
        gt -> ok;  
        lt -> ok  
    end.  
  
cmp(X) when X > 42 -> gt;  
cmp(42) -> eq;  
cmp(X) when X < 42 -> lt.
```

Second demo!

Second example in Core Erlang

```
module example [foo/1] =
  foo/1 = fun (_cor0) ->
    call lists:foreach (fcmp/1, _cor0)

  fcmp/1 = fun (_cor0) ->
    case <apply cmp/1 (_cor0)> of
      <gt>      when true -> ok
      <lt>     when true -> ok
      <_cor1> when true -> FAIL
    end

  cmp/1 = fun (_cor0) ->
    case <_cor0> of
      <X>      when call erlang: '>' (_cor0, 42) -> gt
      <42>     when true                          -> eq
      <X>      when call erlang: '<' (_cor0, 42) -> lt
      <_cor1> when true                          -> FAIL
    end
end
```

Control flow graphs of functions

```

module example [foo/1] =
  foo/1 = fun (_cor0) ->
    call lists:foreach (fcmp/1, _cor0)

```

```

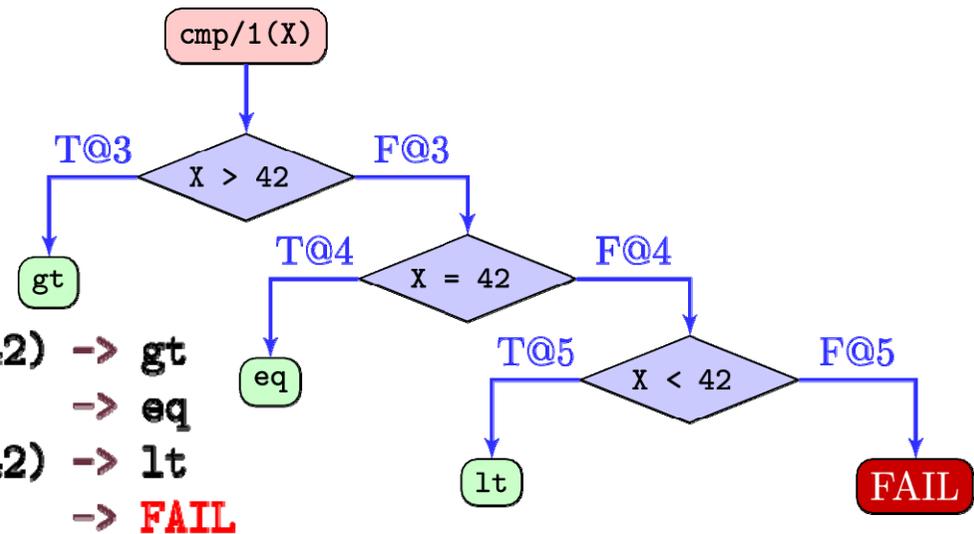
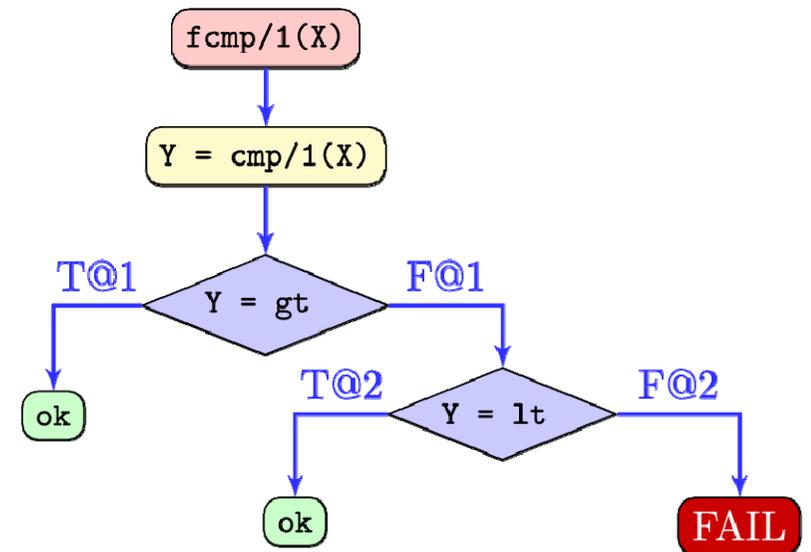
fcmp/1 = fun (_cor0) ->
  case <apply cmp/1 (_cor0)> of
    <gt>    when true -> ok
    <lt>   when true -> ok
    <_cor1> when true -> FAIL
  end

```

