

# **PULSE Tutorial**

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Q

• Very often we write code like:

```
Items = gather_items(),
lists:foreach(fun(I) -> process_item(I) end, Items)
```

or

```
Items = gather_items(),
Res = lists:map(fun(I) -> process_item(I) end, Items)
```

- Each process\_item(I) is independent
- Natural place to parallelize!



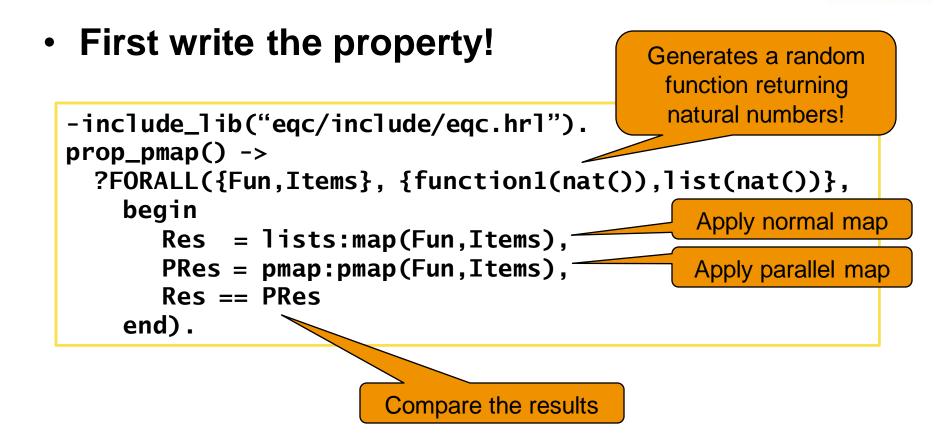
- Q
- It is enough to replace map with parallel map (pmap):

```
Items = gather_items(),
Res = pmap(fun(I) -> process_item(I) end, Items)
```

- Unfortunately there is no standard parallel map in Erlang
- How about implementing one!?



#### **Property Driven Development**





Implementing pmap

• First attempt

```
-module(pmap).
```

```
-export([pmap/2]).
```

```
pmap(F,Ls) ->
Self = self(),
[spawn(fun() -> Self ! F(L) end) || L <- Ls],
[receive Res -> Res end || _ <- Ls].</pre>
```





2>eqc:quickcheck(pmap\_eqc:prop\_pmap()).
OK, passed 100 tests
true

Good, but let's run some more tests...

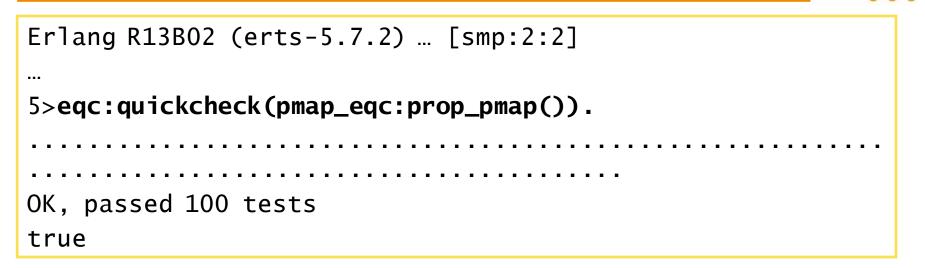
```
3>eqc:quickcheck(eqc:numtests(10000,pmap_eqc:prop_pmap())).
...
OK, passed 10000 tests
true
```



- Perfect! Move on to next problem...
- Or wait a second, what was it we tested?!?
- A **concurrent** implementation on a slow singlecore laptop!
- Not good enough!

