

Making the Web Functional

@chris_mccord / @czaplic

The dysfunctional Web on the server

- There is no more free lunch
- The world is moving from mostly stateless connections to increasingly stateful ones
- This is at odds with most languages' concurrency models
- OO solutions are becoming increasingly more complex to handle the demands of the modern web

The dysfunctional Web on the client

- Fragile tooling
- Competing async models
- Framework churn and framework fatigue as the community seeks ideal architectures



Programming Phoenix

Productive |> Reliable |> Fast

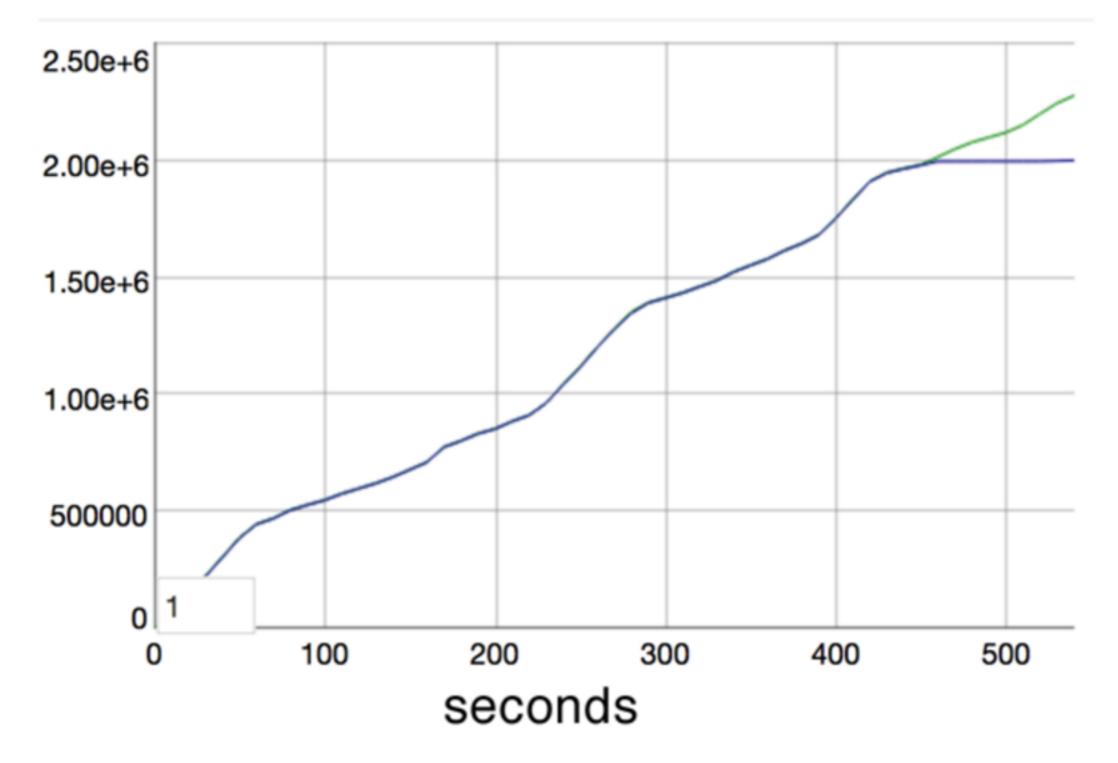




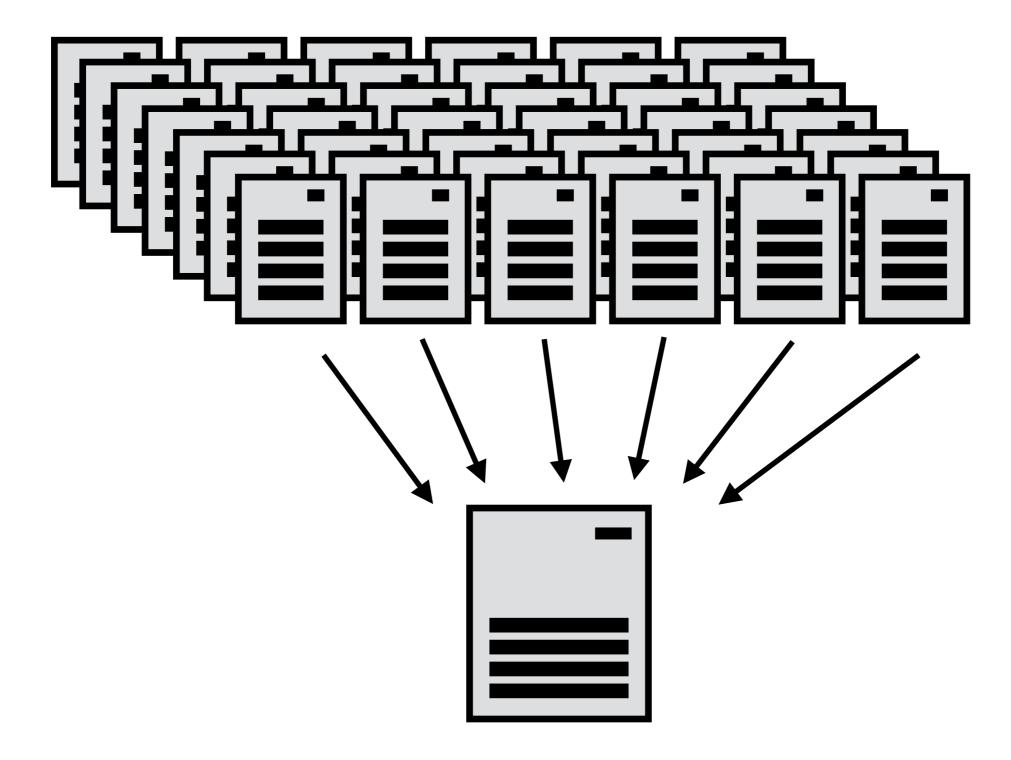
www.dockyard.com

2M channel clients on one server

Simultaneous Users



45 Rackspace 4GB 2vCPU tsung fleet



1 Rackspace 40core 2.8Ghz Xeon 128GB Phoenix server

83 of 128GB used

1 [0.0%	11 []	0.5%	21 [0.0%	31 [0.0%
