

# FOTOS – Fail Over and Take Over System

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# Erlang/OTP selling point

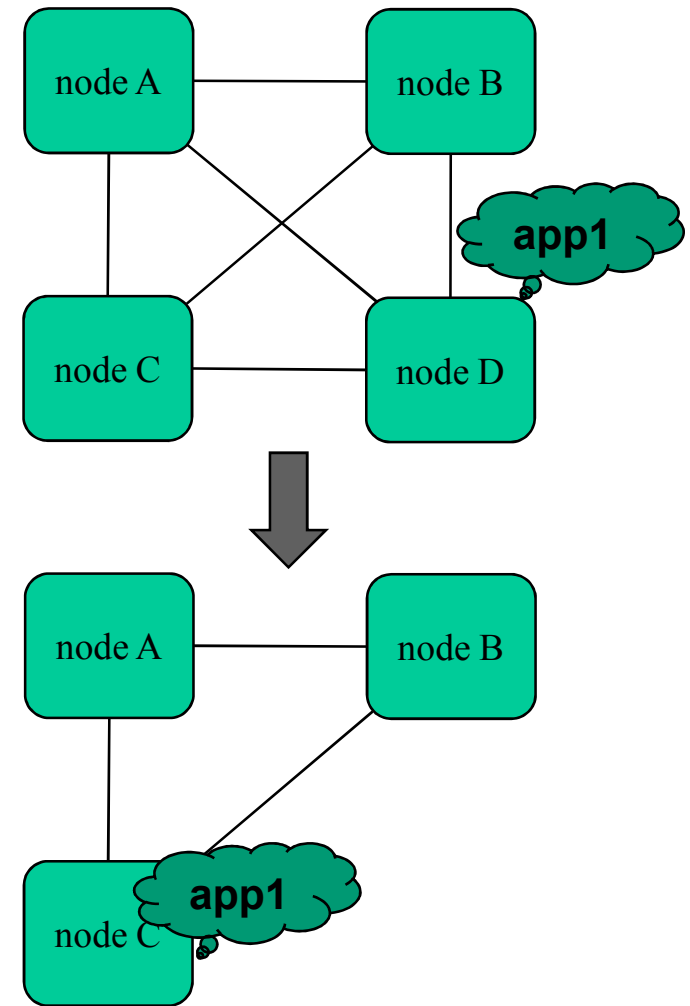
- It is very easy to create multi node systems with load balancing and failover in Erlang/OTP!

or...



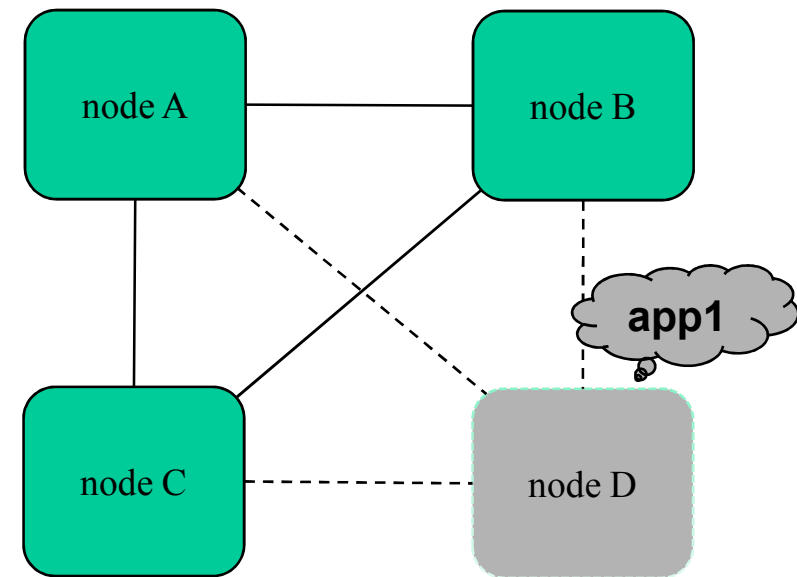
# It is very easy to create...

```
[{kernel,  
  [{distributed,  
    [{app1,  
      [nodeD,  
        nodeC,  
        nodeB,  
        nodeA]}]}]  
}]}
```