# When a test passes, always think about what you just tested!





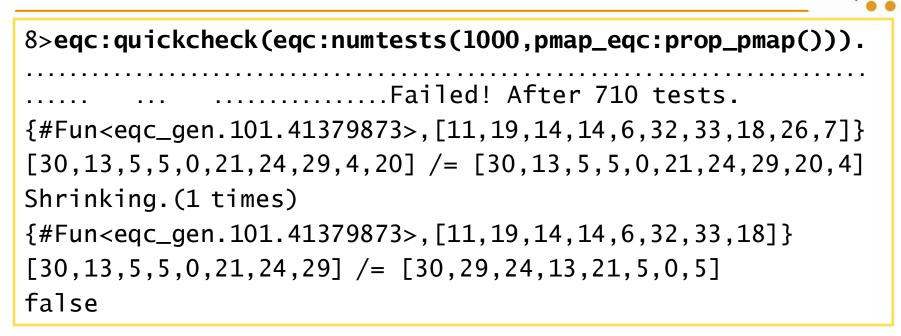
#### Still passes, maybe it is actually correct...

property based testing

- We need more information!
   ?WHENFAIL to run code when a property fail
- We want to see the values of Res and PRes.

```
-include_lib("eqc/include/eqc.hrl").
prop_pmap() ->
    ?FORALL({Fun,Items}, {function1(nat()),nat()},
    begin
        Res = lists:map(Fun,Items),
        PRes = pmap:pmap(Fun,Items),
            ?WHENFAIL(
                io:format("~p /= ~p\n",[Res,PRes]),
                  Res == PRes)
    end).
```









- Switching to multi-core (or enabling SMP) makes concurrency bugs more likely to manifest
- We had to run quite a few tests
- Shrinking didn't work (very well)
  - A small counterexample is often very valuable
  - Shrinking a counterexample is done stepwise
  - Counterexample that 'happens' to fail will not shrink well



- The Erlang scheduler is too deterministic
  - Small tests
  - Low load on system
  - Deterministic even in multi-core systems
  - Large tests are needed to provoke race conditions
  - Many race conditions may not show up until you deploy your system
- With randomized scheduling
  - Small tests are more likely to provoke race conditions
  - Find concurrency bugs early in development process



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- PULSE to the rescue
  - P ProTest
  - **U –** User
  - **L** Level
  - S Scheduler
  - E for Erlang
- PULSE is non-deterministic (random scheduling)
- PULSE can re-run a schedule (repeatable tests)



- Controls the concurrency

   Only one process is executing at a time
- Records all concurrency events
  - Message sending
  - Process spawning
  - Etc...
- **PULSE** can switch to executing another process (simulating context switch) at any time
- We make sure that unlikely scenarios get tested



• pulse\_instrument:

- Instrumentation of the code at compile time

- Implemented as parse\_transform compiler option
- Example:

c(example,[{parse\_transform,pulse\_instrument}]).

Calls to *spawn*, *link* as well as statements *!* and *receive*, etc are replaced by calls handled by PULSE



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• Running instrumented code:

Application **PULSE** must be running: **pulse:start()**.

The **PULSE** application keeps state: last used schedule, random seed, etc, and gives access to event handlers for different kind of output.



8> pulse:start(). Starting eqc version 1.18 ... 9> pulse:run(fun() ->  $pmap:pmap(fun(X) \rightarrow X + 2 end, [1, 2]) end).$ [3,4] scheduling started root spawns pmap <0.234.0> root spawns pmap1 <0.235.0> root blocks pmap sends 3 to root pmap terminated normally root receives 3 return value [3,4] scheduling finished 10>



Q

QuickCheck uses ?PULSE macro:

# ?PULSE( <Pattern bound to result of E>, <Expression E to run in PULSE>, <Property using result of E>

• Normal compilation:

Run code normally

 Compilation with pulse\_instrument, PULSE running: Run code with PULSE scheduler

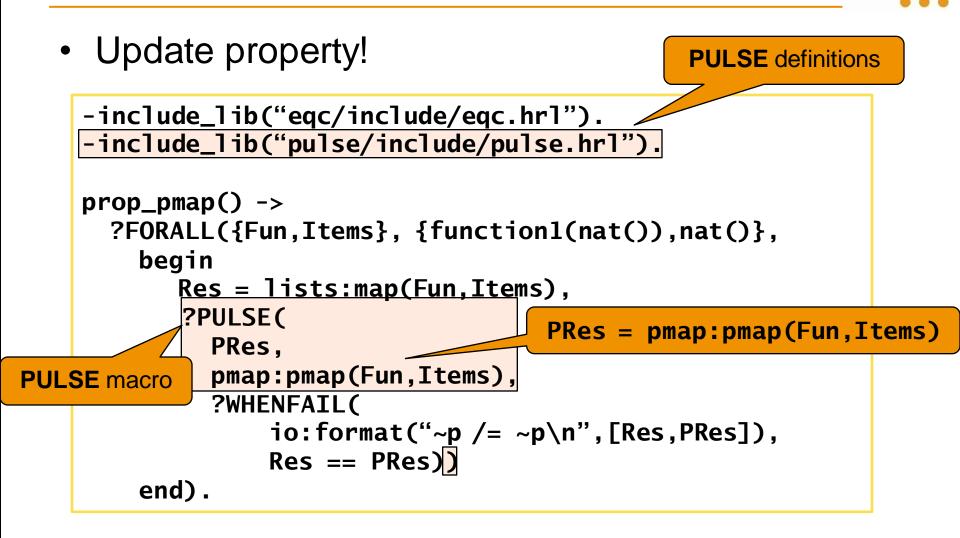


How to use PULSE with QuickCheck

Update property!



#### How to use PULSE with QuickCheck





#### Verbosity in **PULSE**

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- Don't forget the verbosity:
  - pulse:verbose/1.

```
24> pulse:verbose([]).
ok
25> pulse:run(fun() ->
       pmap:pmap(fun(X) \rightarrow X + 2 end, [1, 2]) end).
[3, 4]
26> pulse:verbose([all]).
ok
27> pulse:run(fun() ->
       pmap:pmap(fun(X) \rightarrow X + 2 end, [1, 2]) end).
[3,4]
scheduling started
root spawns pmap <0.234.0>
root spawns pmap1 <0.235.0>
. . .
```



- Verbosity options:
  - **all** All verbosity flags
  - **send** Show sending of messages
  - **'receive'** Show delivery and receiving of messages
  - **procs** Show process events (spawn, link, etc.)
  - **side\_effect** Show (user defined) side effects
- Options are similar to *trace patterns*



```
32> pulse:verbose([]).
ok
33> eqc:quickcheck(pmap_eqc:prop_pmap()).
.....Failed! After 23 tests.
{#Fun<eqc_gen.101.34457915>,[0,2,0]}
{29191,1432,12821}
[3,1,1] /= [1,3,1]
Shrinking...(3 times)
{#Fun<eqc_gen.101.34457915>,[0,1]}
{29191,1432,12821}
[3,0] /= [0,3]
false
```

- Fewer test cases needed
- Shrinking works (for this example)



- What is the error?
- We can use **pulse:rerun\_counterexample/2** to re-run the counterexample with more verbosity
  - Gets the last counterexample from eqc:counterexample/0
  - Uses eqc: check/2 to re-run the property



```
35> pulse:rerun_counterexample([all],pmap_eqc:prop_pmap()).
scheduling started
root spawns pmap <0.244.0>
root spawns pmap1 <0.245.0>
root blocks
pmap sends 3 to root
pmap terminated normal
pmap1 sends 0 to root
pmap1 terminated normal
pmap1 delivers 0 to root
root receives 0
root blocks
pmap delivers 3 to root
root receives 3
return value [0,3]
scheduling finished
Failed!
{#Fun<eqc_gen.101.34457915>, [0,1]}
\{29197, 1532, 821\}
[3,0] /= [0,3]
false
36>
```

- Another way of understanding an error
- We can visualize the schedule to easier understand it!
- Requires pulse\_event\_graph to be added as event handler: pulse\_event\_graph:start().