2 [0.0%	12 []	0.5%	22 [0.0%	32 [0.0%
3 [0.0%	13 [0.0%	23 [0.0%	33 [0.0%
4 [1.0%	14 [0.0%	24 [0.5%	34 [0.0%
5 [0.5%	15 [0.0%	25 [0.0%	35 [0.0%
6 [0.5%	16 [0.0%	26 [0.0%	36 [0.0%
7 [0.0%	17 [0.0%	27 [0.0%	37 [0.0%
8 [1.0%	18 [0.0%	28 [0.5%	38 [0.0%
9 [0.0%	19 [0.0%	29 [0.0%	39 [0.0%
10 [0.0%	20 [0.0%	30 [0.0%	40 [0.0%
Mem[83765/12	28906MB]	Tasks:	22, 150 t	hr; 2 ru	nning
Swp[0/0MB]			Load average: 5.98 5.45 3.98			
				Uptime: 5 days, 11:17:13			

Broadcasting to 2M subscribers

Phoenix Chat \equiv	Phoenix Chat \equiv				
[emanecord] nello wond:	[cnnamecord] neiro wond:				
@ chrismccord	@ gazler				
ho					

Optimizations

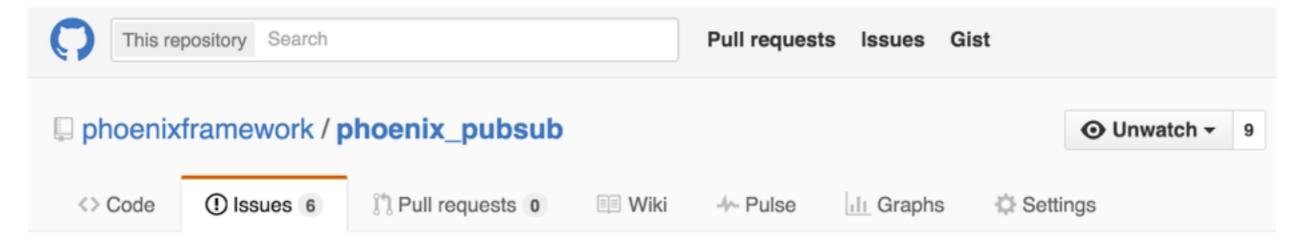
- $30k \rightarrow 60k$ subscribers
 - 14 additions and 69 deletions
- $60k \rightarrow 330k$ subscribers
 - 5 additions and 38 deletions
- 330k \rightarrow 450k subscribers (+10x arrival rate)
 - 1 addition and 1 deletion

