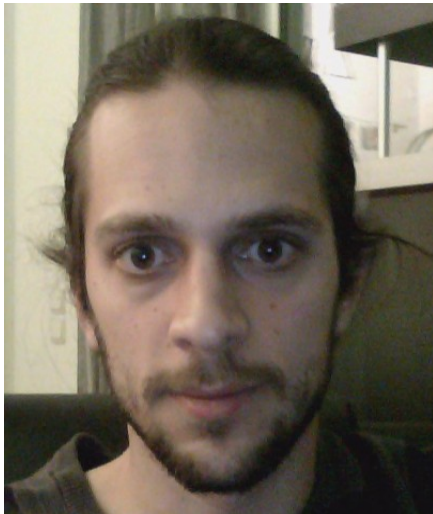




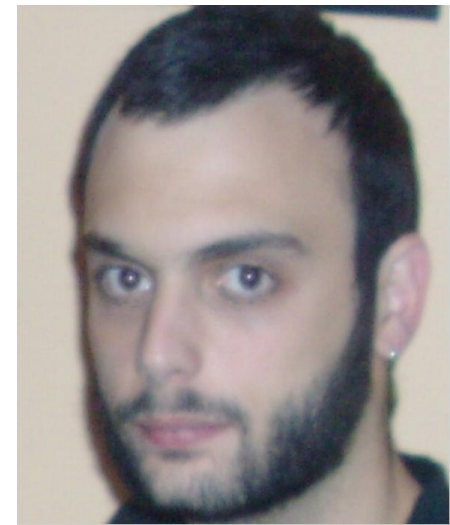
Kostis Sagonas



Chris Stavrakakis

joint work with

and



Yiannis Tsiouris

What is ErLLVM?

- A project aiming to provide multiple back-ends for the **High Performance Erlang (HiPE)** native code compiler of Erlang/OTP ...
- ... using the **Low Level Virtual Machine (LLVM)** compiler infrastructure ...
- ... in order to *improve the performance* of Erlang applications ...
- ... and *ease the maintenance* of its native code compiler.

This talk

- Overview and design
 - HiPE native code compiler
 - LLVM compiler infrastructure
- Architecture and implementation of ErLLVM
 - LLVM extensions
 - New HiPE component
- Performance evaluation
 - vs. BEAM
 - vs. HiPE
 - vs. Erjang
- Current status and future work

HiPE: High Performance Erlang

- Project at Uppsala University started in 1997
- Developed the native code compiler for Erlang
 - Component of Erlang/OTP since 2001
- *A mature* compiler that is *robust* and produces *reasonably efficient* code
- Back ends for
 - SPARC V8+
 - x86 and x86_64 (AMD64)
 - PowerPC and PowerPC64
 - ARM