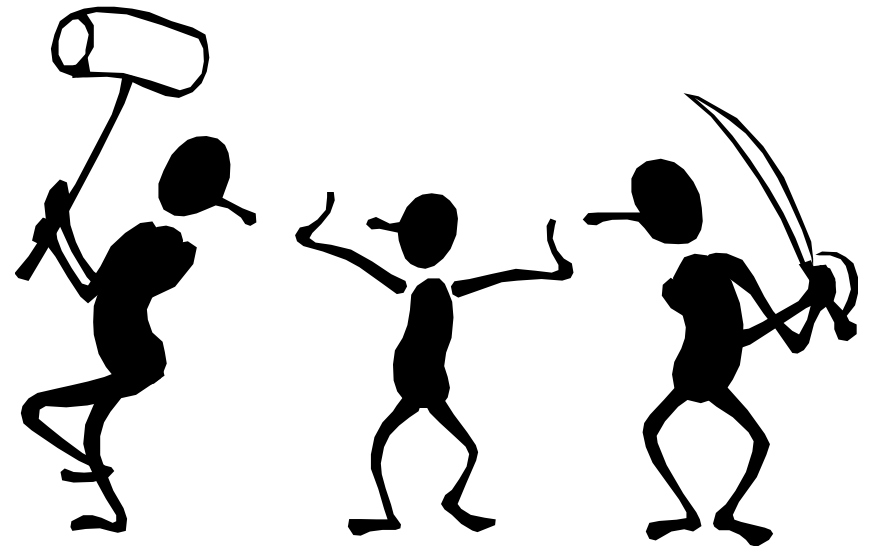
# What is failover?

- The ability of peer nodes to detect the failure of a node and resume its responsibilities.



# Failover...

- Detect failing node, and
- Decide who resumes which responsibilities.



# Failover...

- Decide who resumes which responsibilities.

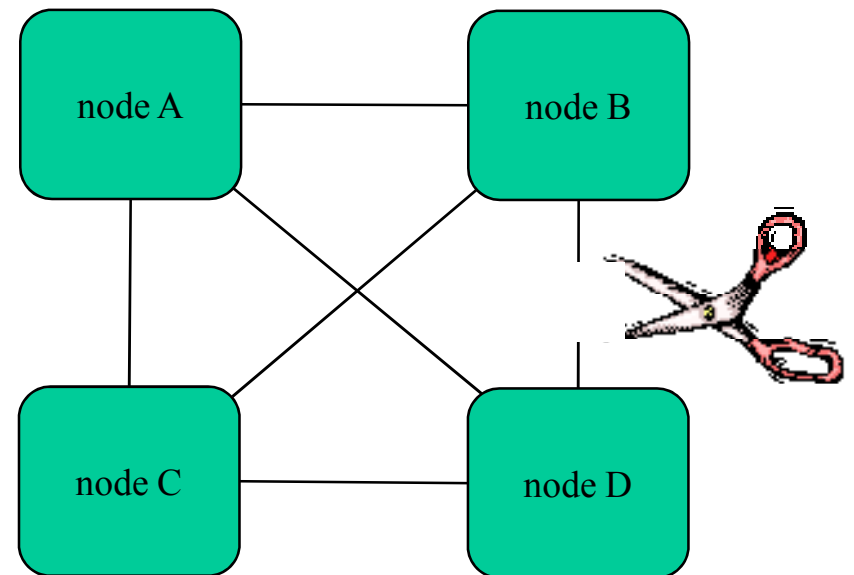


**All nodes must agree on which nodes are in fact part of the network!**



# It is very easy, or...

- The distributed application controller does not handle network failures.



# FOTOS

- **Monitoring of the network and maintain a network of operational nodes.**
  - **Without having to elect a leader.**
- **A simple database with vector clock controlled updates.**