observer

•••				nonode@nohos	st			
	System	Load Charts	Memory Allocators	Applications	Processes	Ţ	ole Viewer	Trace Overview
Pid	N	lame or Initial Func		Reds	Memory	MsgQ	Current Func	tion
<0.61.0	> tir	mer_server		1202342	24776	120	gen_server:loo	op/6
<0.105.0)> ge	en:init_it/6		0	8920	0	wx_object:loo	p/6
<0.97.0	> er	lang:apply/2		0	2752	0	observer_back	kend:flag_holder_proc/1
<0.95.0	> er	lang:apply/2		21645	142808	0	observer_pro_	wx:table_holder/1
<0.94.0	> ge	en:init_it/6		889	24808	0	wx_object:loo	p/6
<0.93.0	> ge	en:init_it/6		0	8920	0	wx_object:loo	p/6
<0.92.0	> ge	en:init_it/6		0	7048	0	wx_object:loo	p/6
<0.91.0	> ge	en:init_it/6		0	8920	0	wx_object:loo	p/6
<0.90.0	> ge	en:init_it/6		0	88728	0	wx_object:loo	p/6
<0.89.0	> er	lang:apply/2		572	8744	0	timer:sleep/1	
<0.88.0)> w	xe_master		0	21680	0	gen_server:loo	op/6
<0.87.0)> w	xe_server:init/1		72924	27664	0	gen_server:loo	op/6
<0.55.0	> Eli	ixir.Logger.Watcher:	init/1	0	2864	0	gen_server:loo	op/6
<0.54.0	> Eli	ixir.Logger.Watcher:	init/1	0	2968	0	gen_server:loo	op/6
<0.53.0	> Eli	ixir.Logger.Watcher		0	7016	0	gen_server:loo	op/6
<0.52.0	> Eli	ixir.Logger.Watcher:	init/1	0	2968	0	gen_server:loo	op/6
<0.51.0	> Eli	ixir.Logger		0	7264	0	Elixir.GenEven	t:fetch_msg/5
<0.50.0	> Eli	ixir.Logger.Supervis	or	0	10880	0	gen_server:loo	op/6
<0.49.0	> ap	plication_master:sta	art_it/4	0	6912	0	application_m	aster:loop_it/4
<0.48.0	> ap	plication_master:ini	t/4	0	2864	0	application_m	aster:main_loop/2
<0.45.0	> Eli	ixir.IEx.Config		0	2824	0	gen_server:loo	op/6
<0.44.0	> Eli	ixir.IEx.Supervisor		0	5872	0	gen_server:loo	op/6
<0.43.0	> ap	plication_master:sta	art_it/4	0	2760	0	application_m	aster:loop_it/4



Was playing around with :observer and spotted a memory leak. Took about 30 seconds to see it was monitor refs. **#myelixirstatus**



memory leak (monitor refs) in Phoenix.PubSub.Local #23

Open adamkittelson opened this issue 32 minutes ago · 1 comment



adamkittelson commented 32 minutes ago

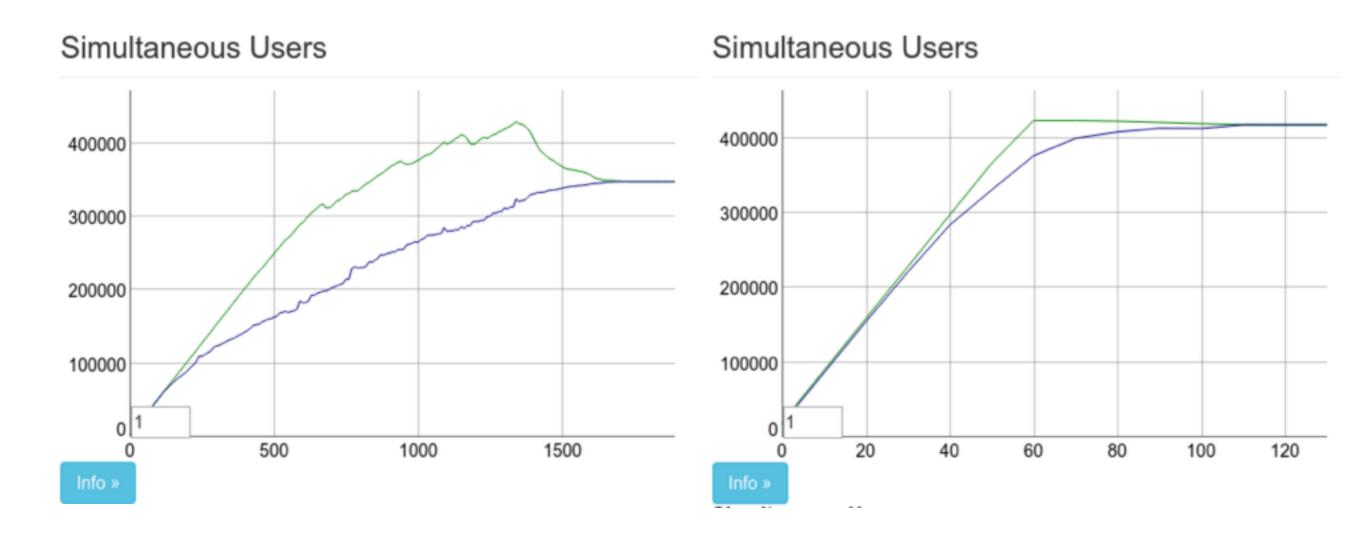
I think I found a memory leak in Phoenix.PubSub.Local with monitor refs and long running subscribers that join topics frequently.

