

More than a century of programming

Joe Armstrong
Robert Virding
Mike Williams



History (early 1980's)

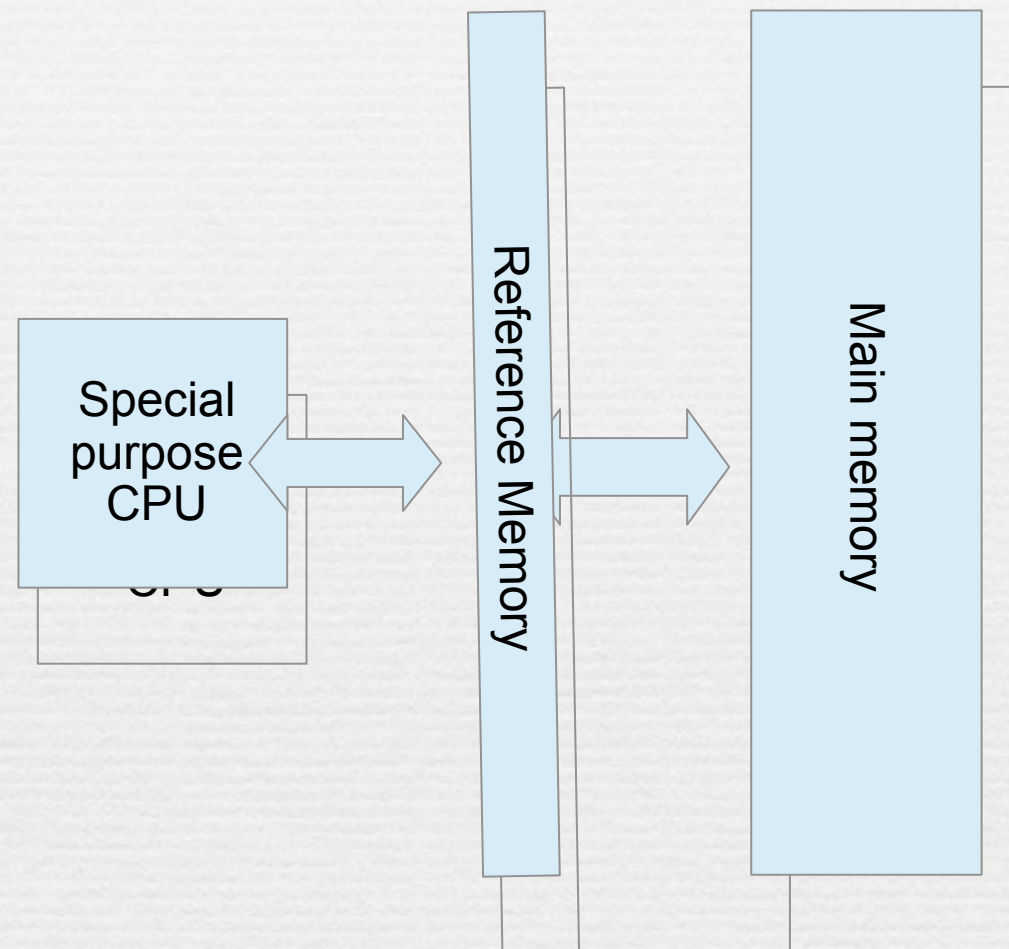
- ❧ Ericsson's “best seller” was AXE telephone exchanges (switches).
- ❧ Required large effort to develop and maintain software.
- ❧ The “job” was to make programming these types of application easier, but keeping the same characteristics



AXE Exchange, 1985
Guangzhou, China

AXE SW Characteristics

- Massive concurrency (thousands of transactions)
- Array bounds and pointer checking in hardware
 - No wild pointers
 - Size changes of statically allocated arrays
 - Re-arrange memory
- Change code at runtime
- Modular
- Error handling and transactions



Special purpose synchronously duplicated AXE hardware

Do the same as AXE, but use

- Conventional processors, easy portability to new processors
- Conventional operating system (type UNIX)
- Distributed multi-processor system enabling scalability (more processing power == more processors)

And

- Make software development effort significantly easier.

How hardware changes how we think

- ❧ Large memory/small memory
- ❧ Always on line
- ❧ Fast CPU
- ❧ Parallel hardware
- ❧ Mobility

If the hardware doesn't
change the software
won't change

Hardware didn't change
much in 1986-2004 so
the software didn't
change much

*apart from clock speed and memory capacity
still Von-Neumann non-distributed non-
connected*

1980 - 1984

- The state of computer science
- 3 MHz clock
- 80 MB disks
- 4MB memory ...
- (language reflect the hardware of the time when they were developed)

Problem Domain

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1. Actions must be performed at a certain point in time or within a certain time.
2. System may be distributed over several computers.
3. The system is used to control hardware.
4. The software system is very large.
5. The system exhibits complex functionality such as feature interaction.
6. The systems should be in continuous operation over many years.
7. Software maintenance (reconfiguration etc.) should be performed without stopping the system.
8. There are stringent quality, and reliability requirements.
9. Fault tolerance both to hardware failures, and software errors, must be provided.
10. The system *must* be able to handle very large numbers of concurrent activities.