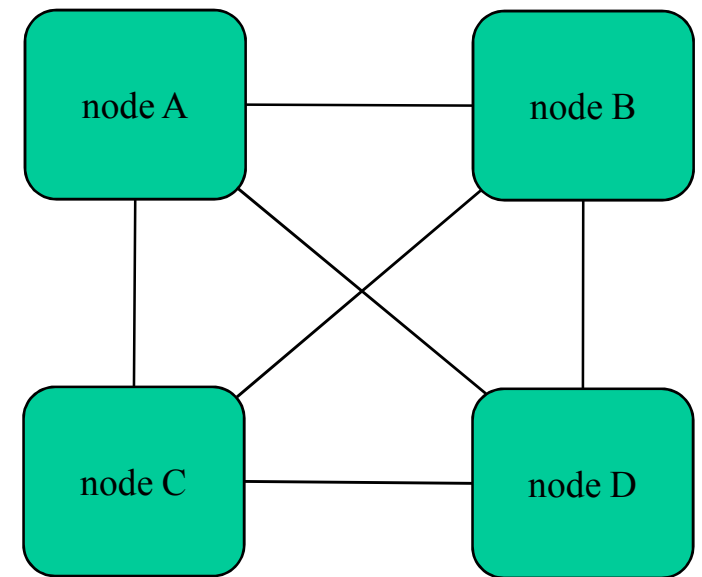
# Assumptions

- All nodes are loyal and follows “the rules”.
- Messages between two nodes are delivered in the same order as sent.
- External means of detecting a partitioned network.

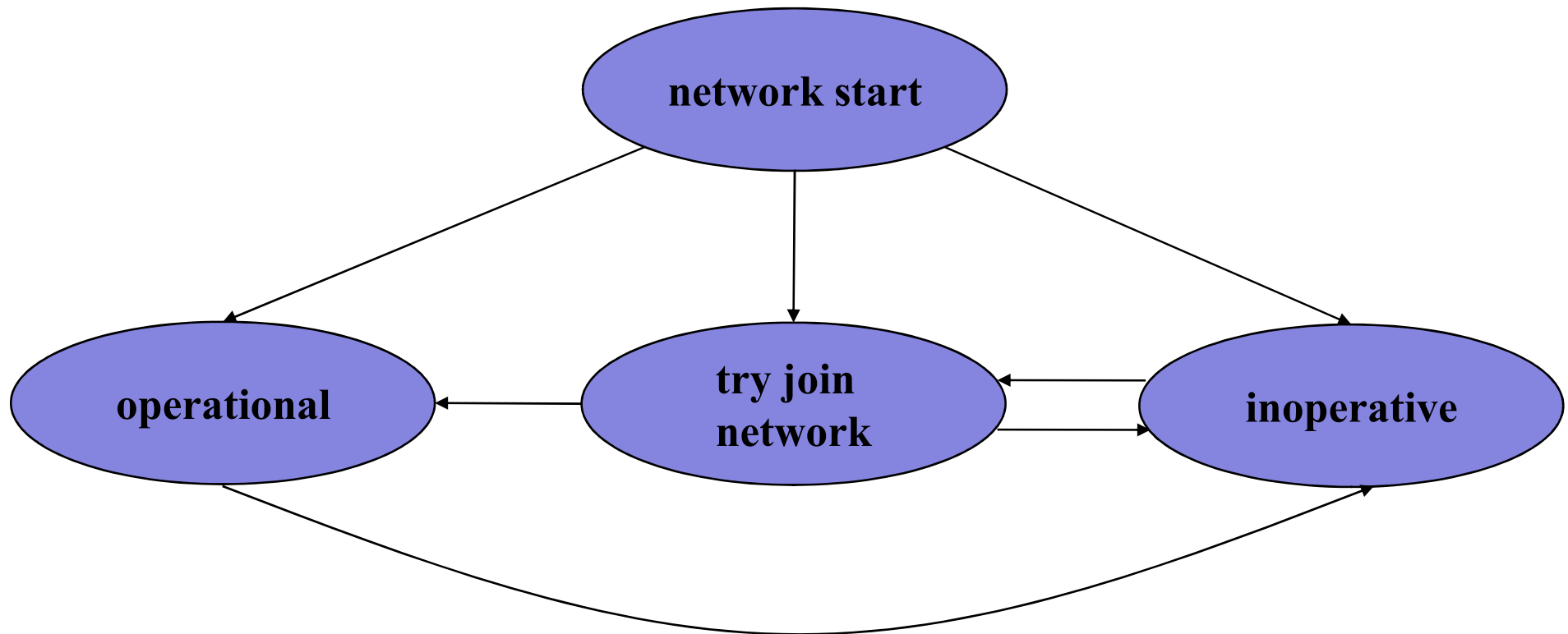


# Monitoring of the network

- **A network of operational nodes**
  - A fully connected mesh of Erlang nodes.
  - Only nodes that can "see" all other participating nodes may participate.