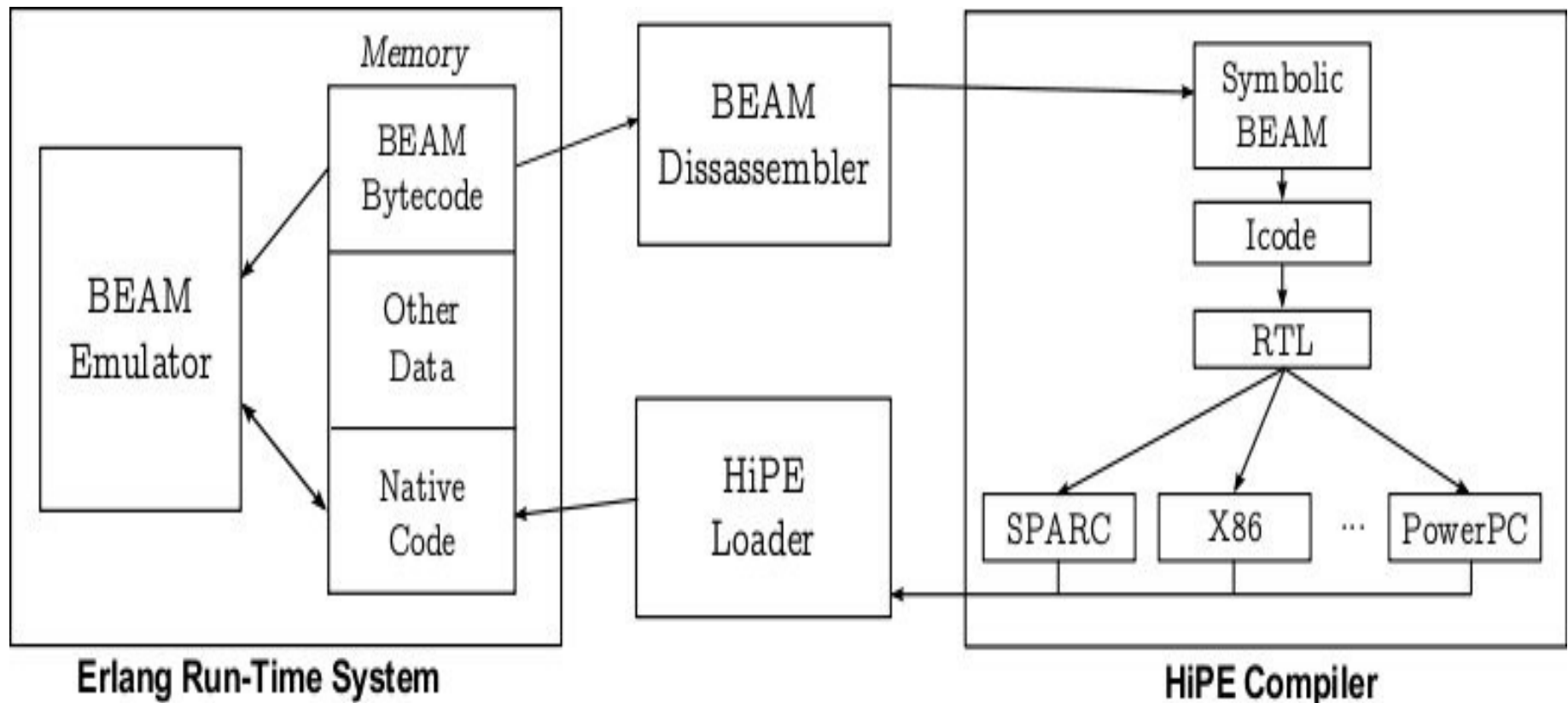
LLVM: Low Level Virtual Machine

- Collection of industrial strength compiler technology
 - Language-independent *optimizer* and *code generator*
 - Many optimizations, many targets, generates good code
 - Clang C/C++/Objective-C front end
 - Designed for speed, reusability, compatibility with GCC
 - Debuggers, “binutils”, standard libraries
 - Providing pieces of low-level tools, with many advantages
- High-level portable LLVM assembly
 - RISC-like instruction set; static type system; SSA form
 - Three forms: human-readable, on-disk, in-memory

Why LLVM?

- Used as a static or JiT compiler and for static analysis
- State-of-the-art software with very active community of developers
- A new compiler = glue code + any components not yet available. Allows choice of the right components for the job, e.g. register allocator, scheduler, optimization order.
- Supports many architectures: x86, x86_64, ARM, PowerPC, SPARC, Alpha, MIPS, Blackfin, CellSPU, Mblaze, MSP430, XCore, ...
- Open source with a *BSD-like License* and many contributors: industry, research groups, individuals

HiPE Architecture in Erlang/OTP



HiPE's Icode

```
1 lists:sum/2(v1, v2) ->
2 %% Info: ['Not a closure', 'Leaf function']
3 1:
4     v3 := phi({1, v1}, {3, v6})
5     v4 := phi({1, v2}, {3, v7})
6     _ := redtest() (primop)
7     goto 4
8 4:
9     if is_cons(v3) then 3 (0.50) else 6
10 3:
11     v5 := unsafe_hd(v3) (primop)
12     v6 := unsafe_tl(v3) (primop)
13     v7 := '+'(v5, v4) (primop)
14     goto 1
15 6:
16     if is_nil(v3) then 5 (0.50) else 2
17 5:
18     return(v4)
19 2:
20     v10 := function_clause
21     fail(error, [v10])
```

```
sum([H|T], A) -> sum(T, H+A) ;
sum([], A) -> A.
```

Figure 2. The Icode (in SSA form) of a `lists:sum/2` function

HiPE's RTL

```
1 {lists,sum,2}(v19, v20) ->
2 ;; Leaf function
3 ;; Info: []
4 .DataSegment
5 .CodeSegment
6 L1:
7     v21 <- v20
8     v22 <- v19
9     goto L2
10 L2:
11     v23 <- phi({1, v21}, {11, v32})
12     v24 <- phi({1, v22}, {11, v33})
13     %fcalls <- %fcalls sub 1 if lt then L4 (0.01) else L6
14 L4:
15     <- suspend_0() then L6
16 L6:
17     r25 <- v24 'and' 2 if eq then L7 (0.50) else L8
18 L7:
19     v26 <- [v24+-1]
20     v27 <- [v24+7]
21     r28 <- v26 'and' v23
22     r29 <- r28 'and' 15
23     if (r29 eq 15) then L13 (0.99) else L12
24 L13:
25     r34 <- v23 sub 15
26     v35 <- v26 add r34 if not_overflow
27     then L11 (0.99) else L12
28 L11:
29     v31 <- phi({13, v35}, {12, v30})
30     v32 <- v31
31     v33 <- v27
32     goto L2
33 L12:
34     v30 <- '+'(v26, v23) then L11
35 L8:
36     if (v24 eq -5) then L15 (0.50) else L16
37 L15:
38     return(v23)
39 L16:
40     v36 <- atom_no('function_clause')
41     <- erlang:error(v36)
42     return(15)
```