Each time a process subscribes to a topic Process.monitor is called at https://github.com/phoenixframework/phoenix_pubsub/blob/9e0e079bcb1a94 6a5c/lib/phoenix_pubsub/local.ex#L194 causing the Local GenServer to create subscriber pid regardless of whether it's already monitoring that pid as a result from another (or the same) topic.

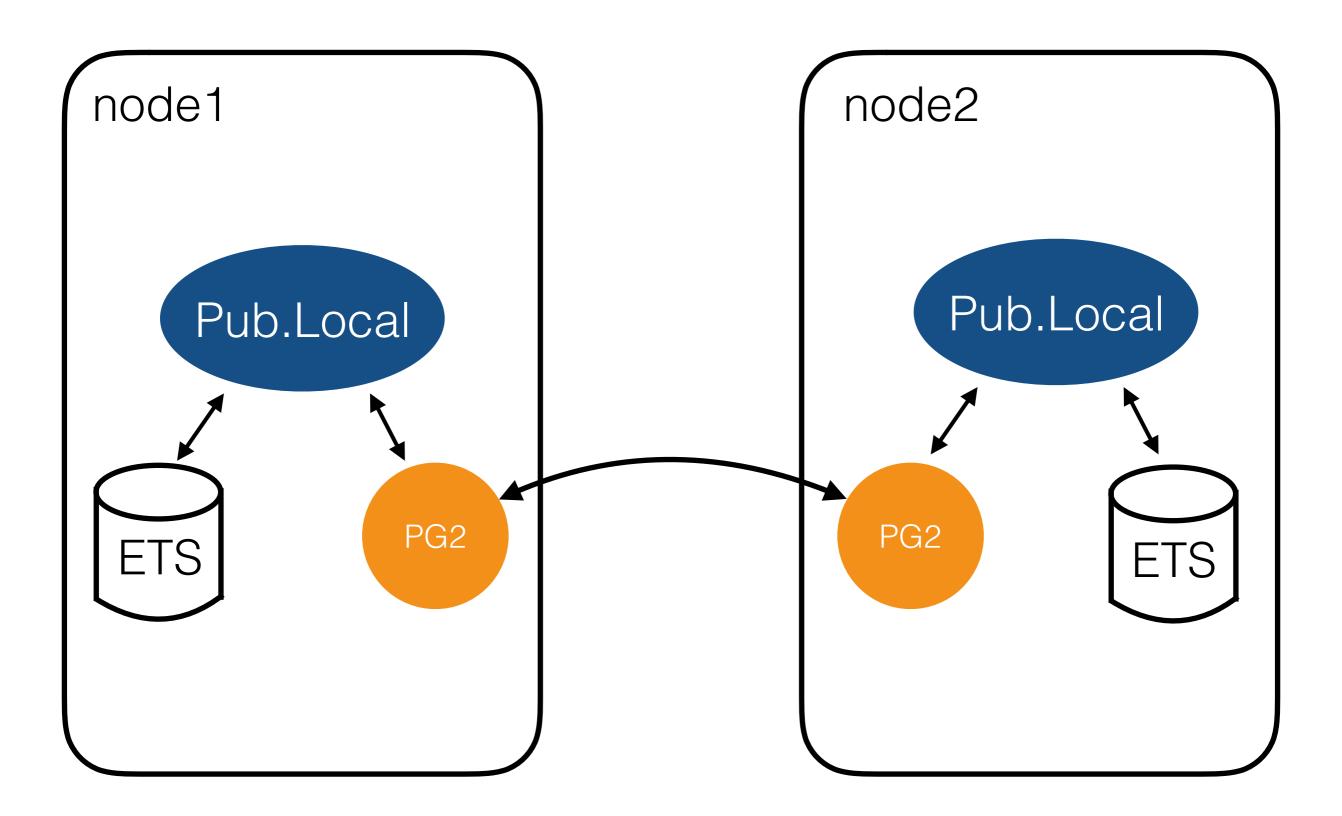


10x increase in arrival rate

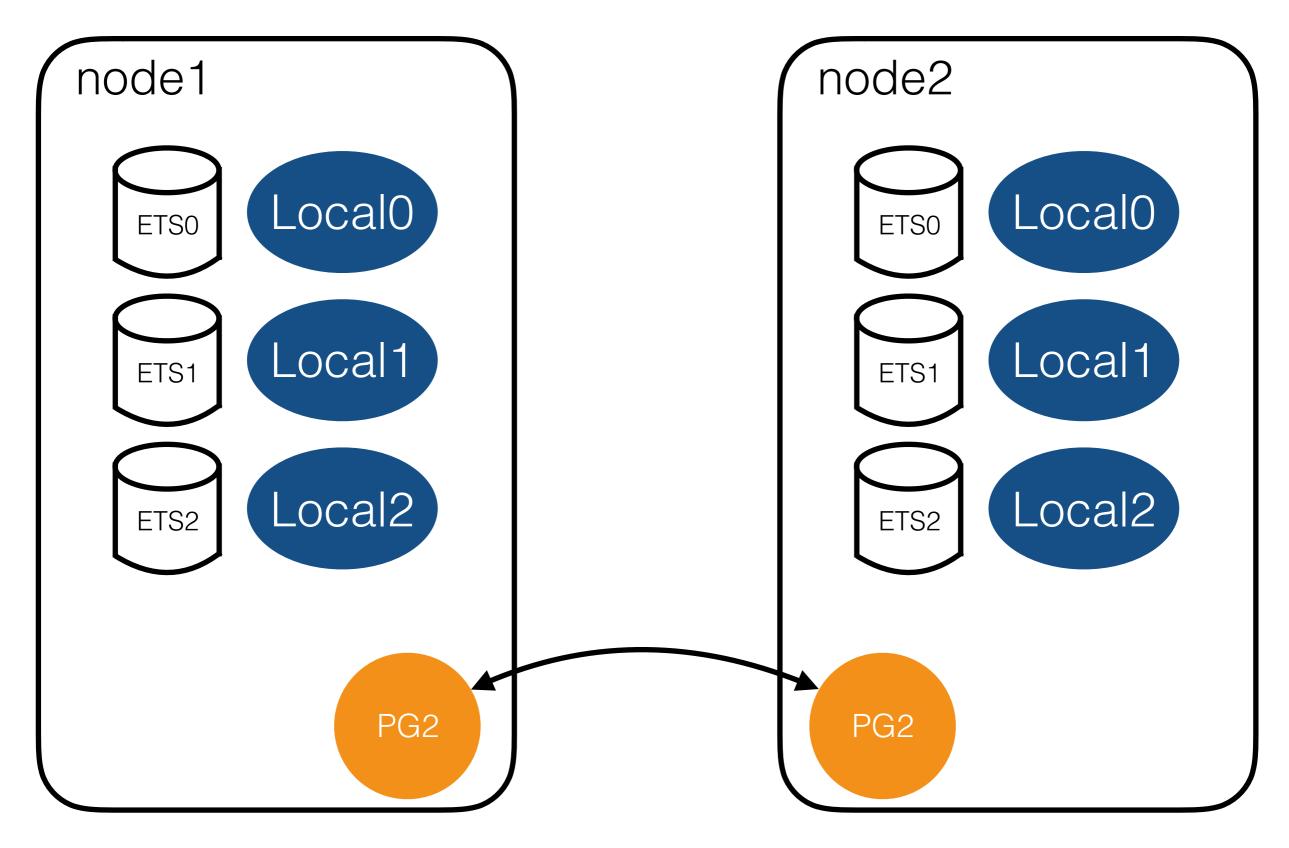
- ^local = :ets.new(local, [:bag, :named_table, :public, + ^local = :ets.new(local, [:duplicate_bag, :named_table, :public,