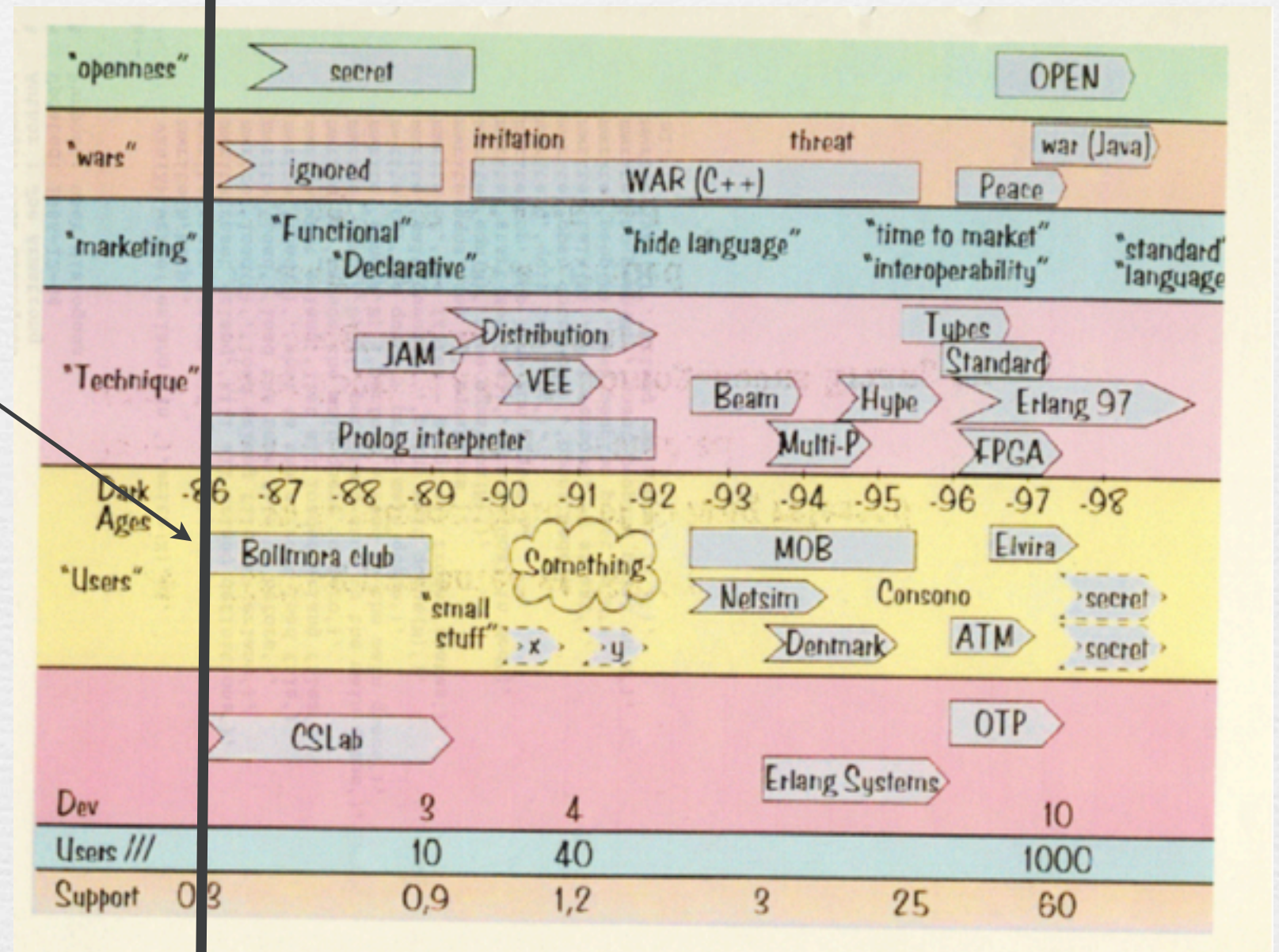
Bjarne Däcker. November 2000 – Licentiate thesis

Reality check

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Ericsson Business Networks AB



MD110

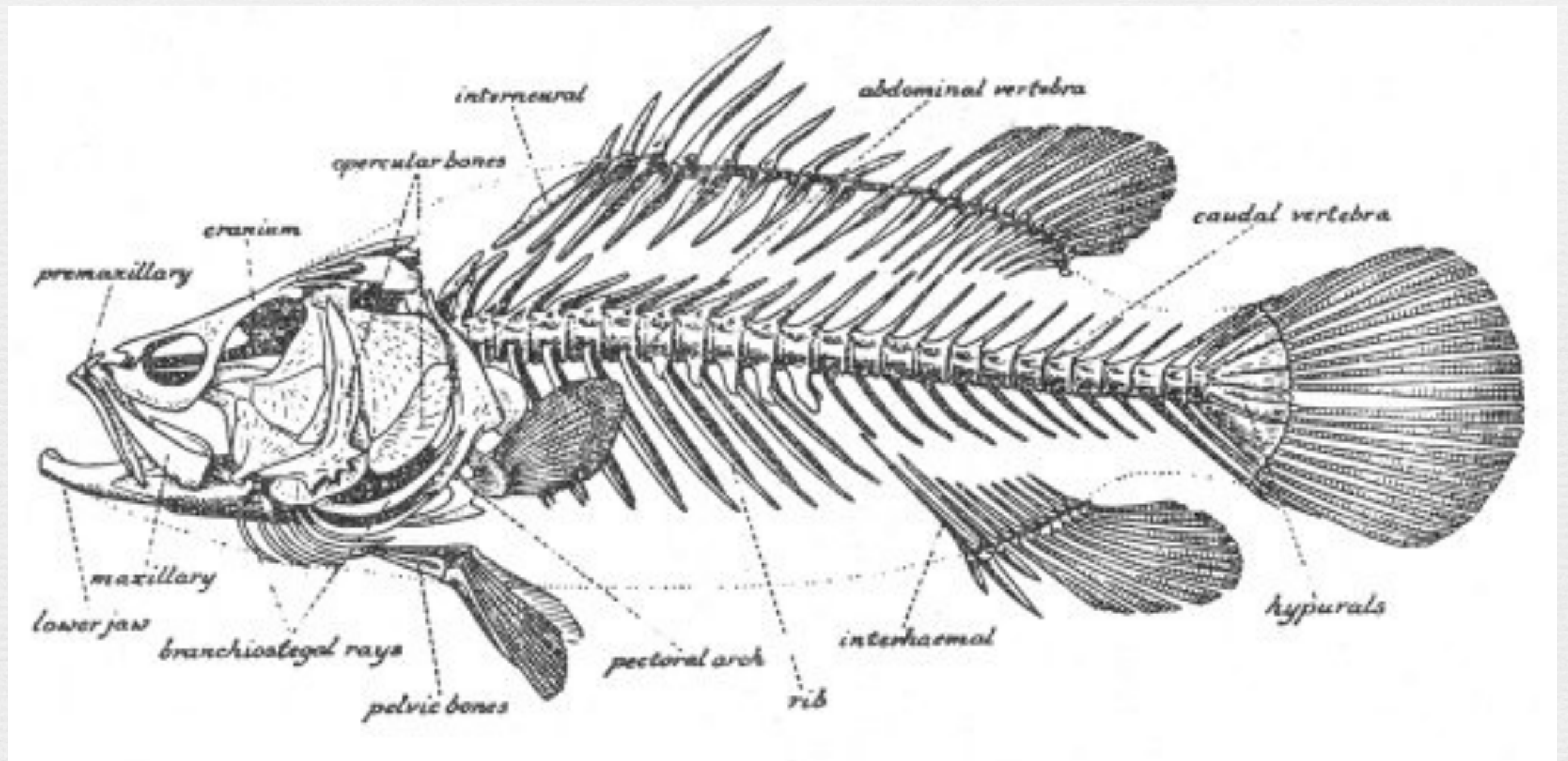
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“It’s the hardware stupid”