```

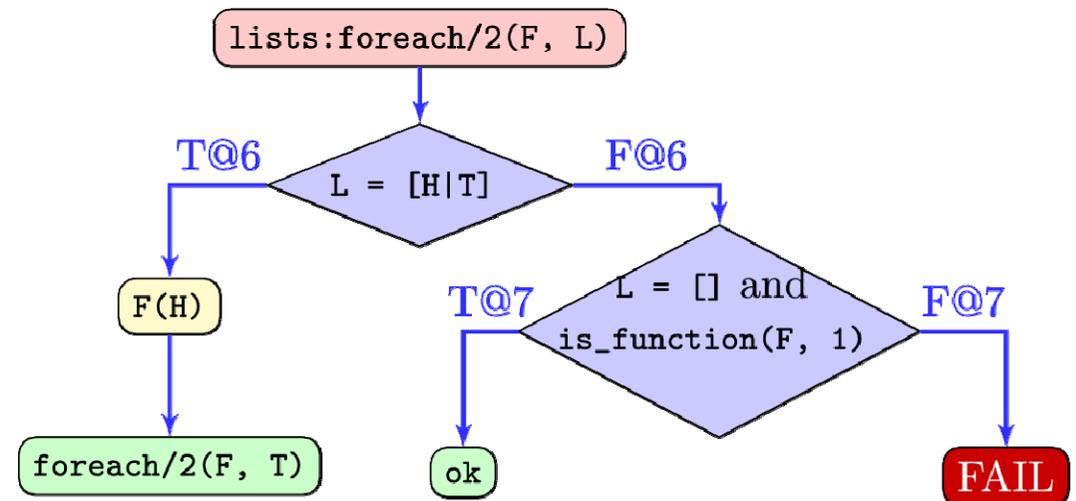
cmp/1 = fun (_cor0) ->
  case <_cor0> of
    <X>    when call erlang:'>' (_cor0, 42) -> gt
    <42>   when true -> eq
    <X>    when call erlang:'<' (_cor0, 42) -> lt
    <_cor1> when true -> FAIL
  end

```



Control flow graphs of functions