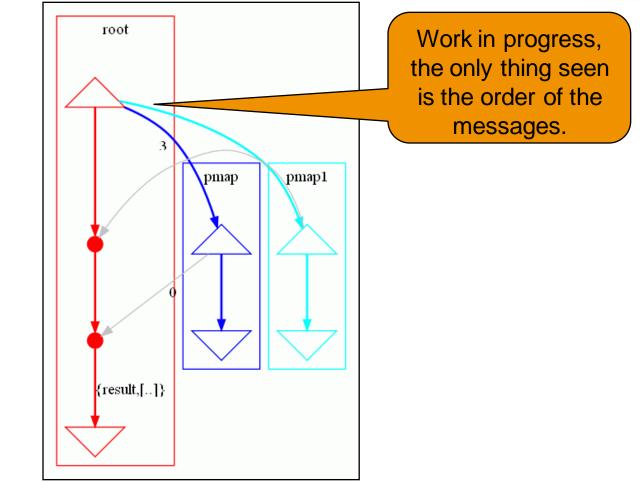
```
36> pulse_event_graph:start([]).
ok
37> pulse:rerun_counterexample([],pmap_eqc:prop_pmap()).
pulse_event_graph set verbose to []
pulse_event_terminal set verbose to []
Failed!
...
```

• Every scheduled run now creates a graph.dot file!



#### Visualization



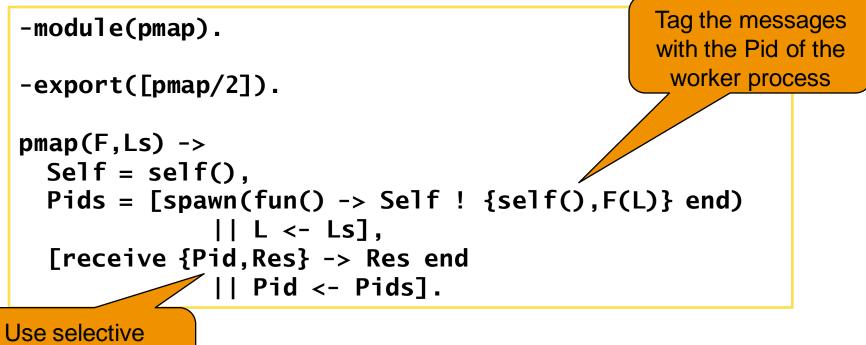


Requires GraphViz to be installed. In particular the program **dot** *http://www.graphviz.org/* 



pmap 2<sup>nd</sup> attempt

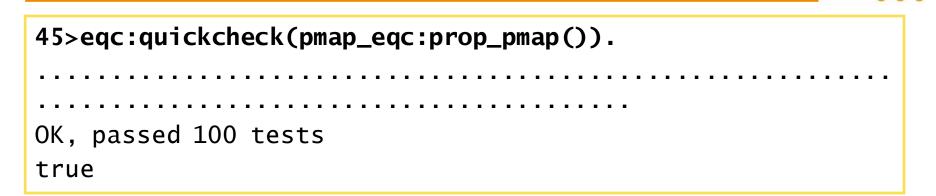
- Q
- We need to ensure the order of the results:



receive to fetch the results in order



#### Testing the new implementation



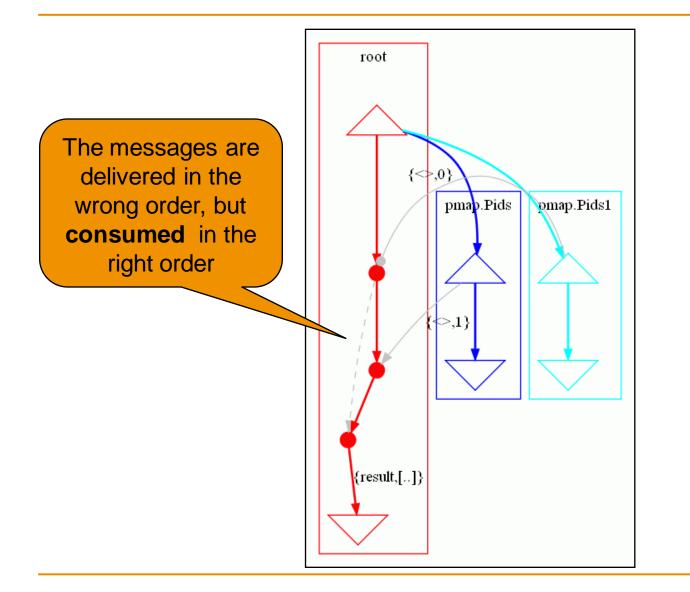
Good, but again, let's run some more tests...