“Hardware is funny” - RV

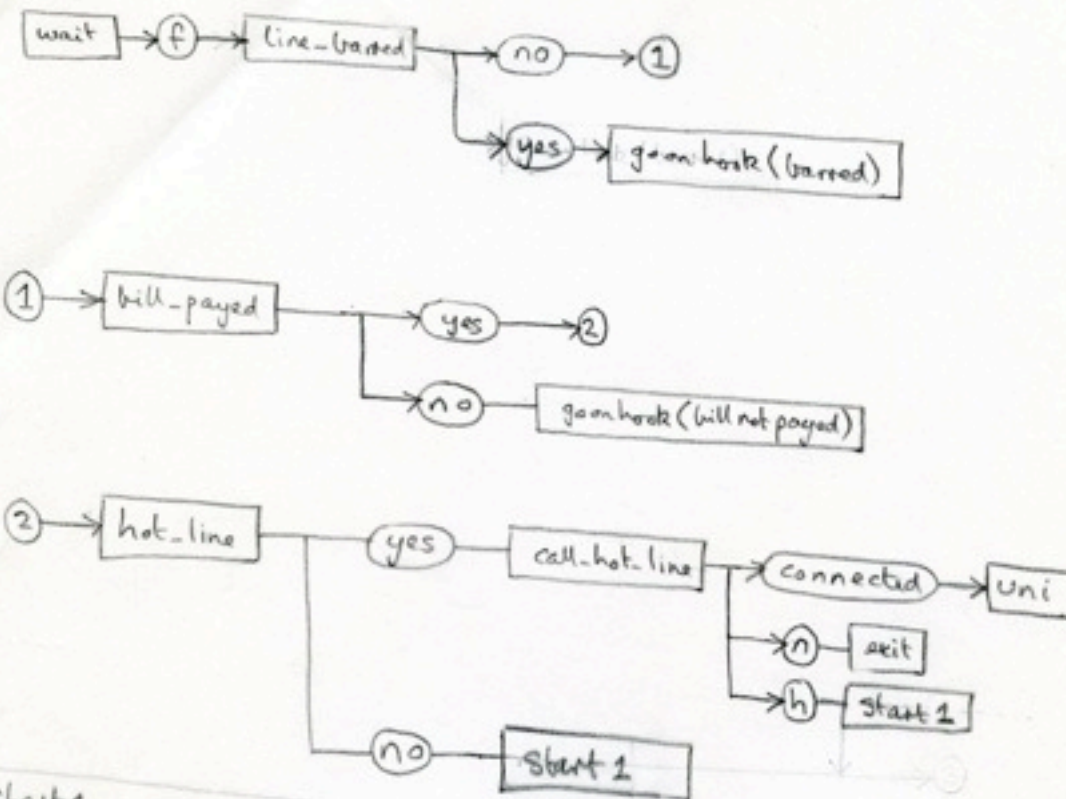
Erlang version 0



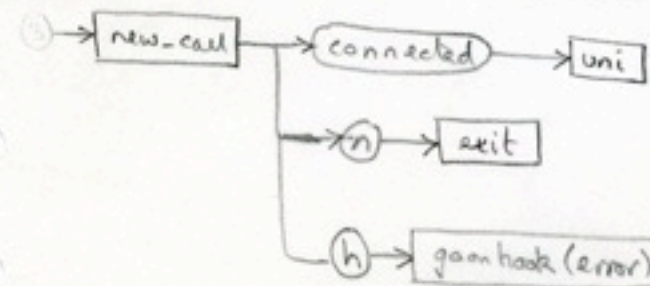


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Start (start it off)



Start 1

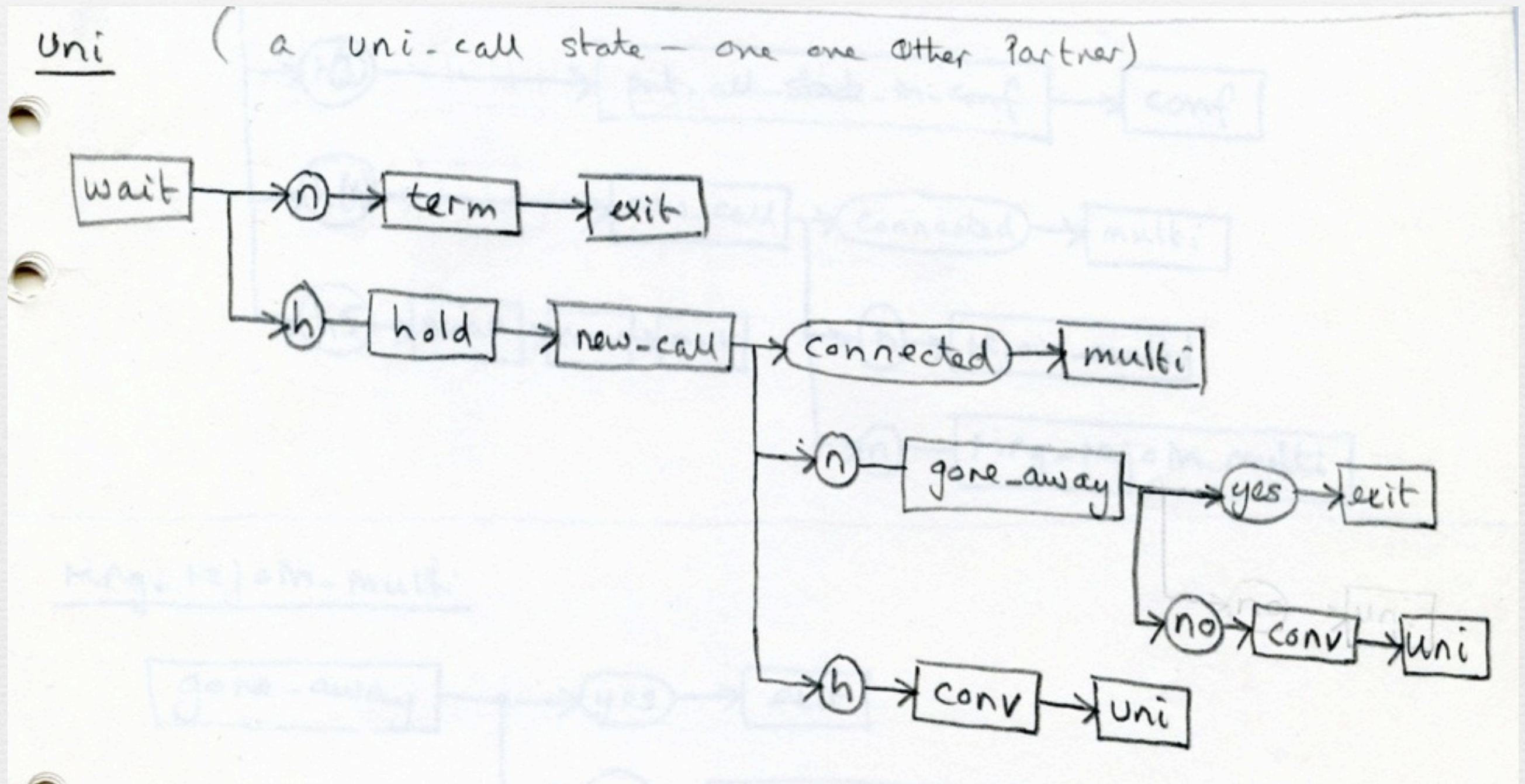


Notes

- 1) an exit from `call-hot-line` of `h` allows the user to attempt a new call

case f +

Fish bone diagrams




```
function uni returns none.
```

```
1 # uni --->
```

```
    case(wait, [
```

```
        n => [term, exit],
```

```
        h => [hold,
```

```
            case(new_call, [
```

```
                connected => multi,
```

```
                n => case(gone_away, [
```

```
                    yes => exit,
```

```
                    no => [conv, uni]
```

```
                ]),
```

```
                h => [conv, uni]
```

```
            ])
```

```
    ]
```

```
)).
```