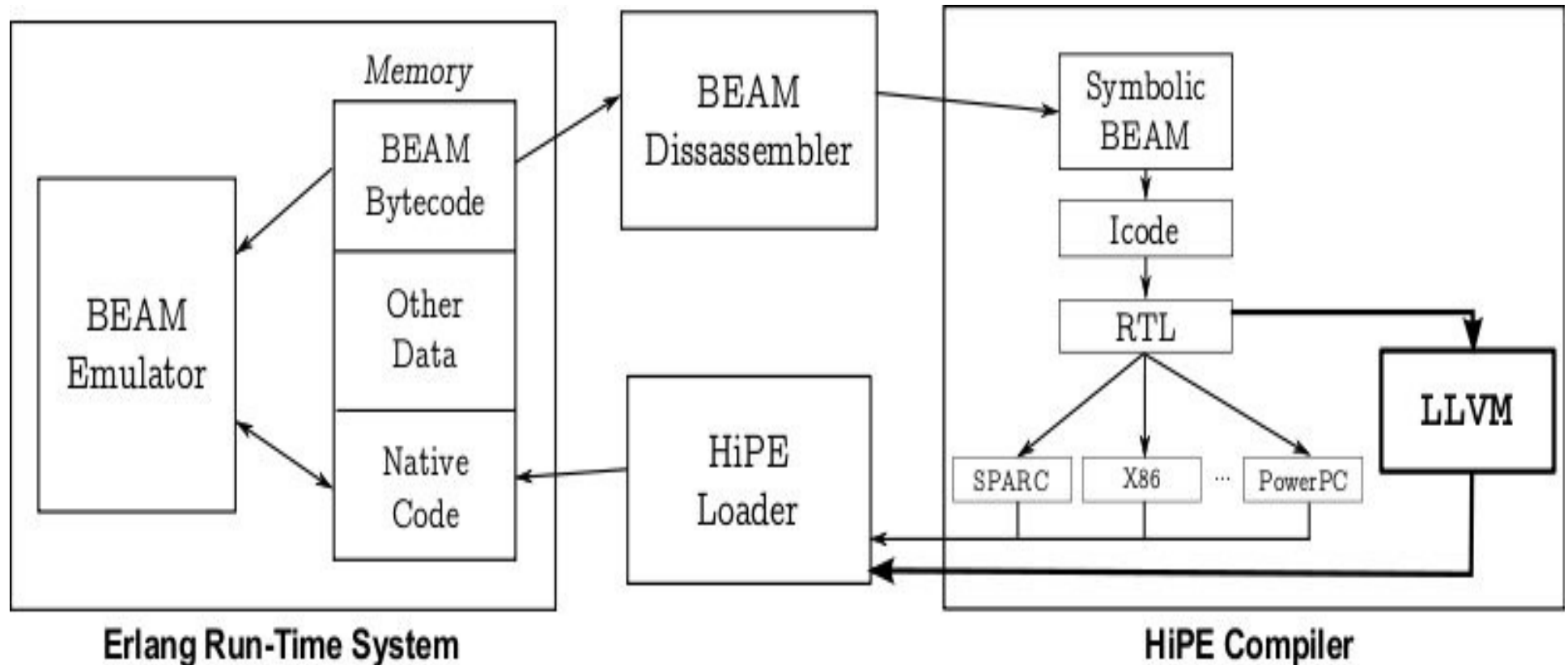
Structure of HiPE's Back-ends

- Register allocation
 - Many choices; default iterated register coalescing
- Frame management
 - Check for stack overflow
 - Set up frame
 - Create stack descriptors
 - Add “special” code for tailcalls
- Code linearization
- Assembly

Why use LLVM as a Back-end?