```
module lists [..., foreach/2, ...] =  
  ...  
  foreach/2 = fun (_cor1, _cor0) ->  
    case <_cor1, _cor0> of  
      <F, [H|T]> when true ->  
        do apply F (H)  
          apply foreach/2 (F, T)  
      <F, []> when call erlang:is_function (_cor0, 1) -> ok  
      <_cor3, _cor2> when true -> FAIL  
    end  
  ...
```



L ↦ [17] ; L

L ↦ [17] ; L

X ↦ 17 ; hd(L)

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X ↦ 17 ; hd(L)

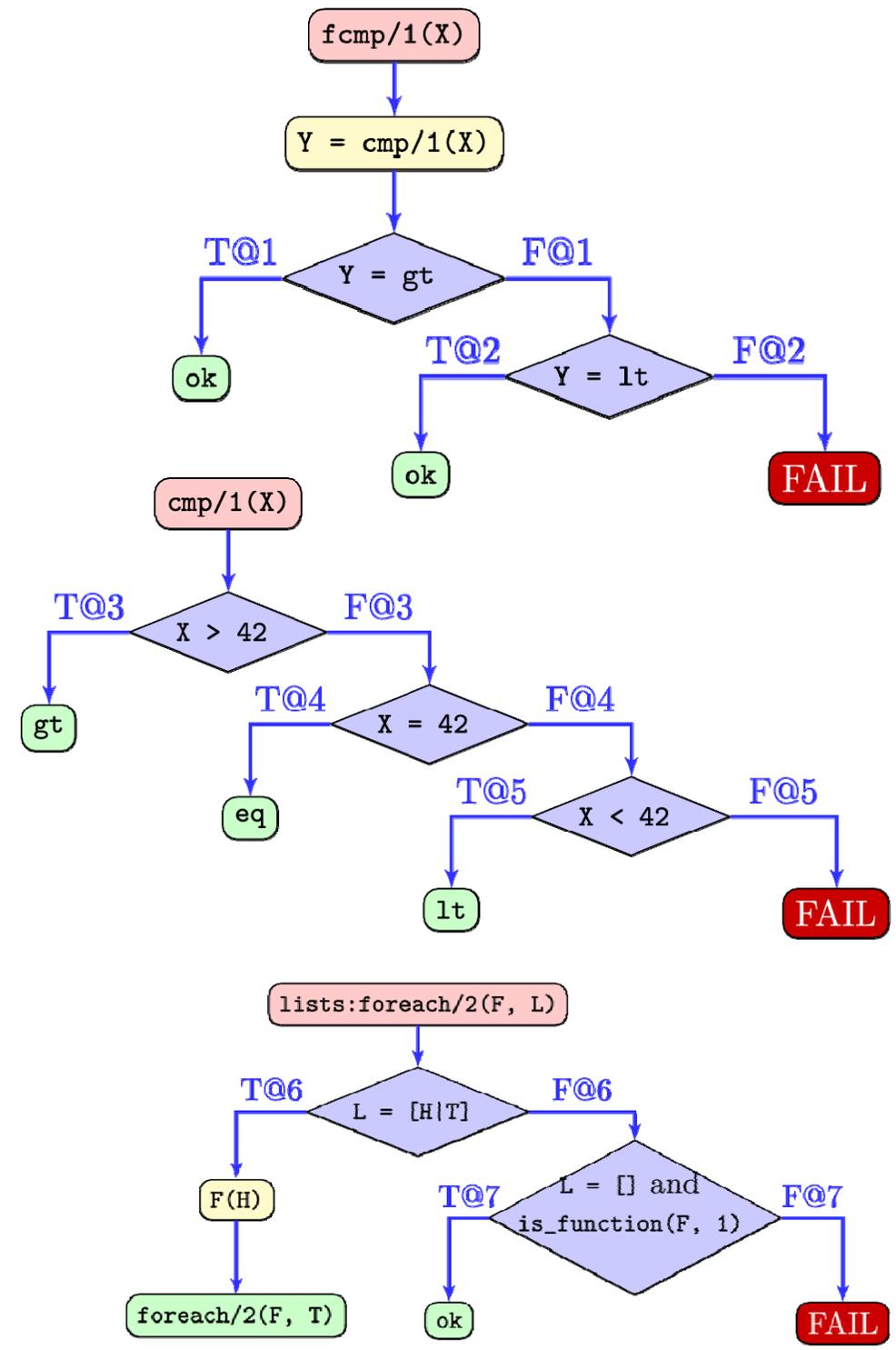
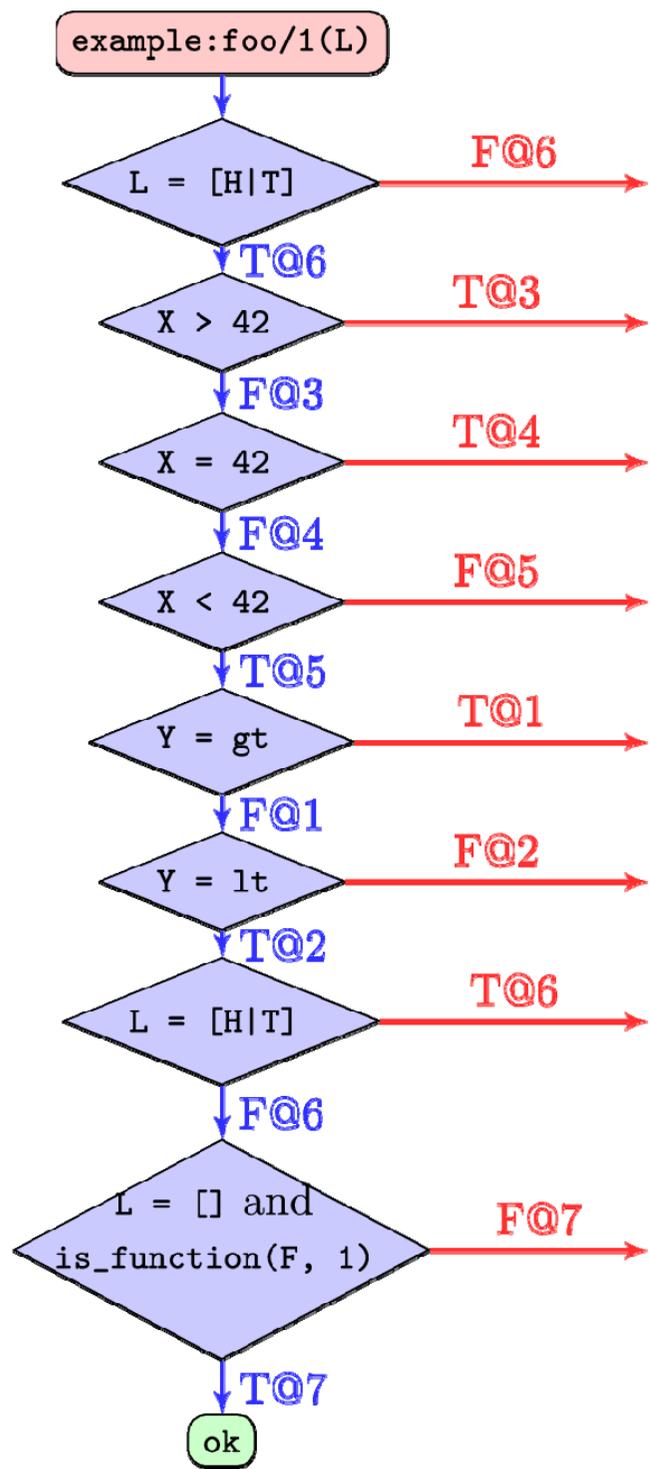
Y ↦ lt

Y ↦ lt

L ↦ [] ; t1(L)

L ↦ [] ; t1(L)

F ↦ fcmp/1



L ↦ [17]; L

example:foo/1(L)

L ↦ [17]; L

L = [H|T] → F@6 → L = [H|T] ∧ L = []



X ↦ 17; hd(L)

X > 42 → T@3 → L = [H|T] ∧ hd(L) > 42



X ↦ 17; hd(L)

X = 42 → T@4 → L = [H|T] ∧ ¬hd(L) > 42 ∧ hd(L) = 42

X ↦ 17; hd(L)

X < 42 → F@5 → X > 42 ∧ X = 42 ∧ X < 42

Y ↦ lt

Y = gt → T@1



Y ↦ lt

Y = lt → F@2

L ↦ []; t1(L)

L = [H|T] → T@6

L ↦ []; t1(L)
F ↦ fcmp/1

L = [] and is_function(F, 1) → F@7

ok

Search strategy

- Which decision node to reverse?
- We use two metrics:
 - If a decision node exists whose reversed (red) label has not yet been visited, reverse it
 - Else reverse the decision node which is closer to the root
- Stop when there are no decision nodes left to reverse

Depth-bounded search

- **Depth** counts **case** constructs that precede the decision node
- All constraints related to patterns and guards of a specific **case** construct are considered to be at the same depth
- Prune decision nodes whose depth exceeds a threshold

Support for type specifications

- Type specifications impose additional constraints on program inputs
- For the first demo program:

```
-type ret() :: 'negative' | 'small' | 'big'.
```

```
-spec classify([number()]) -> ret().
```

- For the second demo program:

```
-spec foreach(fun((T) -> term()), [T]) -> ok.
```

```
-spec foo([term()]) -> ok.
```

```
-spec foo([integer()]) -> ok.
```

-spec foo([term()]) -> ok.

$$\boxed{is_list(L)} \wedge L = [H|T] \wedge hd(L) > 42$$

$$\boxed{is_list(L)} \wedge L = [H|T] \wedge hd(L) > 42 \wedge hd(L) = 42$$

$$\boxed{is_list(L)} \wedge L = [H|T] \wedge L = []$$

$$\boxed{is_list(L)} \wedge X = hd(L) \wedge X > 42 \wedge X = 42 \wedge X < 42$$

-spec foo([integer()]) -> ok.

$$\boxed{is_integer_list(L) \implies L = [] \vee is_integer(hd(L)) \wedge is_integer_list(tl(L))}$$

$$X = hd(L) \wedge \boxed{is_integer(X)} \wedge X > 42 \wedge X = 42 \wedge X < 42$$

Third demo!

The first example with some twists

- A program unit:

```
classify(L) when length(L) < 4 -> tiny;
classify(L) ->
  case lists:foldl(fun erlang:'+' / 2, 0, L) of
    S when S < 0 -> negative;
    S when S < 4711 -> small;
    S when S > 4711 -> big
  end.
```

A function that classifies a list of numbers

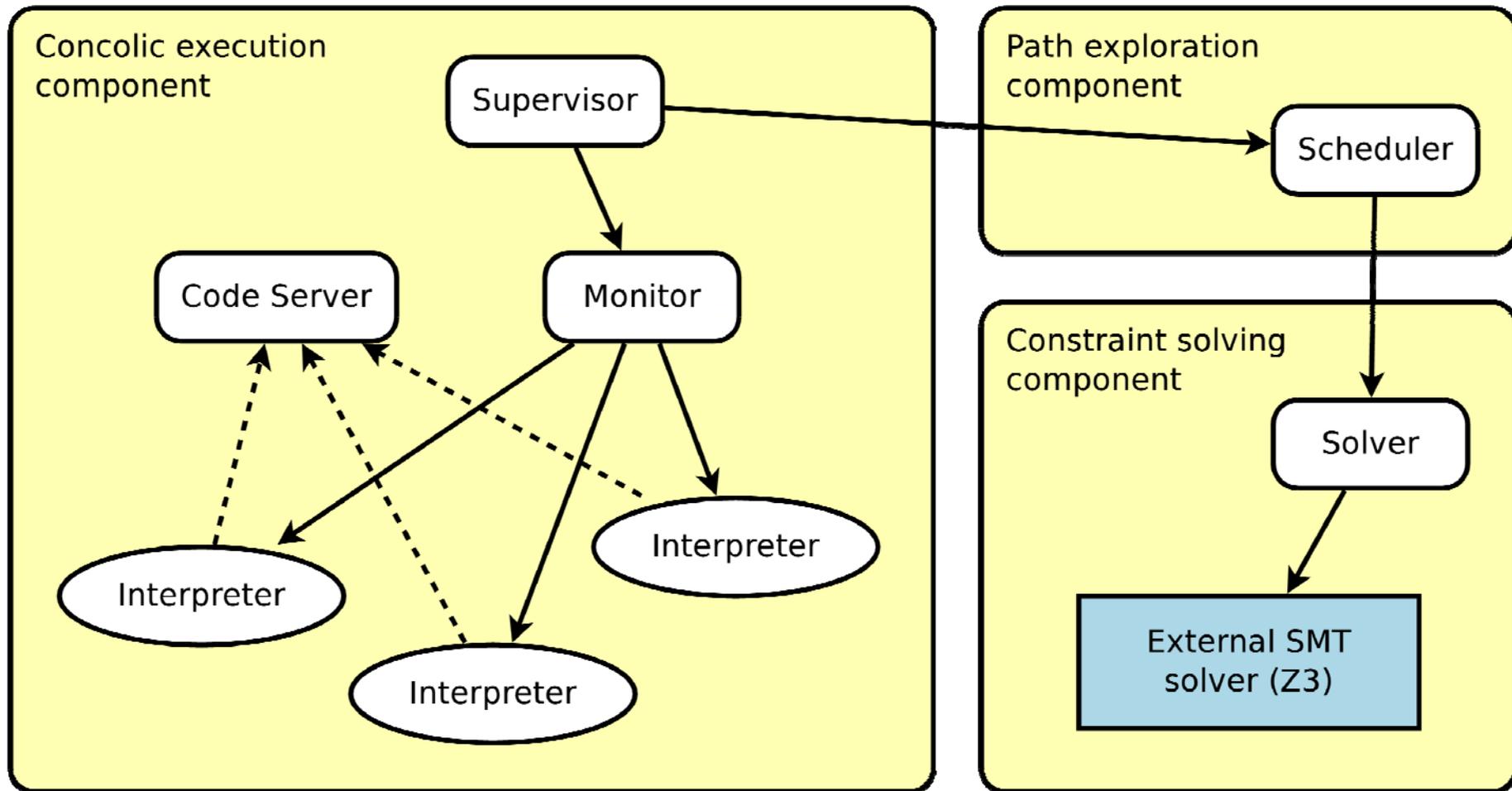
One more twist

- A program unit:

```
classify(_, L) when length(L) < 4 -> tiny;
classify(F, L) ->
  case lists:foldl(F, 0, L) of
    S when S < 0 -> negative;
    S when S < 4711 -> small;
    S when S > 4711 -> big
  end.
```

A function that classifies a list of numbers

CutEr: Concolic Unit Testing for Erlang



SMT solving with Z3

- Define the most general type, i.e. `term()`

```
Term, TList, IList = Datatypes('Term, TList, IList')

Term.declare('int', ('ival', IntSort()))
Term.declare('real', ('rval', RealSort()))
Term.declare('atm', ('aval', IList))
Term.declare('lst', ('lval', TList))
Term.declare('tpl', ('tval', TList))