Today

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```
uni() ->
  receive
    n ->
      term(), exit();
    h ->
      case new_call() of
        connected ->
          multi();
        n ->
          case gone_away() of
            yes -> exit();
            n -> conv(), uni()
          end;
        h ->
          conv(), uni()
      end
  end
end.
```


1988 -
Robert
Virding joins
the team

4 days for a
rewrite

Not so fast

88/12/16
12:44:20

erlang.pl

1

```

/*
 * $HOME/erlang.pro
 *
 *          Copyright (c) 1988 Ericsson Telecom
 *
 * Author: Joe Armstrong
 * Creation Date: 1988-03-24
 * Purpose:
 *   main reduction engine
 *
 * Revision History:
 *   88-03-24      Started work on multi processor version
 *                 of erlang
 *   88-03-28      First version completed (Without timeouts)
 *   88-03-29      Correct small errors
 *   88-03-29      Changed 'receive' to make it return the pair
 *                 msg(From,Mess)
 *   88-03-29      Generate error message when out of goals
 *                 i.e. program doesn't end with terminate
 *   88-03-29      added trace(on), trace(off) facilities
 *   88-03-29      Removed Var := {...} , this can be achieved
 *                 with {...}
 *   88-05-27      Changed name of file to erlang.pro
 *                 First major revision started - main changes