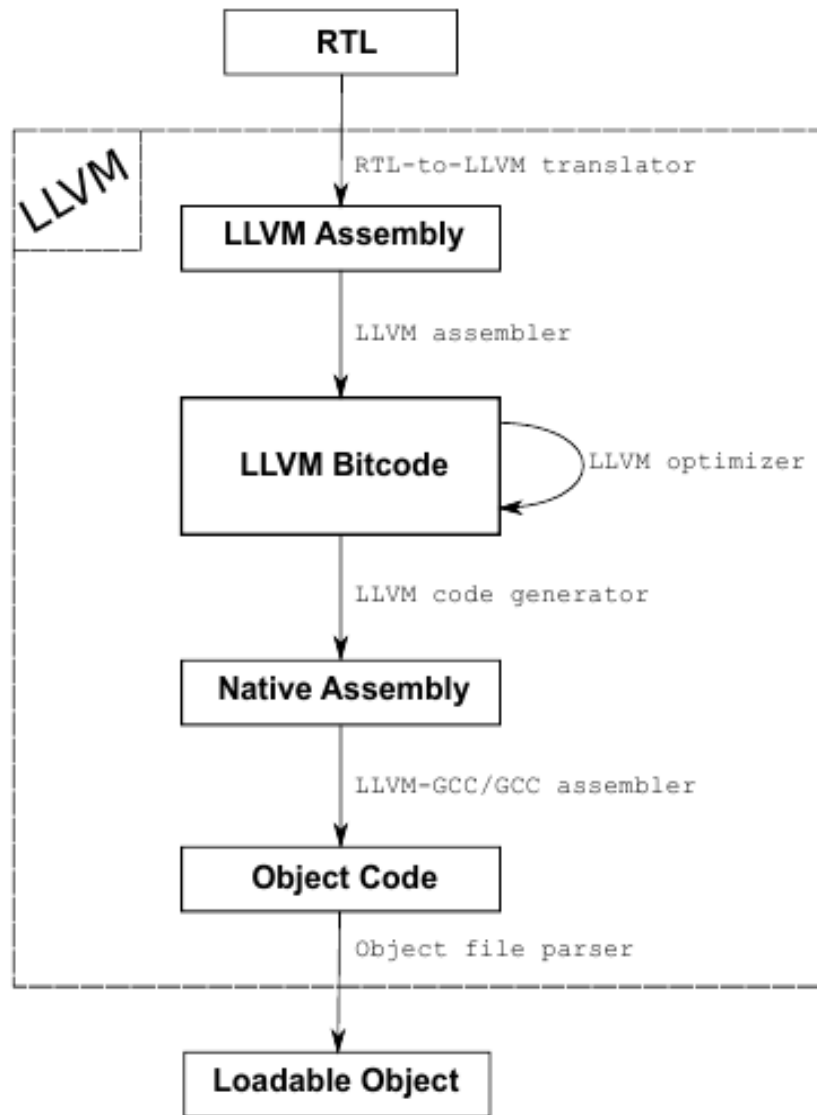
- Curiosity: perform a research experiment
- Easier maintenance of existing back-ends
 - One instead of six
 - Small-sized, straightforward code
 - Outsource implementation and further optimization
- Get more back-ends “for free” (well, almost...)
- Possibly improve performance
 - Outsource target-related optimizations

HiPE Architecture in ErLLVM



- Use existing HiPE Loader and ERTS support
 - Be **ABI compatible!**

The LLVM Component



`hipe_rtl2llvm` Create human-readable LLVM assembly (.ll)

`llvm-as` Human-readable assembly (.ll) → LLVM bitcode (.bc)

`opt` Optimization *Passes*, supports standard groups (-O1, -O2, -O3) (.bc → .bc)

`llc` Bitcode (.bc) → Native assembly (.s), impose rules about memory model, stack alignment, etc.

`llvm-gcc` Create object file (.s → .o)