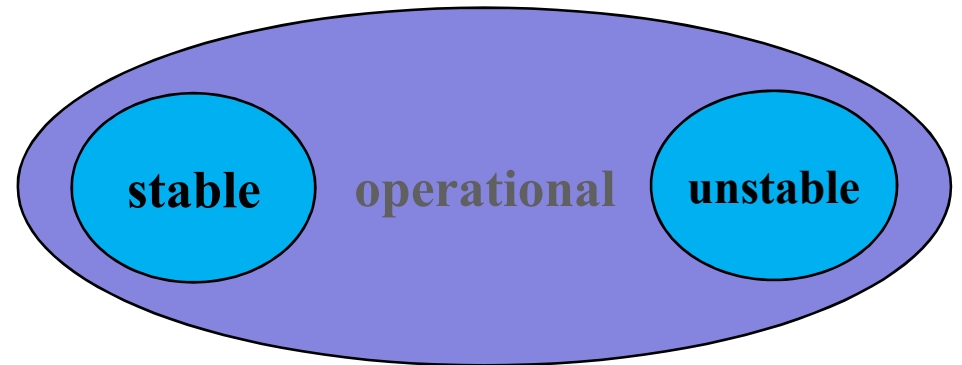


# The node state machine



# The node state machine

- An operational node may be either stable or unstable.



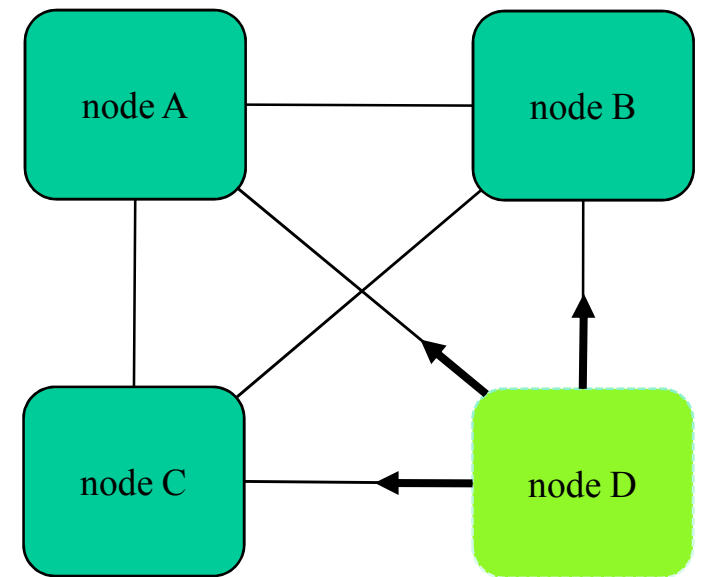
# Network start

- The first thing a newly started node does is to try to start a network of operational nodes.
  - It contacts all known nodes to see if they want to help create a network of operational nodes.



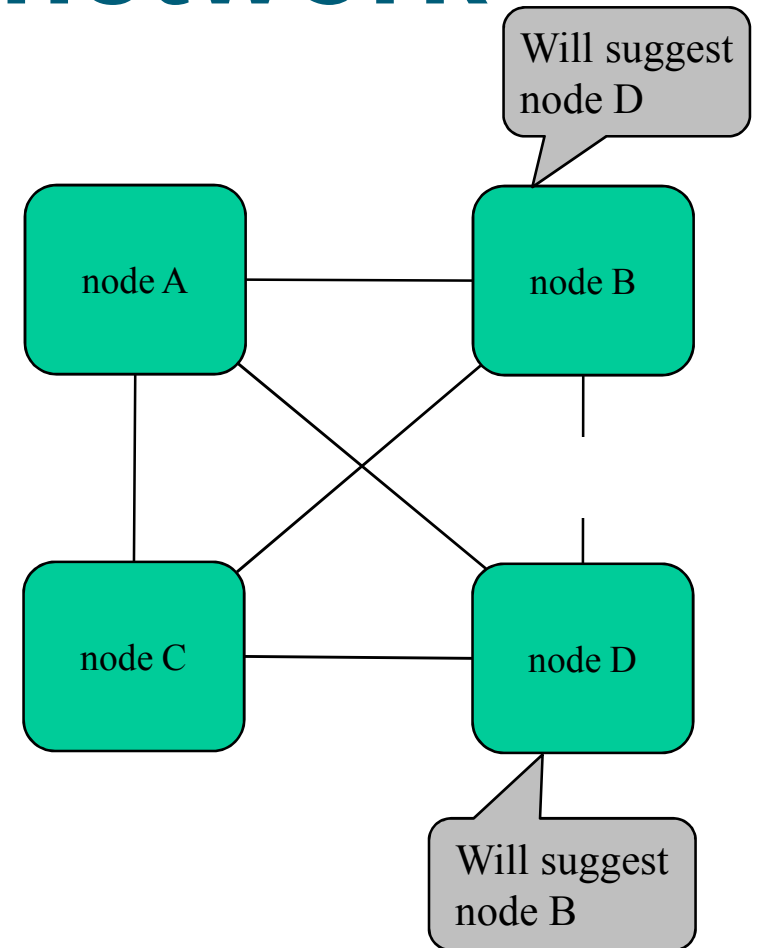
# Joining a network of operational nodes

