Erlang/OTP and how the PRNGs work



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What is random number generator?

Generating sequence of discrete numbers Two types of RNGs:

- •"True" RNGs: data from physical phenomena
- Pseudo RNGs: computed from a seed seed: initial vectors of tables of the internal state
- In Erlang/OTP, two modules of RNGs
 - crypto: OpenSSL API (NIFs from R14B)
 - •random: Wichmann-Hill AS183 (in 1982)



Requirements of RNGs

Uniform deviates

- Each of possible values is **equally probable**
- The building block for other deviates

Each number in the sequence must be statistically independent

- •Non-deterministic (unpredictable from past)
- •Non-periodic (no same sequence reappears)
- Fast enough to supply the demand
 - Generation speed could be a bottleneck



"True" RNG hardware examples

Collecting physical randomness / entropy

- Avalanche diode noise
- Free-running oscillators
- Atmospheric noise (random.org uses this)

Slow and expensive

- •The generation process does **not** guarantee if the output is equally probable and statistically independent
- •The output should be continuously verified and calibrated if the offset of the output deviate is large

See <u>RFC4086 Section 3 and Section 4</u> for the details

- Not repeatable (at least theoretically)
 - Practically used for seeding PRNGs for cryptography



Avalanche diode RNG circuit example

Example at https://github.com/jj1bdx/avrhwrng/ Speed: ~10kbps (or even slower for accuracy)

Arduino Duemilanove shield schematics for a hardware random number generator by Kenji Rikitake 6-APR-2009 Vin = +12V or +13.8V (+9V didn't work)+5V 0k 4.7mH 2.2u Vnd Vnd Vin Vnd 2.2k 10k <10k 47k 47k 470n 470n 100n 2.2u 100n 100nF: 50V ceramic AIN1/ AIN0 PD6 470nF: film Dig.6 Dig. 7 2.2uF: 50V electrolytic 2SC1815 x 2

2SC1815 x 2



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Arduino RNG looks like this





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Characteristics of pseudo RNGs

- Computed number sequences
 - Deterministic by definition
 - given the same seed, the same results show up
 - •Very long period but periodic anyway Longer period needed for larger scale application
 - Faster and more efficient than "True" RNGs
- Practical use: simulation and modeling
 - •random sampling / hashing / testing Load balancing, DHT, Monte Carlo method, etc.



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Cryptographic strength of PRNGs

Cryptographically-strong PRNGs must:
use the algorithm to prevent future data from the past generated data (with AES, SHA, etc.)
maintain collection of entropy pools from the various sources (network activities, etc.)
virtual machines: less entropy will be obtainable
secure the seeding process to prevent injection

attempts from the attackers

Use well-established methods for security

- •OpenSSL uses /dev/urandom on FreeBSD
- Accuracy transcends speed

Expect a lot of time to obtain sufficient random bits



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So what kind of RNGs in Erlang/OTP?

crypto module

- •rand_bytes/1, rand_uniform/2 OpenSSL API functions
- Always use crypto functions for security

random module: Wichmann-Hill AS183

- •period is very short (~ 7 x 10^12) [1]
- Written solely in Erlang

[1] B. A. Wichmann, I. D. Hill, "Algorithm AS 183: An Efficient and Portable Pseudo-Random Number Generator", Journal of the Royal Statistical Society. Series C (Applied Statistics), Vol. 31, No. 2 (1982), pp. 188-190, Stable URL: <u>http://www.jstor.org/stable/2347988</u>



Original AS183 code in FORTRAN

C IX, IY, IZ SHOULD BE SET TO INTEGER VALUES C BETWEEN 1 AND 30000 BEFORE FIRST ENTRY

- IX = MOD(171 * IX, 30269)IY = MOD(172 * IY, 30307)
- IZ = MOD(170 * IZ, 30323)
- RANDOM = AMOD(FLOAT(IX) / 30269.0 +
 FLOAT(IY) / 30307.0 + FLOAT(IZ) /
 30323.0, 1.0)

Source: Microsoft, Description of the RAND function in Excel http://support.microsoft.com/kb/828795



random module code of AS183

%% from lib/stdlib/src/random.erl
%% of Erlang/OTP R14B02

```
uniform() ->
    \{A1, A2, A3\} = case get(random_seed) of
                          undefined -> seed0();
                          Tuple -> Tuple
                      end,
    B1 = (A1*171) \text{ rem } 30269,
    B2 = (A2*172) \text{ rem } 30307,
    B3 = (A3*170) \text{ rem } 30323,
    put(random_seed, {B1,B2,B3}),
    R = A1/30269 + A2/30307 + A3/30323,
    R - trunc(R).
```