`elf64_format` Extract executable code and relocations

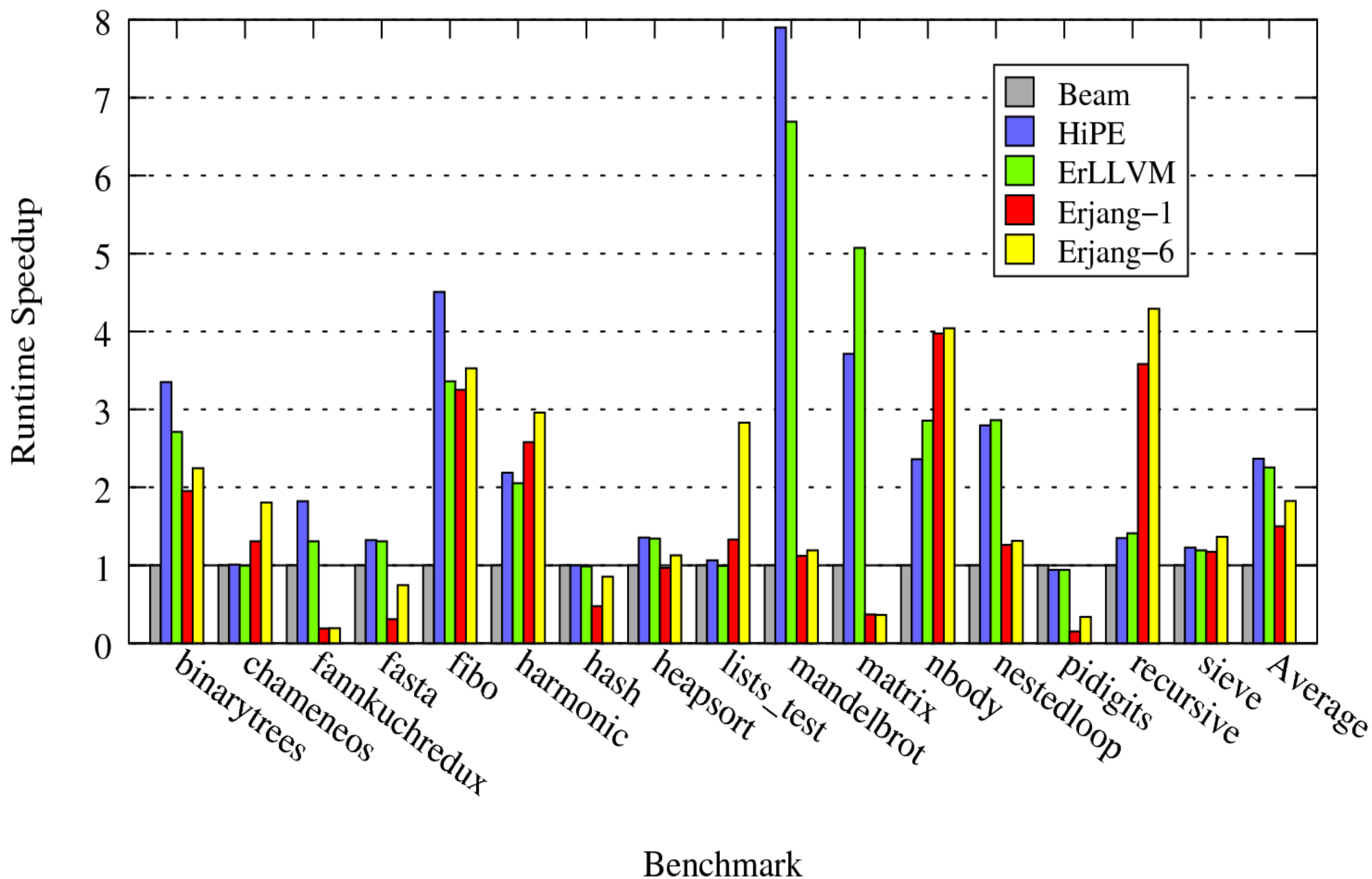
Subtle Points of Using LLVM

- Calling convention
 - VM “special” registers, args and return values
 - callee-/caller-save registers, callee pops args
- Explicit frame management
 - In-lined code for stack overflow checks in assembly prologue
- Stack descriptors
 - Exception handling
 - **Precise** garbage collection