Sharding Subscriptions

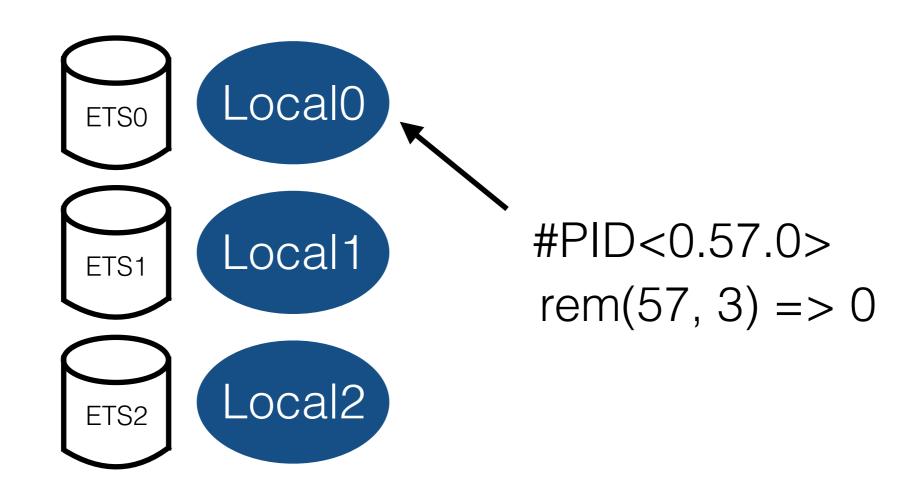


Sharding Subscriptions



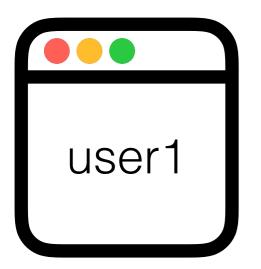
defp pid_to_shard(pid, shard_size) do pid |> pid_id() |> rem(shard_size) end defp pid_id(pid) do bin = :erlang.term_to_binary(pid) pre = (byte_size(binary) - 9) * 8 <<_::size(pre), id::size(32), _::size(40)>> = bin id

end



Phoenix.Presence

A look at the problem



Online Users

- user1
- user2



Solution?: e server that monitors and broadcas

Presence server that monitors and broadcasts joins/ leaves

```
def join("lobby", _, socket) do
    Presence.add(self(), socket.assigns.user)
    {:ok, %{users: Presence.list(), socket}
    end
```

```
defmodule Presence do
  def handle call(:list, , users) do
    {:reply, users, users)
  end
  def handle call({:add, pid, user}, , users) do
    ref = Process.monitor(pid)
   broadcast("lobby", "join", user)
    {:reply, :ok, Map.put(users, ref, user)}
  end
  def handle info({:DOWN, ref, :process, , }, users) do
    broadcast("lobby", "leave", Map.get(users, ref))
    {:noreply, Map.delete(users, ref)}
  end
end
```