- A newly started node tries to join a network if it "sees" all nodes.
- An operational node may only object to a JOIN REQUEST if it does not include exactly all nodes that are operational, or
- If it is participating in another procedure currently.

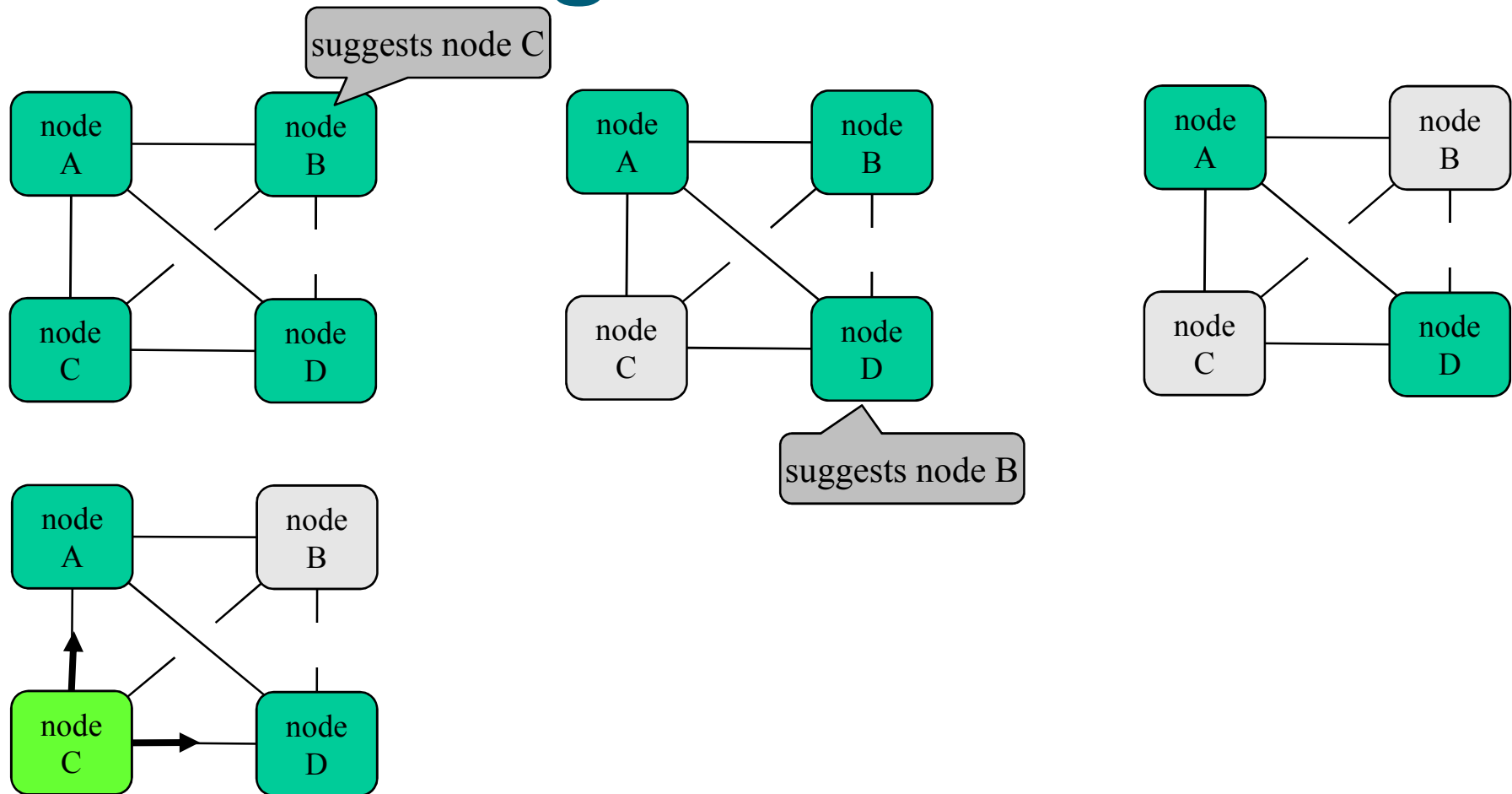


# Monitoring of the network

- A node that detects a peer node failure suggests its removal.
  - A node may not be removed unless all other nodes (except the failed node) agrees.



# Monitoring of the network



# Single failure optimization

- A node may only run one procedure at a time.
  - Removing or joining one node at a time.
- If multiple simultaneous failures, no guarantee the result is the most optimal possible network.

# Network unstable

- **As soon as a node learns of that a participating node is not visible from one or several nodes – the network is deemed unstable.**