AS183 512x512 bitmap pattern test





(this looks well-randomized visually)

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What weak or bad RNGs will cause

- Vulnerability by predictable choice
 - •DNS UDP source port numbers
 - Precisely guessing cross-site state through JavaScript Math.random() method [2]
- Non-uniform bias on simulation
 - •Which may show up on a short-period RNG
 - Assumption of uniform deviate may fail

[2] A. Klein: Temporary user tracking in major browsers and Cross-domain information leakage and attacks, Trusteer, 2008, URL: <u>http://www.trusteer.com/list-context/publications/temporary-user-tracking-major-browsers-and-cross-domain-information-leakag</u>



rand(0,1) on PHP 5 Windows



Another popular example of bad RNG

%% originally from <u>http://xkcd.com/221/</u>
%% converted(?) to Erlang by Kenji Rikitake

```
-module(get_random_number).
-export([rand/0]).
```

```
rand() ->
% Chosen by fair dice roll.
% Guaranteed to be random.
4.
```

%% DO NOT USE THIS FOR A REAL APPLICATION!



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Issues needed to be solved

For security, crypto functions are must

- •In ssh module of R14B02 only AS183 found
- Longer period for non-crypto RNGs
 - •AS183 is good, but we need something better

$7 \ge 10^{12}$ period only holds ~81 days, if you generate 1 million random numbers for each second

- Faster generation for non-crypto RNGs
 - Faster algorithm for integer use
 - Maybe even faster with NIFs



SIMD-Oriented Mersenne Twister

A very good and fast PRNG

- •A revised version of Mersenne Twister
- very good = very long generation cycle typical: 2^19937 - 1, up to 2^216091 - 1 (depending on the internal state table size)
- Supporting SSE2/altivec SIMD features
- •Open source and (new) BSD licensed
- Implementations of (SF)MT available for:

C, C++, Gauche, Java, Python, R, etc.

URL: http://www.math.sci.hiroshima-u.ac.jp/~m-mat/MT/SFMT/index.html



So why SFMT on Erlang?

- The PRNG quality is well proven
 - survived the DIEHARD test
- It would be fast if implemented with NIFs
 - and that's what I have done
- SFMT RNG parameters are tunable
 - multiple algorithms generating independent streams possible if needed



PRNG enhancements with sfmt-erlang

SFMT implementation

Making the C code reentrant <u>http://github.com/jj1bdx/sfmt-extstate</u>

- of five (5) different periods with NIFs ~40 times faster than the non-NIF code it's even faster than random module
- Wichmann-Hill 2006 generator [3]
 - Called random_wh06 module
 - •A better RNG when NIFs can't be used

[3] B.A. Wichmann, I.D. Hill, Generating good pseudo-random numbers, Computational Statistics & Data Analysis, Volume 51, Issue 3, 1 December 2006, Pages 1614-1622, ISSN 0167-9473, DOI: <u>10.1016/j.csda.2006.05.019</u>.



SFMT Step 1: reentrant C code

Revised the SFMT reference code

- •Removed all static arrays The internal state table was defined as static the ultimate form of **the shared memory evil!**
- Removed the altivec and 64bit features no testing environment available
- SSE2 code removed

crashes for an unknown reason 128-bit alignment issue of enif_alloc()?

• Rewritten the code so that the internal state tables must be passed by the pointers Allowing concurrent operation of the functions



SFMT Step 2: pure Erlang version

- Literal translation from the revised C code
- SFMT itself can be written as a recursion
 - a[X] = r(a[X-N], a[X-(N-POS1)], a[X-1], a[X-2])
- Extensive use of head-and-tail lists
 - •Adding elements to the heads and do the lists:reverse/1 made the code 50% faster than using the ++ operator
- Still ~300 times slower than the C Code
 - •But it worked! (And that's what is important)





C to Erlang conversion tips

- Erlang integers are **BIGNUMs**
 - •Explicitly limit the result bit length by band each time after bsl and any other operation which may exceed the given C integer length
- Erlang bsr is arithmetic shift right
 - •e.g., -1 =:= -10 bsr 4 is true
- The array module object is **immutable**
 - •i.e., array:set/3 makes a modified copy



SFMT Step 3: writing a NIF version

- NIF modules are full of C static code
 - It's a shared-everything world as default
 - •When a NIF fails, it crashes the BEAM
- The fastest way to learn the NIF coding:
 - read the manual of erl_nif (under erts)
 - read the R14 crypto module
 - try first from smaller functions, step-by-step
 - •Use regression testing tools (e.g., eunit)