A look at the problem

Online Users

• user2

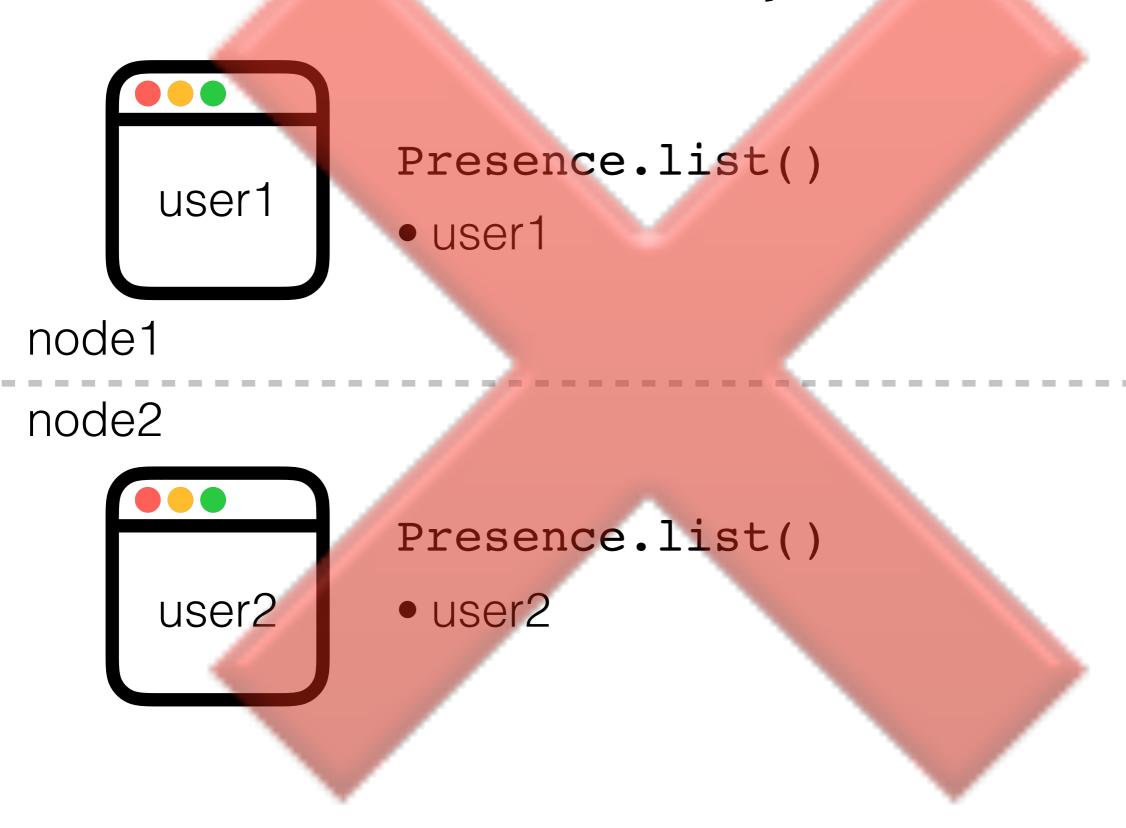


"join", user1 "join", user1 "join", user1 "leave", user1



"join", user2

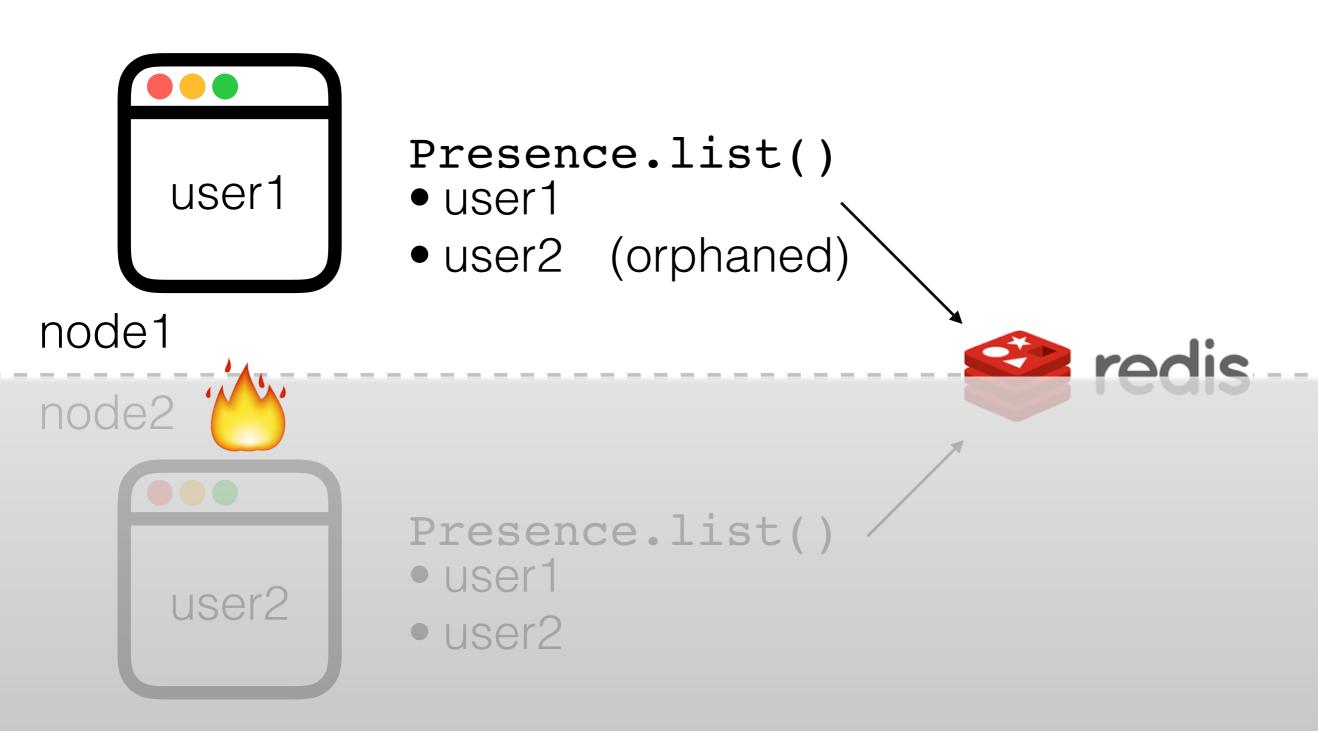
A look at the problem



Solution?: Presence backed by shared database!



A look at the problem



The Problem

- Local-node concerns
 - must account for unique presences for same user
- Multi-node concerns
 - must handle node-down events and clean up local state for presences belonging to downed node
 - must replicate data across cluster