48>eqc:quickcheck(eqc:numtests(10000,pmap\_eqc:prop\_pmap())).
...
OK, passed 10000 tests
true

Done! But now we should add some fault tolerance, etc...



#### Visualization – A correct run







## Short break!

## Try it yourselves!

### Next: User defined side effects



- Concurrency errors can be caused by modules interacting with other modules.
- Example: writefile

```
prop_writefile() ->
    ?FORALL({Text1,Text2}, {string(), string()},
    begin
        ok = file:write_file(?TESTFILE,Text1),
        ok = file:write_file(?TESTFILE,Text2),
        {ok,Bin} = file:read_file(?TESTFILE),
        binary_to_list(Bin) == Text2
    end).
```

Not very interesting, since it is sequential it works. How about parallel file writing?

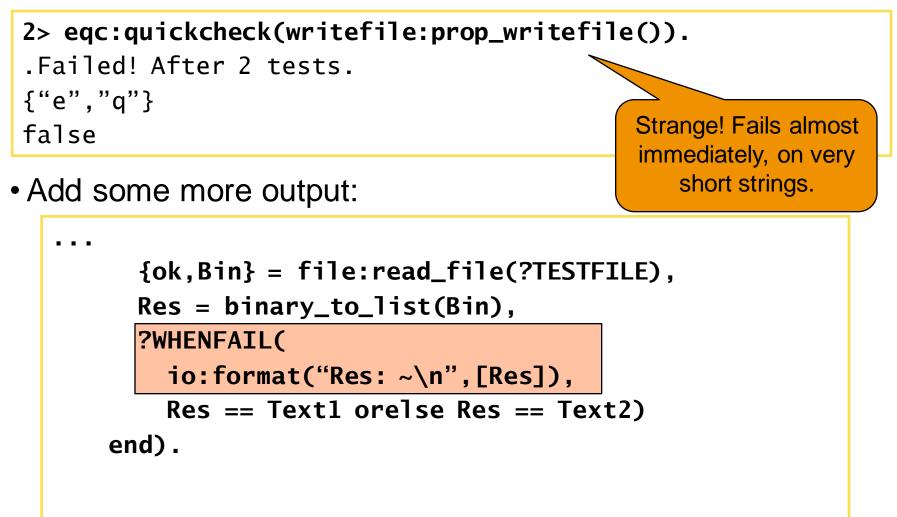




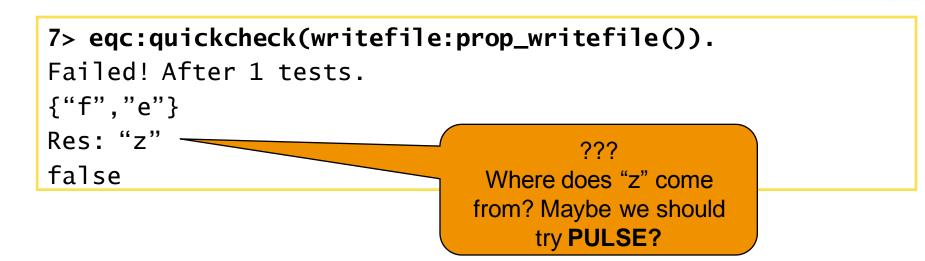
```
-define(PAR(E1,E2),
        begin
          spawn(fun() -> E1 end),
          spawn(fun() -> E2 end)
        end).
prop_writefile() ->
                                                    Write files in
  ?FORALL({Text1,Text2}, {string(), string()},
                                                      parallel
    begin
      ?PAR(file:write_file(?TESTFILE,Text1),
            file:write_file(?TESTFILE,Text2)),
      {ok,Bin} = file:read_file(?TESTFILE),
                                                  The result should
      Res = binary_to_list(Bin),
                                                   be either of the
      Res == Text1 orelse Res == Text2
                                                      strings
    end).
```