 *                 Complete change from process to channel
 *                 based communication
 *                 here we (virtually) throw away all the
 *                 old stuff and make a bloody great data base
 *   88-05-31      The above statements were incorrect much better
 *                 to go back to the PROPER way of doing things
 *                 long live difference lists
 *   88-06-02      Reds on run([et5]) = 245
 *                 changing the representation to separate the
 *                 environment and the process - should improve things
 *                 It did .... reds = 283 - and the program is nicer!
 *   88-06-08      All pipe stuff working (pipes.pro)
 *                 added code so that undefined functions can return
 *                 values

```


1988

- ❧ Let's make a product
 - ❧ Documentation ...
 - ❧ Community ...
 - ❧ Performance ...
 - ❧ Courses ...

Documentation

erlang vsn 1.05

h	help
⊗ reset	reset all queues
reset_erlang	kill all erlang definitions
load(F)	load erlang file <F>.erlang
load	load the same file as before
load(?)	what is the current load file
what_erlang	list all loaded erlang files
go	reduce the main queue to zero
send(A,B,C)	perform a send to the main queue
send(A,B)	perform a send to the main queue
cq	see queue - print main queue
wait_queue(N)	print wait_queue(N)
cf	see frozen - print all frozen states
eqns	see all equations
eqn(N)	see equation(N)
start(Mod,Goal)	starts Goal in Mod
top	top loop run system
q	quit top loop
open_dots(Node)	opens Node
talk(N)	N=1 verbose, =0 silent
peep(M)	set peeping point on M
no_peep(M)	unset peeping point on M
vsn(X)	erlang vsn number is X

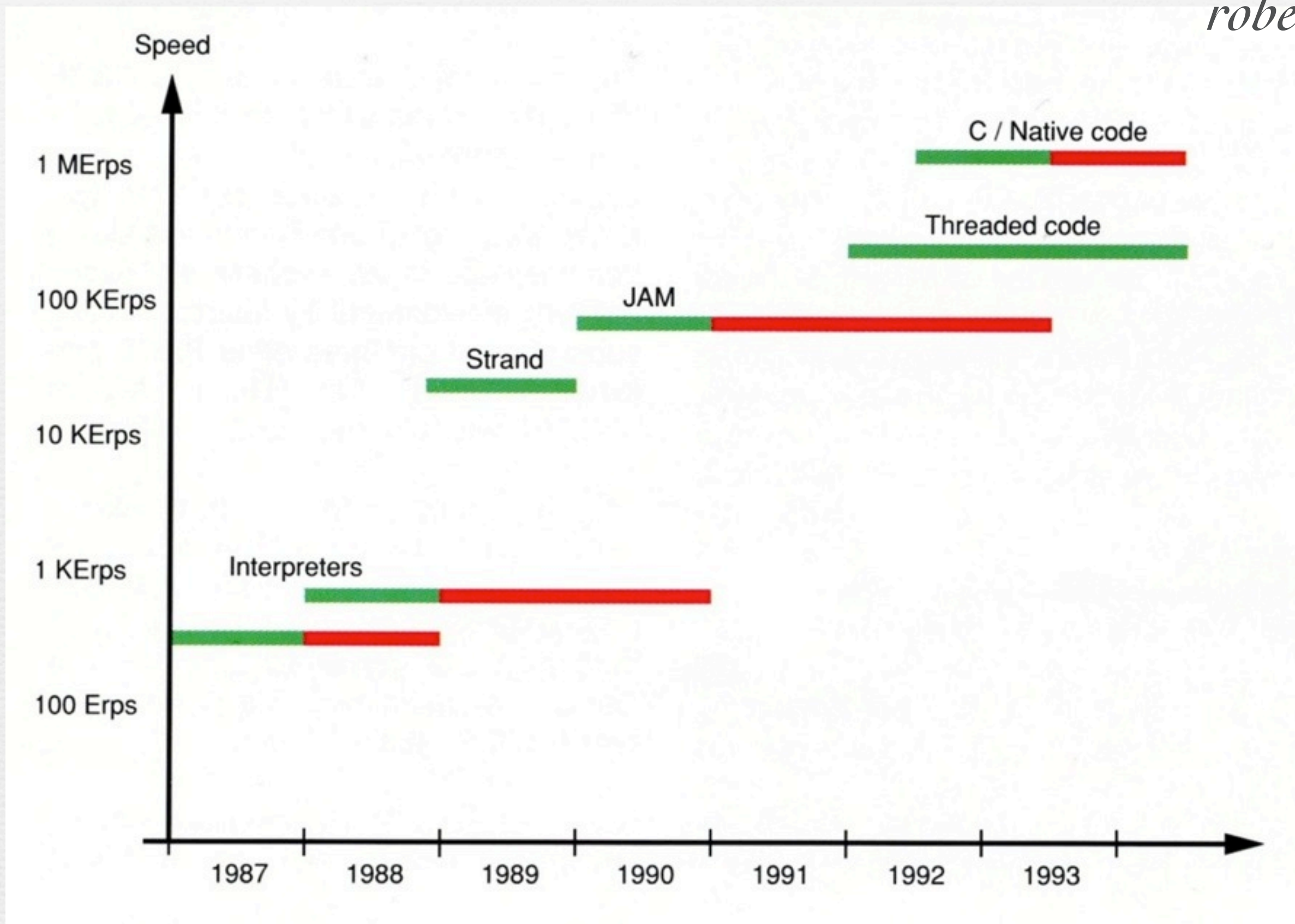


Performance

Why own VM?

- ❧ Performance
- ❧ Semantics (code change)/isolation/real-time GC
- ❧ How?
 - ❧ byte code interpreter (inspiration P-code, WAM)





Improving Performance

- ❧ Experiments
- ❧ Parlog
- ❧ Strand
- ❧ ...
- ❧ Own VM
- ❧ Compiler and Emulator in Prolog



35 ERPS

“do it in C”

Mike Williams
reads Joe's C
and declares it it be "the
worse program ever
written"

The JAM Erlang virtual machine

- JAM = Joe's abstract machine
- Joe → Compiler (Prolog → Erlang) + Architecture
- Mike → VM in “CV”
- I thought I really knew how to program “C” until I started to program the JAM
- VM
 - Byte code instructions
 - 32 bits: 8 bits tag, 24 bits data/pointer
 - Each Erlang process has it's own separate stack and heap
 - Garbage collection per Erlang process

JAM

First version on a VAX 11/750

3 MHz clock, 8 MByte memory, about 300 Mbyte disk

Second version on SUN workstation

Motorola 68K processor

Later Sparc

First use in product on “Mobility Server”