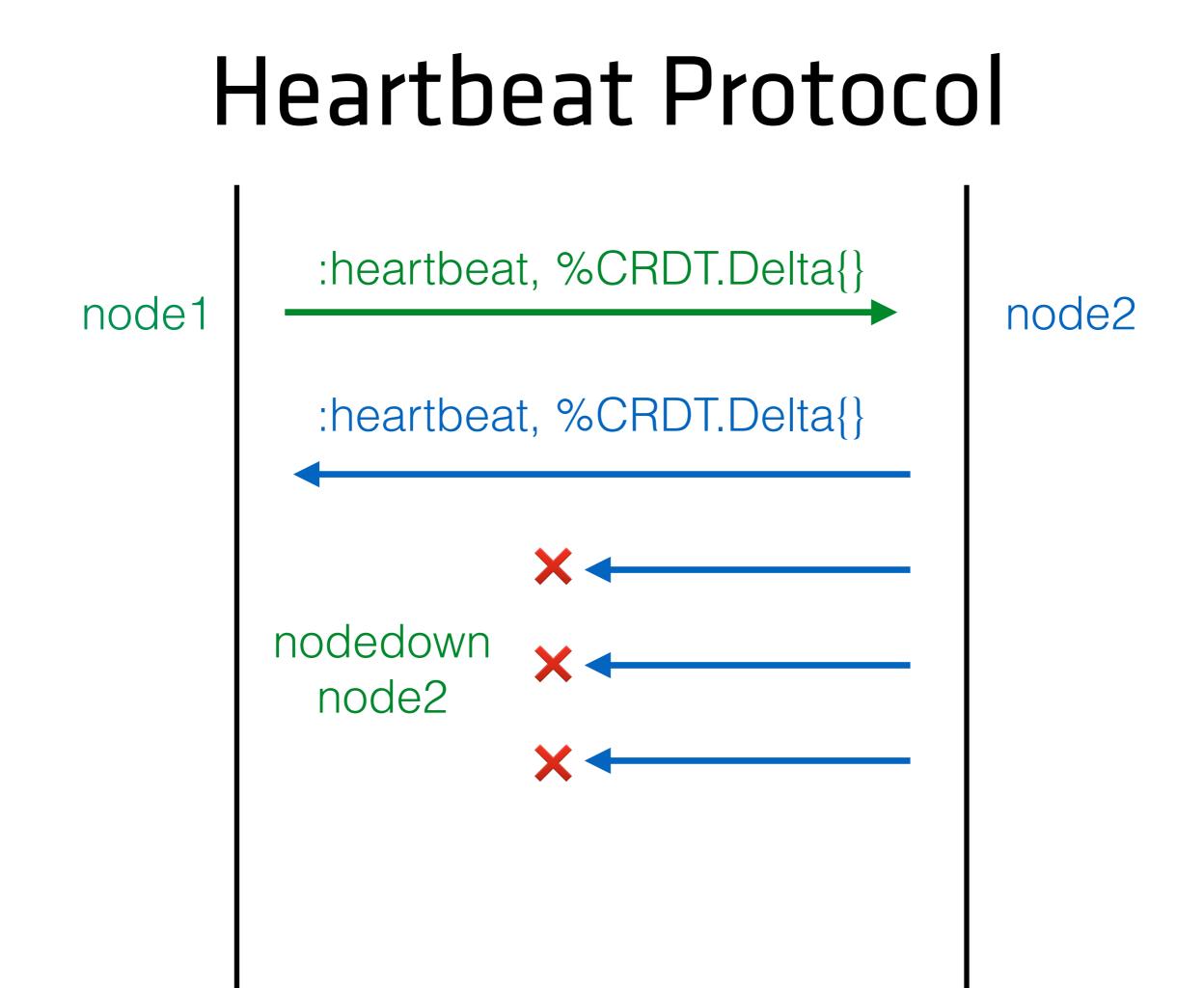
The Solution

- Ideal Architecture
 - no single source of truth
 - no single point of failure
- CRDT (Conflict-free Replicated Data Type)
- Heartbeat/Gossip protocol

CRDT

conflict-free, replicated data type

- Strong eventual consistency
- Replicate presence join and leave events across the cluster without merge conflicts
- Conflicts are *mathematically impossible*
- Supports replication without remote synchronization



Vector Clocks Catching up on missed deltas

node1

n2: $1 \rightarrow 1$, %{n3: 1, n4: 2} n3: $1 \rightarrow 2$, %{n2: 2, n4: 2} n4: $2 \rightarrow 3$, %{n2: 2, n3: 2}

has updates has updates

only n4 is selected to request updates since n3 is contained in its future

Server API

```
defmodule RoomChannel do
  def join("rooms:" <> room_id, _, socket) do
    send self(), :after_join
    {:ok, socket}
    end
```

```
def handle_info(:after_join, sock) do
    id = sock.assigns.user_id
    Presence.track(sock, id,%{status: "avail"})
    push sock, "presences", Presence.list(sock)
    {:noreply, sock}
    end
end
```

Client API

```
import {Socket, Presence} from "phoenix"
let socket = new Socket("/socket")
let room = socket.channel("rooms:" + id)
let presences = {}
room.on("presences", state => {
 Presence.syncState(presences, state)
})
room.on("presence diff", diff => {
  Presence.syncDiff(presences, diff)
})
console.log("users", Presence.list(presences)})
```

Demo

Making The Web Functional

- Good platforms drive you toward optimal solutions
- You can trust that following the principles laid out produces fast, maintainable programs
- Fast code does not have to equal dense code
- Productive code does not have to equal slow code
- Good platforms let you to focus on what matters your application

"writing great code should be easy... now it is" – elm-lang.org