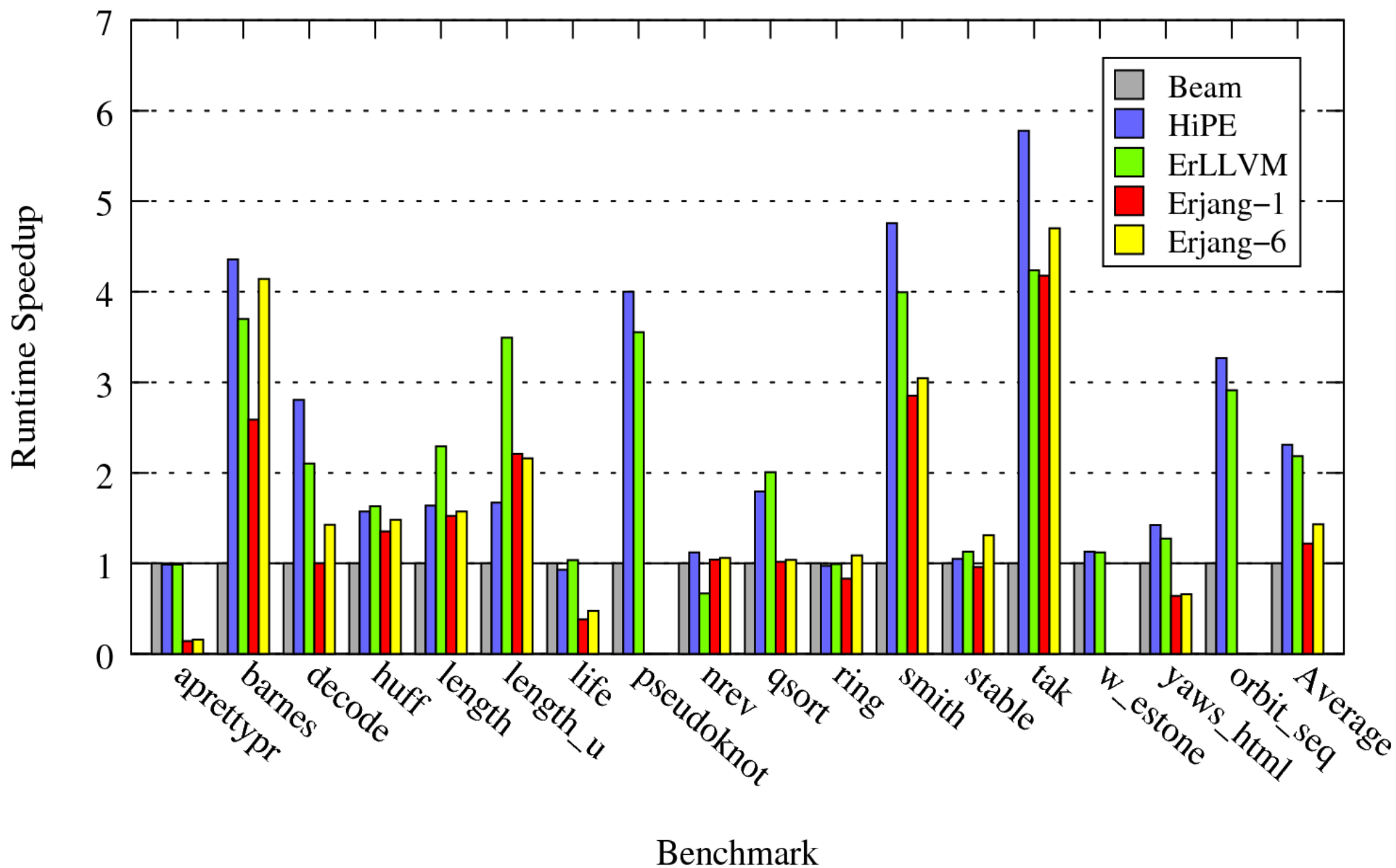
Current Status of ErLLVM

- Patches to LLVM
 - Custom calling convention & register pinning
(Already part of LLVM's code base!)
 - GC plugin to write GC information in object file
Use `elf_format` to parse .o file and extract the info
 - Function pass to emit custom prologue
- New HiPE component on top of R15B “maint”
 - Support for x86 and x86_64
 - Support for *accurate GC*: mark stack slots not live when variables that “inhabit” them are no longer live
 - About 5000 LOC
- Very robust and ready to use in production!

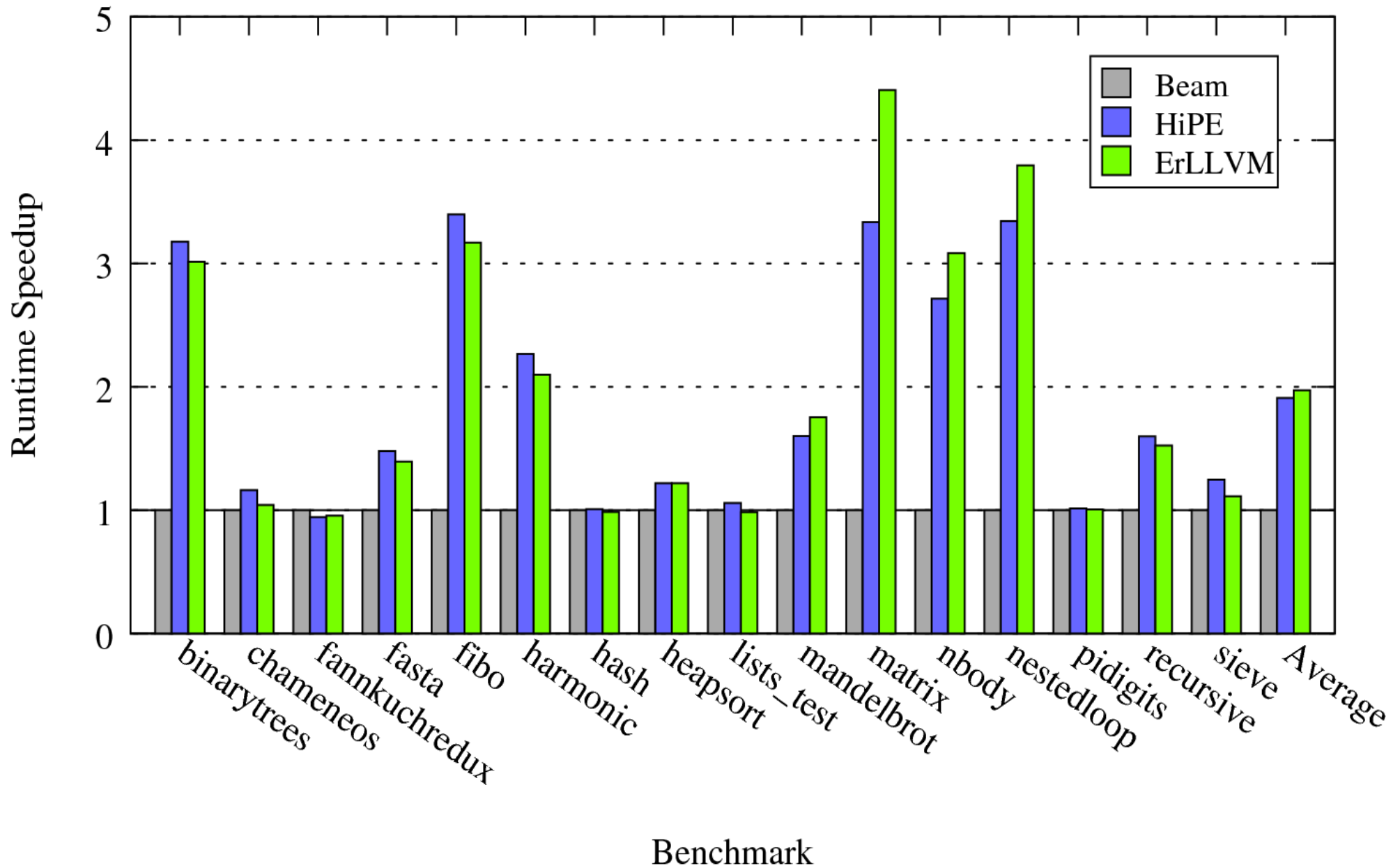
Performance on x86_64 (Beam=1)