OS – VXWorks, processor 68 K

ETS (Erlang Term Storage) added later

Enabler for the Mnesia real time fault tolerant database.

Some Later products

Anx – ADSL DSLAM

AXD 301 ATM Switch

SGSN MME (Data access for GSM GPRS, WCDMA and LTE)

Other VMs

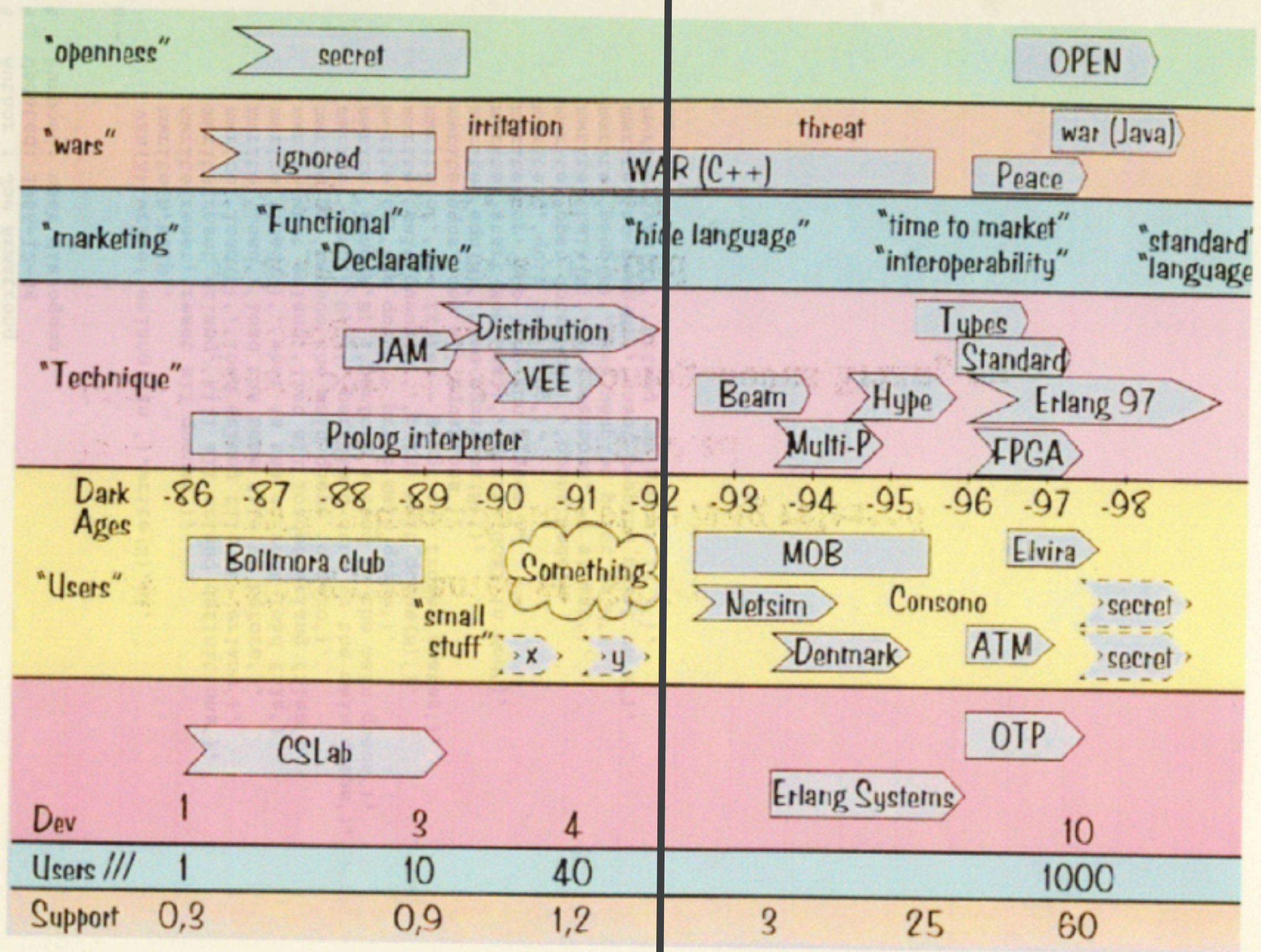
- VEE = Viriding's Erlang Engine
- BEAM = Bogdan's (nowadays Björn's) Erlang abstract Machine
 - BEAM has replaced JAM in all Ericsson products

Fault and Failure Handling

- Fault = bug in code
- Failure = hardware breaks
- Concept:
 - Faults cannot be handled in the same context (i.e. Erlang process) as they occur
 - Failures cannot be handled in the same hardware which is broken
 - Code which handles faults and failures must be as simple as possible.
- Error handling concepts inspired by the “C” wire in ancient relay based telephone exchanges
- Concept of linked process means that if one of them crashes (fault or failure) they all terminative
 - Except super simple recovery processes which receive information about the fault/failure and take remedial action.
- Often used principle:
 - Put steady state data in the Mnesia, let failing transactions crash, recovery processes use data in Mnesia to restore stable state.

Things you may not have thought about

- Dynamic typing
 - Makes tracing and debugging a lot easier as lot of symbolic information is retained