TList.declare('nil')
TList.declare('cons', ('hd', Term), ('tl', TList))

IList.declare('anil')
IList.declare('acons', ('ahd', IntSort()), ('atl', IList))
```

Encoding values in Z3

- For example, for the terms:
 - 42
 - [17, 42]
 - {42, ok}

```
t1 = Term.int(42)
t2 = Term.lst(TList.cons(Term.int(17),
    TList.cons(Term.int(42), TList.nil)))
t3 = Term.tpl(TList.cons(Term.int(42),
    TList.cons(Term.atm(
        IList.acons(111, IList.acons(107, IList.anil))
    ), TList.nil)))
```

Encoding axioms in Z3

- For example, the path constraint:

$$\neg (\mathbf{L} = [\mathbf{H}|\mathbf{T}]) \wedge \neg (\mathbf{L} = [])$$

```
Not(And(Term.is_lst(L), TList.is_cons(Term.lval(L))))  
Not(And(Term.is_lst(L), TList.is_nil(Term.lval(L))))
```

CutEr

- Available on GitHub:

<https://github.com/aggelgian/cuter>

- Requires Erlang/OTP 17.x or 18.x

Current known limitations:

- Does not support maps (yet!)
- Support for recursive types is still incomplete

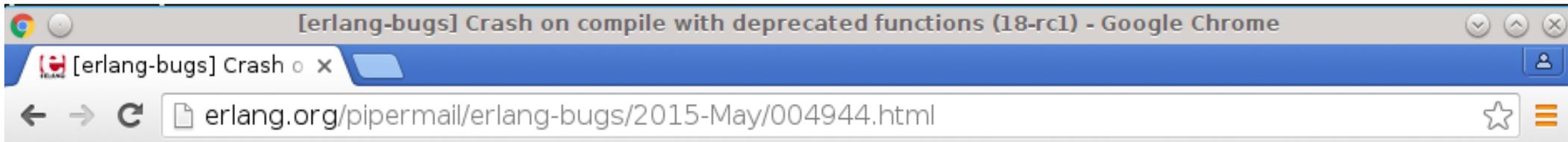
A bigger unit to test

- A post in the erlang-bugs mailing list:

```
http://erlang.org/pipermail/erlang-bugs/2015-May/004944.html
```

- Module `otp_internal` from Erlang/OTP 18.0-rc1

```
https://github.com/erlang/otp/blob/OTP-18.0-rc1/lib/stdlib/src/otp\_internal.erl
```



[erlang-bugs] Crash on compile with deprecated functions (18-rc1)

Loïc Hoguin <essen@ninenines.eu>

Sat May 2 17:11:30 CEST 2015

- Previous message: [\[erlang-bugs\] FreeBSD FPE issue on ERTS_FP_CHECK_INIT Re: ERTS_FP_CHECK_INIT error of HiPE in 18.0-rc1 running on FreeBSD 10.1-STABLE](#)
- Next message: [\[erlang-bugs\] Crash on compile with deprecated functions \(18-rc1\)](#)
- **Messages sorted by:** [\[date \]](#) [\[thread \]](#) [\[subject \]](#) [\[author \]](#)

Hello,

Some of my applications don't compile anymore because they have the `ssl:negotiated_next_protocol` instead of `negotiated_protocol`.

Problem is there is a crash instead of a nice error:

```
src/gun.erl: internal error in lint_module;  
crash reason: {case_clause,{deprecated,{ssl,negotiated_protocol}}}  
  
in function otp_internal:obsolete/3 (otp_internal.erl, line 33)  
in call from erl_lint:deprecated_function/5 (erl_lint.erl, line 3551)  
in call from erl_lint:check_remote_function/5 (erl_lint.erl, line 3527)  
in call from erl_lint:expr/3 (erl_lint.erl, line 2166)  
in call from erl_lint:expr/3 (erl_lint.erl, line 2111)  
in call from erl_lint:expr/3 (erl_lint.erl, line 2250)  
in call from erl_lint:exprs/3 (erl_lint.erl, line 2044)  
in call from erl_lint:icrt_clause/3 (erl_lint.erl, line 3029)
```

Demo!

Concluding remarks

- This presentation:
 - Concolic testing for the “functional” subset of Erlang
 - CutEr: a tool that implements this approach

<https://github.com/aggelgian/cuter>

- Future Work

- Better search strategies
- Experiment with more SMT solvers
- Handle concurrency



Thanks!