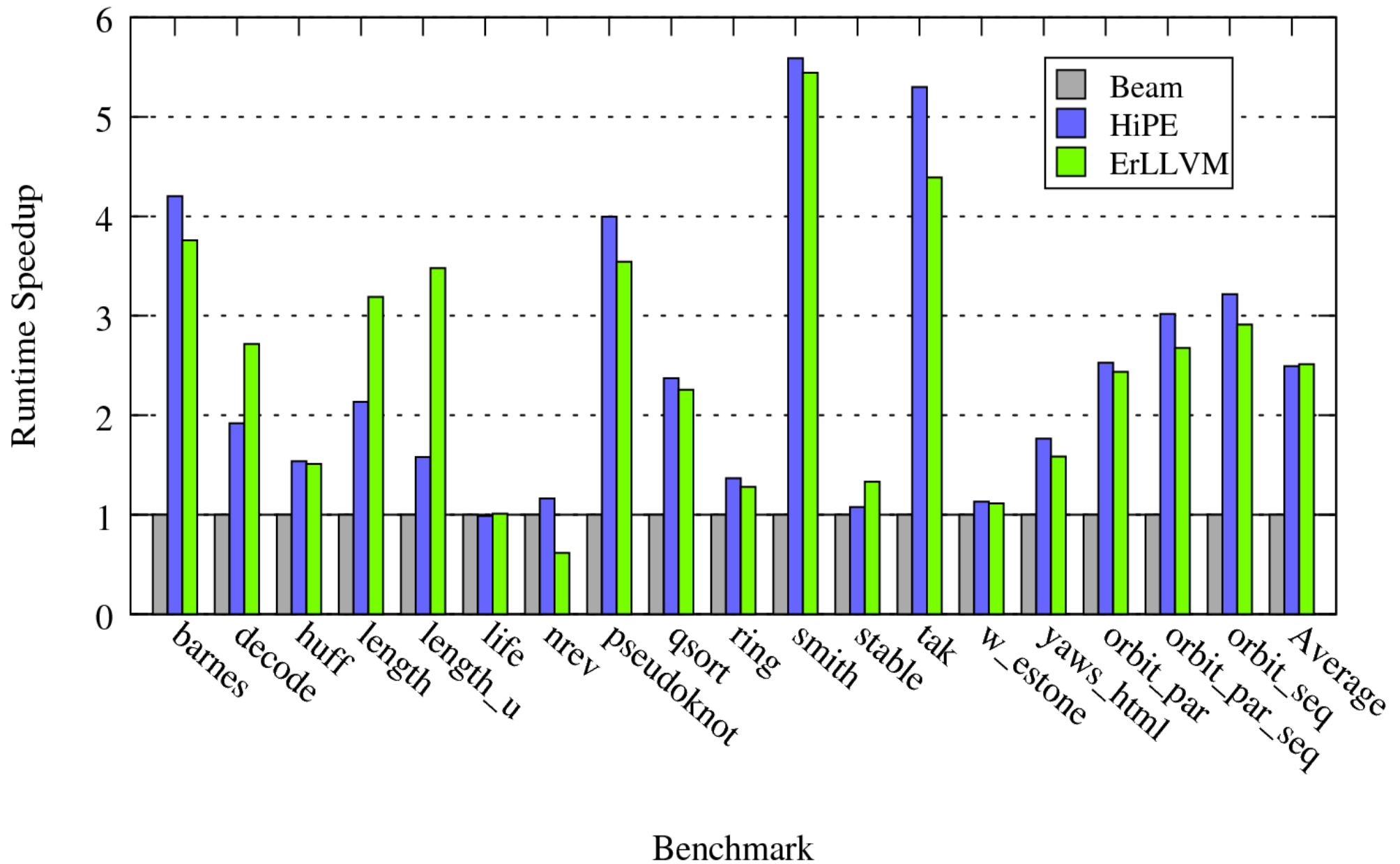
Performance on x86_64 (Beam=1)



Performance on x86 (Beam=1)



Performance on x86 (Beam=1)



Now what?