NIF programming tips

- It's hard-core C programming
 - Put all functions in the same .c file Remember how static scope works
 - Make the copy first before modifying a binary Without this you may face a heisenbug Erlang binaries are supposed to be immutable; so the content must stay unmodified!
 - •Learn the enif_*() functions first they will make the code efficient and terse



A case study: table handling on SFMT

- Case 1: list processing
 - •NIF: internal table -> integer list
 - •generating PRN by [head|tai]] operation
- Case 2: random access through NIF
 - •generating PRN each time by calling a NIF with the internal table and the index number
- Result: Case 1 is faster than Case 2
 - •on a 2-core SMP VM parallelism discovered?
 - •Lesson learned: profile before optimize



For the efficient Erlang + C coding

- Use a decent syntax highlighter
 - •erlang-mode and cc-mode on Emacs
- Use dev tools as much as possible
 - •eunit, fprof, rebar, escript, etc.
- Automate the documentation
 - •EDoc (for Erlang) and Doxygen (for C)
 - Learn the Markdown format
 It's much easier than to write HTML by hand



So how fast the SFMT NIF code is?

- Wall clock time of 100 * 100000 PRNs
 - on Kyoto University ACCMS Supercomputer Thin Cluster node (Fujitsu HX600)

AMD Opteron 2.3GHz amd64 16 cores/node RedHat Enterprise Linux AS V4 Erlang R14B01, running in a batch queue

sfmt: gen_rand_ list32/2	sfmt: uniform_s /1	random: uniform_s /1	random_wh 06: uniform_s /1	sfmt: gen_rand3 2_max/2	random: uniform_s /2	random_wh 06: uniform_s /2
240ms	2600ms	7110ms	11220ms	2440ms	7720ms	11790ms
x1.0	x10.8	x29.6	x46.8	x10.2	x32.2	x49.1
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speed of random .vs. random wh06

For 100000 calls of OWN time measured by fprof on R14B01 System details:

- reseaux: Core2Duo E6550 2.3GHz FreeBSD/i386 8.2-RELEASE
- leciel: Atom N270 1.6GHz FreeBSD/i386 8.2-RELEASE
- thin: Opteron 8356 2.3GHz RHEL AS V4 on amd64 This set of results suggest:
- The speed overhead from random to random_wh06 for CPUs with sufficient floating-point calculation support: < 10%
- On a CPU with lesser capability such as Atom, the overhead will increase to > 60%

	random:uniform_s/1	random_wh06:uniform_s/1	ratio of random_wh06 / random
reseaux	544.9ms	487.9ms	0.895
leciel	1400.3ms	2274.8ms	1.625
thin	309.2ms	331.2ms	1.071
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Total exec time of sfmt:gen_rand32_max .vs. SFMT internal table length



Total wall clock time[ms]

30

SFMT gen_rand32_list/2 performance



SFMT gen_rand_all/1 performance



OWN time for each call [microsecond]

Conclusion and future works (1)

- SFMT NIF: >x3 faster than AS183
- It's also better for simulation and modeling
- SFMT NIF behavior for period length
 - Shorter period causes larger calling overhead
 - •gen_rand32_list/2 exec time is ~ constant
 - •gen_rand_all/1 exec time is proportional to the internal state table size for a large period

random_wh06: 10~60% slower than AS183

• more room to optimize for slower CPUs Full 32bit integer is BIGNUM for 32bit Erlang VM



Conclusion and future works (2)

- Future works: exploring parallelism
 - •SFMT is inherently sequential/iterative
 - Looking for a new algorithm is needed There are parallelism-oriented PRNG algorithms Simplistic algorithms: LShift, XOR32, etc.
- Review of Erlang/OTP code for the secure usage of PRNGs is needed
 - •Very few network modules use crypto RNG
 - Analysis on Windows and other OSes needed



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It's more cost effective than building an amd64 test environment on an independent PC

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References

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- Press et al, Numerical Recipes (Third Edition), Cambridge Press, 2007, ISBN 9780521880688, Chapter 7 "Random Numbers", see <u>http://www.nr.com/</u>
- <u>http://www.diigo.com/user/jj1bdx/random</u>
 - > My bookmarks about random number generation
- Ferguson et al, Cryptography Engineering, Wiley, 2010, ISBN 9780470474242, Chapter 9 "Generating Randomness"