- **In first “person” (‘DOWN’ message).**
- **“Rumor” from another node that wishes to remove the now failed node from the network of operational nodes.**

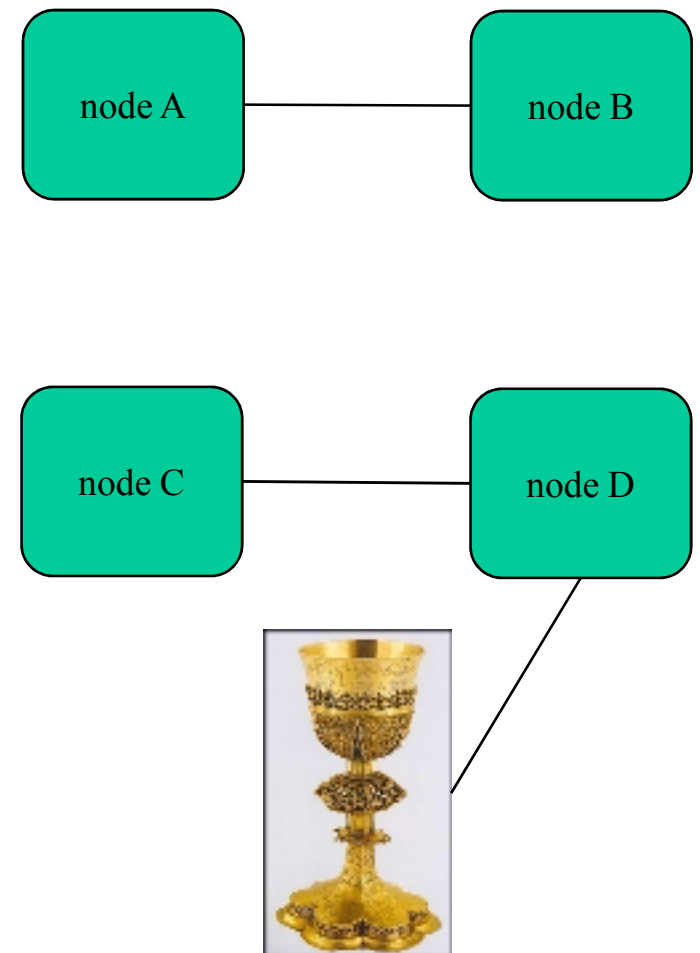
# There is no magic!

- You must still make fail-over and load balancing decisions!
  - Based on stable/unstable.
  - Based on changes of the set of participating nodes.



# You must provide a “holy grail”

- In case of a completely partitioned network – a “holy grail” is needed to decide which network is the “real network”.



# Some spin-offs

- **2 Phase Commit**
  - 2 phase commit with rendezvous
- **A Vector Clock**

# FOTOS summary

- **Subscribe to stability status.**
- **Subscribe to changes in the set of operational nodes.**
- **Optimized for single failure.**
- **A simple database with vector clock protected updates.**