Demo time!

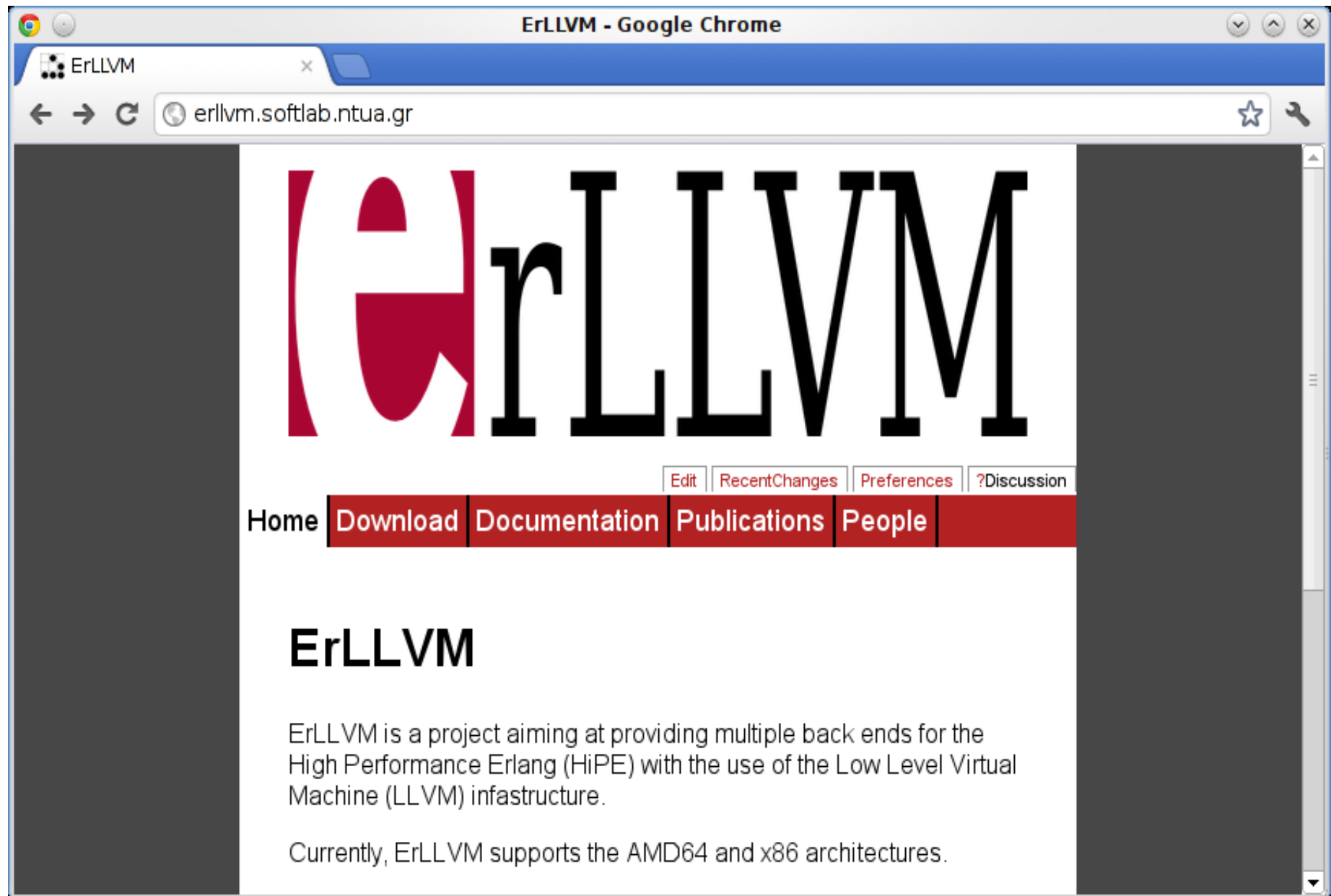
Current Status: Pros

- **Complete & robust:** handles all Erlang programs
- **ABI compatible:**
 - smooth integration with BEAM and HiPE code
- **Performance:**
 - much better than BEAM
 - almost as good as HiPE
- Smaller and simpler code base for the back-ends
- Possibility to target more architectures
- LLVM back-end improvements now also improve performance of Erlang applications!

Current Status: Cons & Future Work

- Cons:
 - Need to download and install custom LLVM
 - Slightly longer compilation times
 - Not a problem in practice
 - Erlang LLVM bindings to the rescue??
- Future Work
 - Finish pushing LLVM patches upstream
 - Port to ARM
 - Improve GC support

Where can I find ErLLVM?



Users are now welcome!

- Install following the instructions at:
<http://erllvm.softlab.ntua.gr>
- Grab code from `github`
- Test and measure!
- Report experiences
- Contribute to the project