property based testing











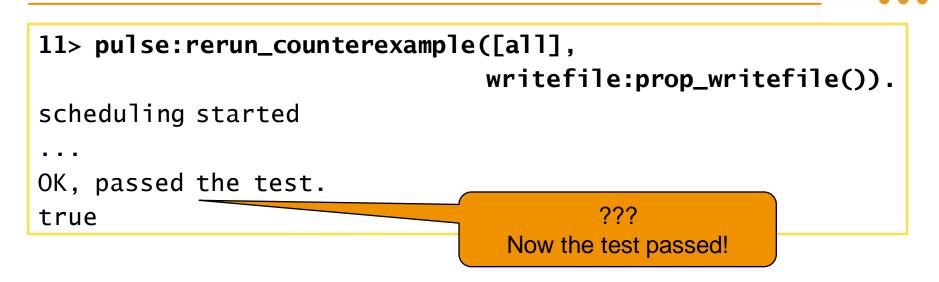
• Add ?PULSE to the property:

```
prop_writefile() ->
  ?FORALL({Text1,Text2}, {string(), string()},
    ?PULSE(
      Res.
      begin
        ?PAR(...),
        {ok,Bin} = file:read_file(?TESTFILE),
        binary_to_list(Bin),
      end,
      ?WHENFAIL(io:format("Res: ~\n",[Res]),
                Res == Text1 orelse Res == Text2)).
```



```
9> pulse:start(),pulse:verbose([all]).
10> eqc:quickcheck(writefile:prop_writefile()).
scheduling started
root spawns 'prop_writefile.Res' <0.1528.0>
root spawns 'prop_writefile.Res1' <0.1529.0>
return value "k"
'prop_writefile.Res1' terminated normal
'prop_writefile.Res' terminated normal
                                               Doesn't tell us very
scheduling finished
                                                much more, we
Failed! After 1 tests.
                                               know that write file
{"f", "e"}
                                               is a side-effect, but
                                               PULSE does not...
{8534,66433,27482}
Res: "k"
false
```





- Important PULSE fact: PULSE does not control the universe!
- **PULSE** cannot re-run a schedule (faithfully) when the environment has changed (new files are written etc...)





• We want **PULSE** to show an event when we perform a file operation.

All calls to module file are considered side effects: c(writefile,

[{parse\_transform,pulse\_instrument},

{pulse\_side\_effect,[{file,'\_','\_'}]}]).

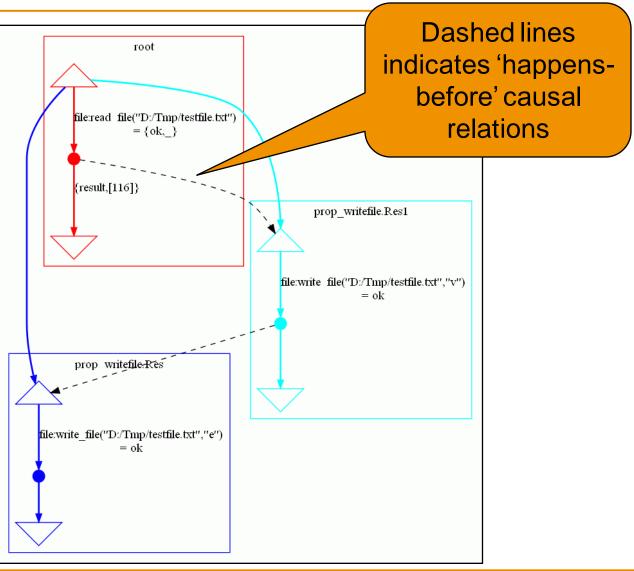
Matching module, function, arguments



```
13> eqc:quickcheck(writefile:prop_writefile()).
scheduling started
root spawns 'prop_writefile.Res' <0.1832.0>
root spawns 'prop_writefile.Res1' <0.1833.0>
root yields
                                          Aha! We are reading
'prop_writefile.Res1' yields
                                         the file before either of
'prop_writefile.Res' yields
                                          the writers has written
                                               anything!
root continues
root side-effect file:read_file(
    "D:/Tmp/testfile.txt") result in {ok,<<"x">>}
return value "x"
'prop_writefile.Res1' continues
'prop_writefile.Res1' side-effect file:write_file(
    "D:/Tmp/testfile.txt") result in ok
```



## Writefile – Visualization





- Q
- PAR spawns two processes, but a third process is also running in parallel to them!

Solution: Synchronize

```
-define(PAR(E1,E2),
    begin
        Self = self(),
        spawn(fun() -> E1, Self!done end),
        spawn(fun() -> E2, Self!done end),
        receive done ->
            receive done -> ok end
        end
        end
        end).
```



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## Exercise: Master-slave workers



- N workers: one master and N-1 slaves
- Process registry is used to identify the master
- Functionality in: master.erl
- Test case in: master\_eqc.erl
- There is a race condition in the code, which is hard to provoke with a test case
- Hint: use pulse\_side\_effect



- Two useful tricks
- Performance with **PULSE**
- A success story
- Availability of **PULSE**
- The future of **PULSE**

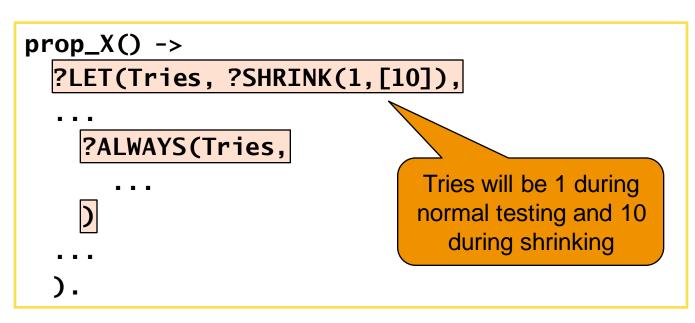




- What to do when shrinking doesn't work?
- Mostly important in larger more complex examples
- Even with pulse the counterexamples can be large
- Visualization is also useful, but graphs quickly gets quite large



- Q
- Idea 1: ?ALWAYS(N,Property)-macro tries the property N times, and fails if any of the tries fails
- Idea 2: Try the property many times while shrinking to increase the chance of hitting the bug





Comparing performance

- Used parallel map as benchmark
  - Short computations: fib(N) where N = 10-15
  - Long computations: fib(N) where N = 30-35
- Single core:
  - With longer computations **PULSE** is faster!
  - With short computations, communication dominates and **PULSE** is (much) slower
- Multi-core:
  - PULSE is always slower, since it only uses one of the cores.



- Performance is very application dependent
- Communication bound applications could be x100 slower.
- A 'normal' distributed application is likely to be x10 slower
  - Due to not using multi-core
  - and slower communication



- Real industrial example
- An optimized process registry
- Concurrency errors found by stress testing in 2006 (very large counterexamples)
- Nobody was able to track down the errors, so the component was shelved
- With PULSE we got shorter counterexamples
- With PULSE and the visualizer we could explain the error
- Described in paper at ICFP 2009



- Two versions:
  - Open source version (BSD license)
    - Developed at Chalmers
    - Work in progress (ProTest)
    - Not very user-friendly
    - No public release yet
  - Commercial version
    - Available as part of Quviq QuickCheck
    - Package PULSE in application
    - Integrates QuickCheck and PULSE



- Missing features (multi-node support etc)
- Improve shrinking of traces
- Re-write the core for a more modular design (already started)
- Support for testing timing dependent code
   (receive after X -> ...)
- Package and release open source version