 - Makes marshalling of data for inter-machine communication easy at runtime
- Being able to change code “on the fly” at runtime greatly speeds up the
code->test->debug->correct
cycle
- Distribution is transparent in nearly all the code!
- Selective message reception greatly simplifies state machine code
- You can implement synchronous interprocess communication on top of asynchronous communication, the the inverse is very much harder!



Rapid Prototyping c.1992

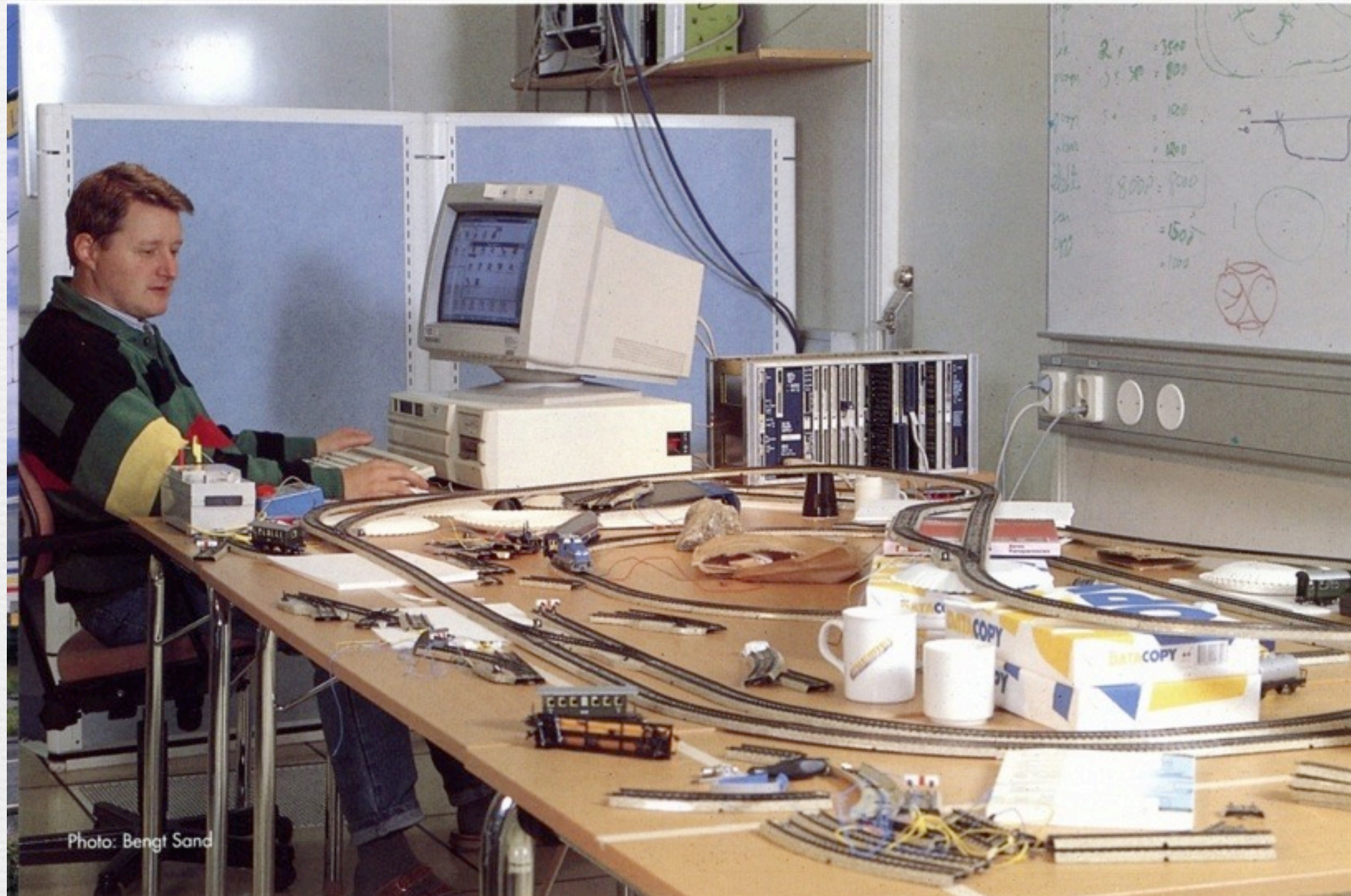
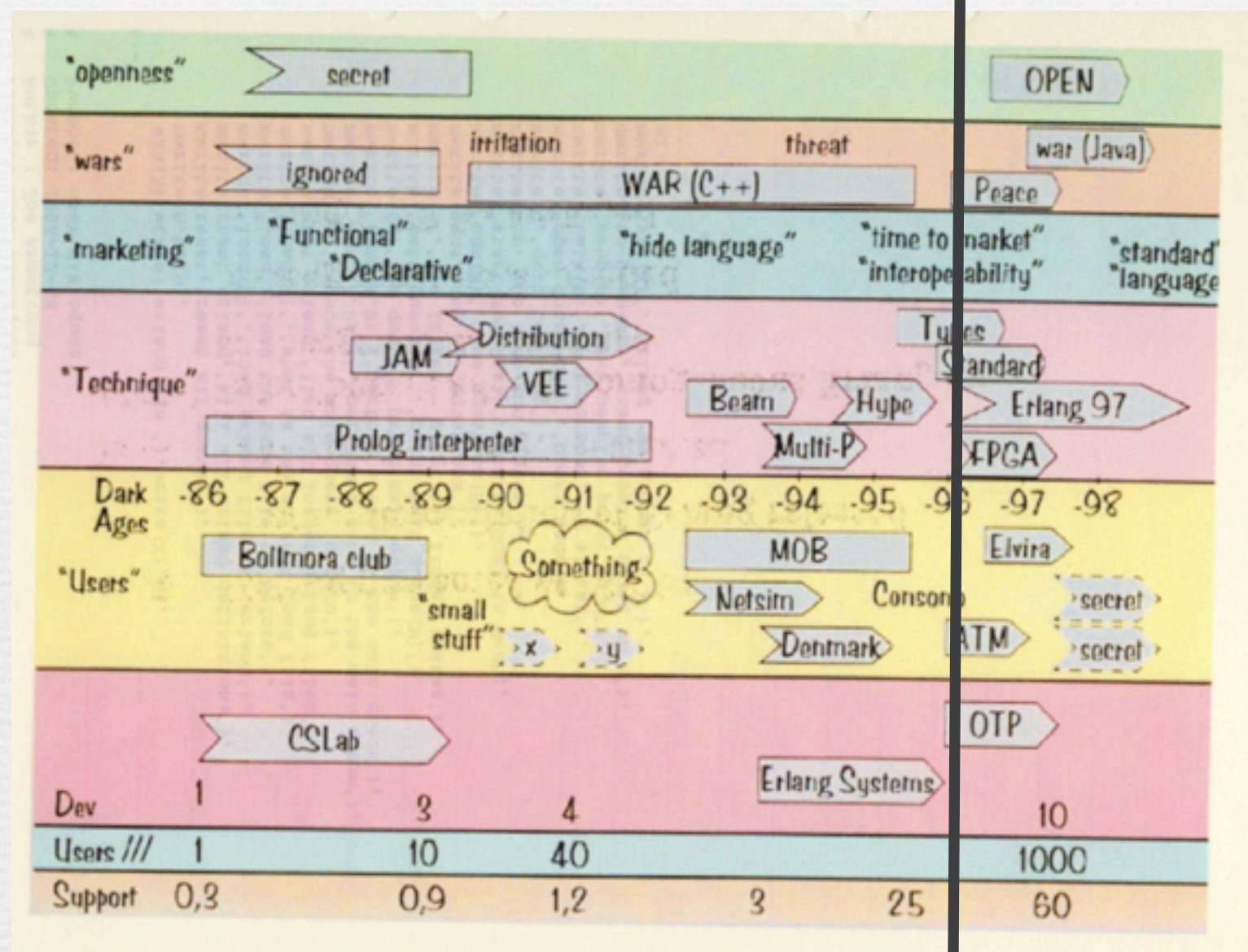


Photo: Bengt Sand

1992 - 1995
nothing much
happens ...

robert

8 Dec 1995
AXE-N Cancelled ...



1996 AXD 301 starts

Lot's of stuff happens quickly

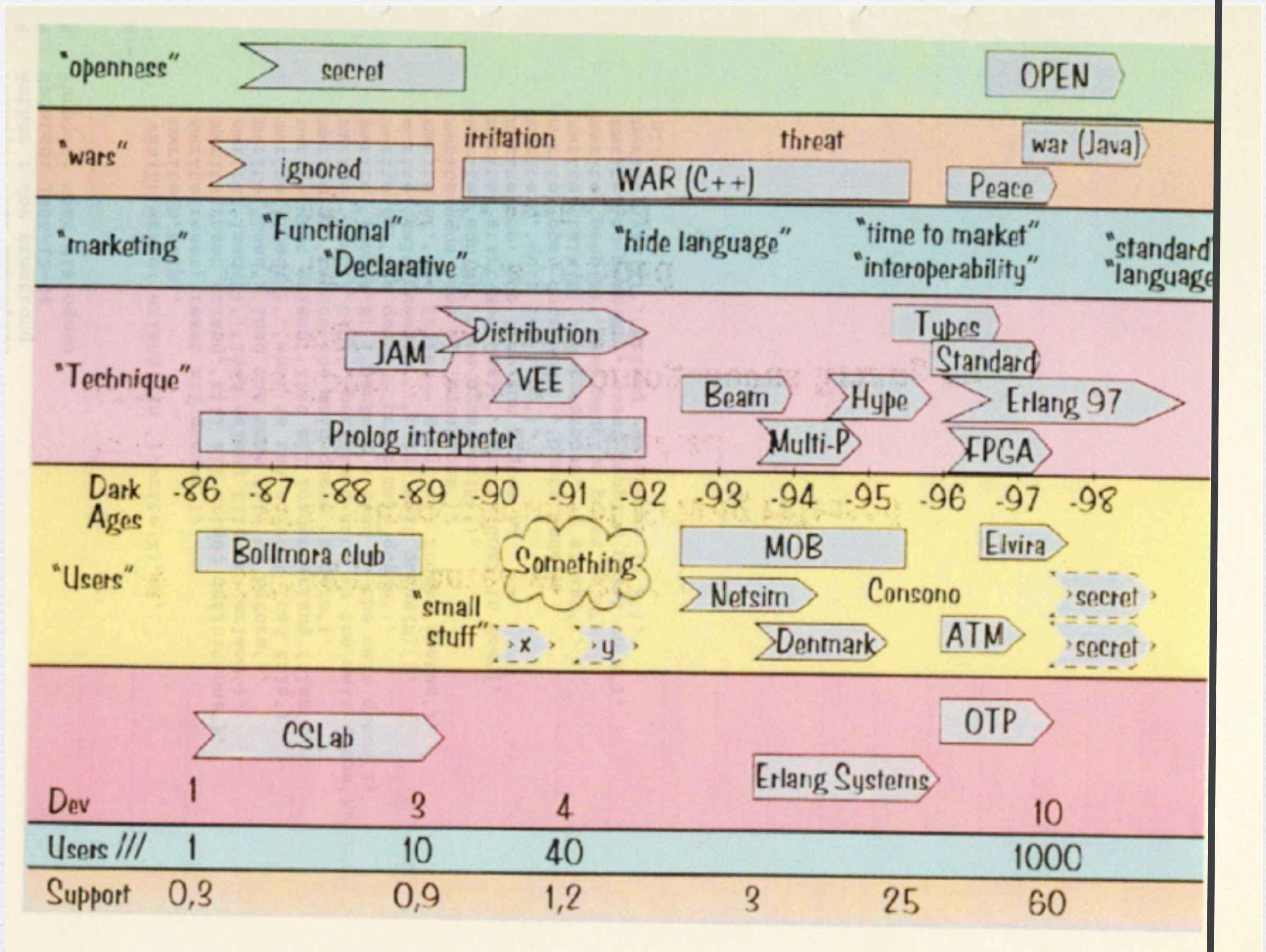
robert

1996 - 1998
nothing much
happens ...

robert

1998
AXD 301 is a great
success ...

Still in use today in BT network



BANNED

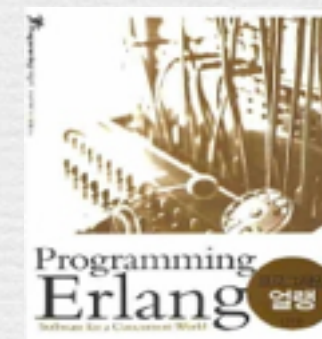
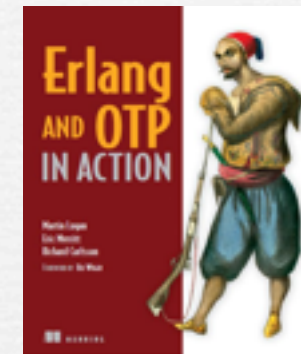


1998 Stuff Happens

- Lot's of stuff happens quickly
- Erlang becomes Open Source ...
- Four days later ... Bluetail AB ...

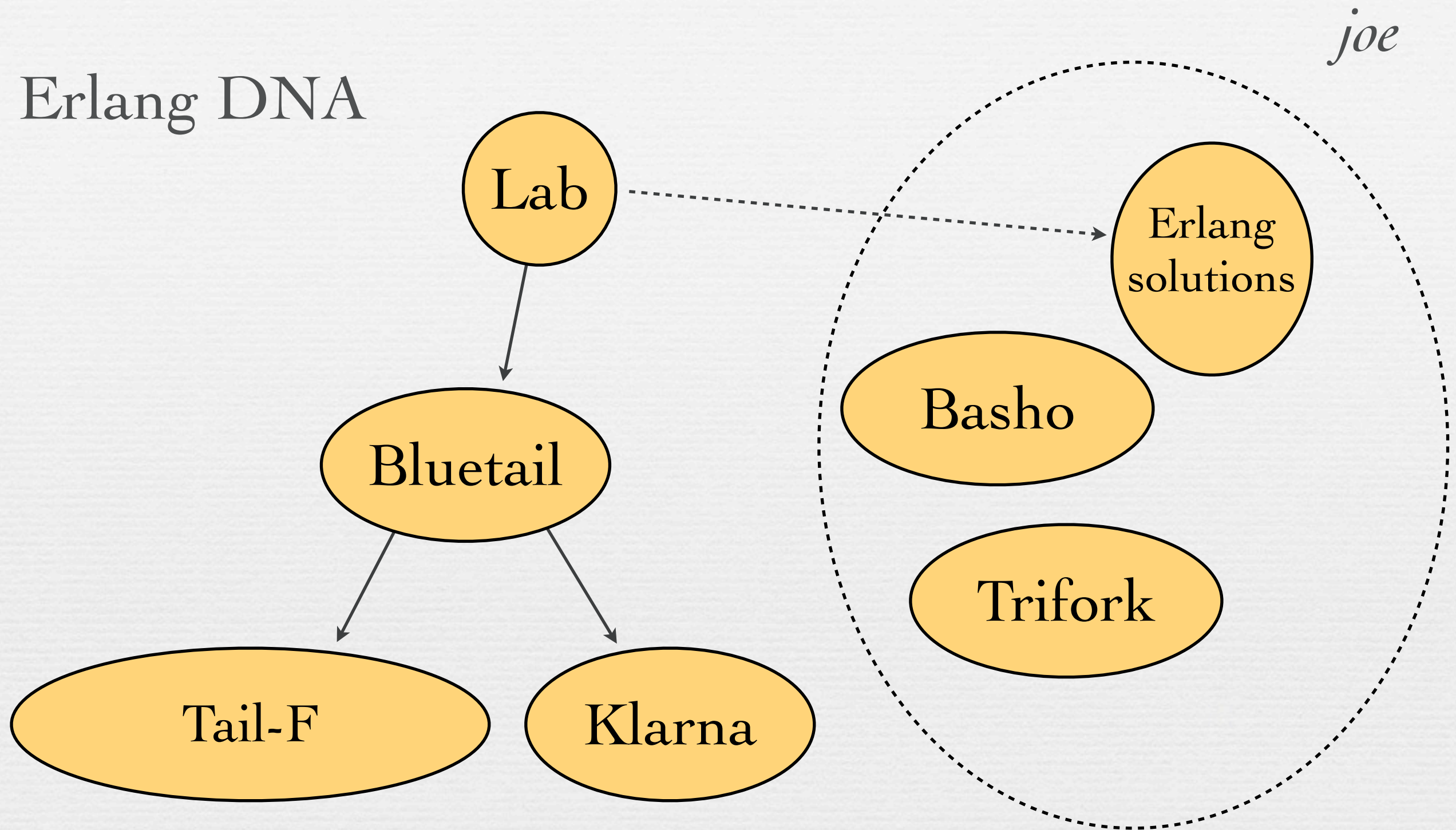
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10 years later...





Erlang DNA



And loads more that we don't know about ... (ask Francesco)

The Future

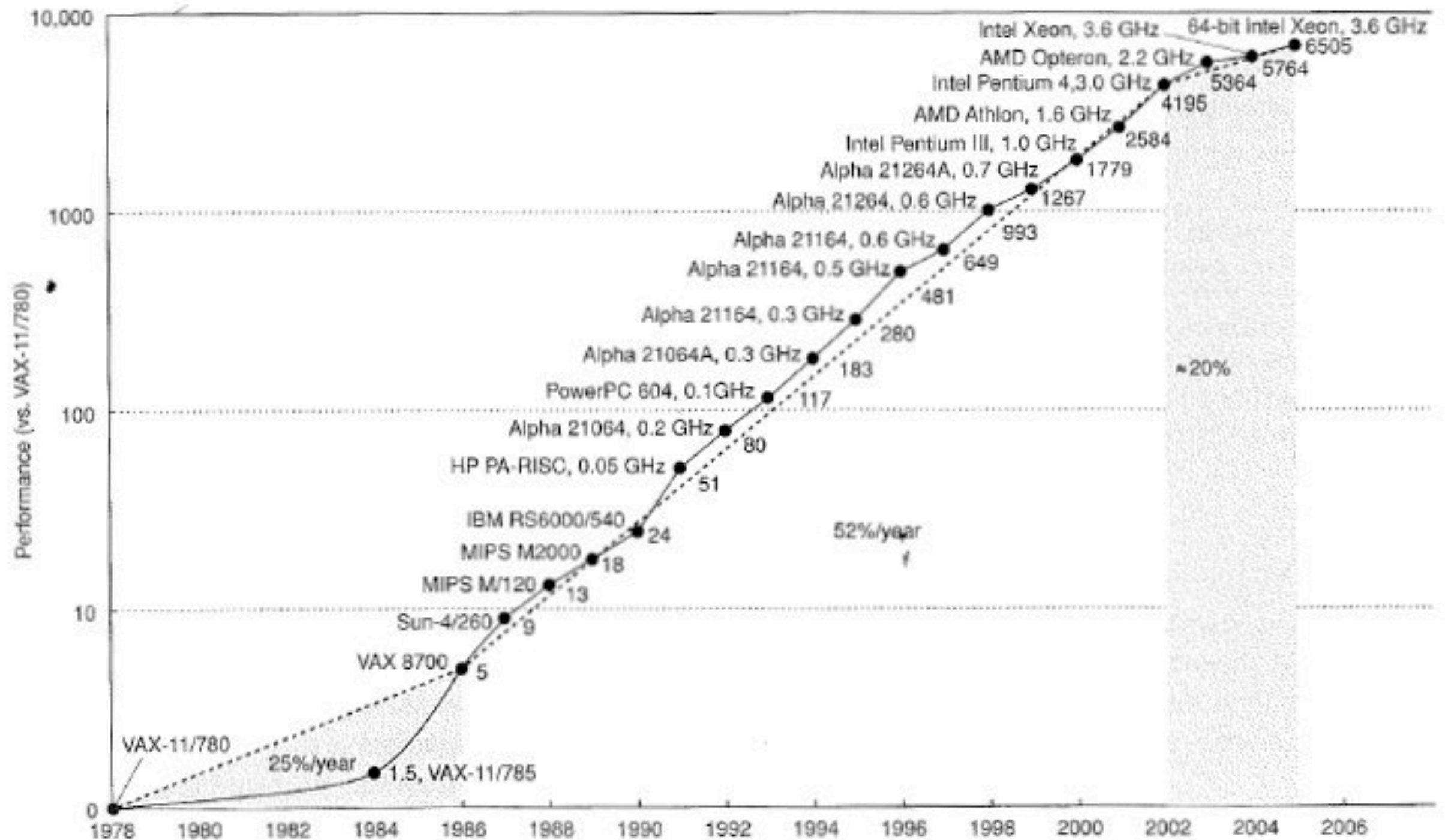
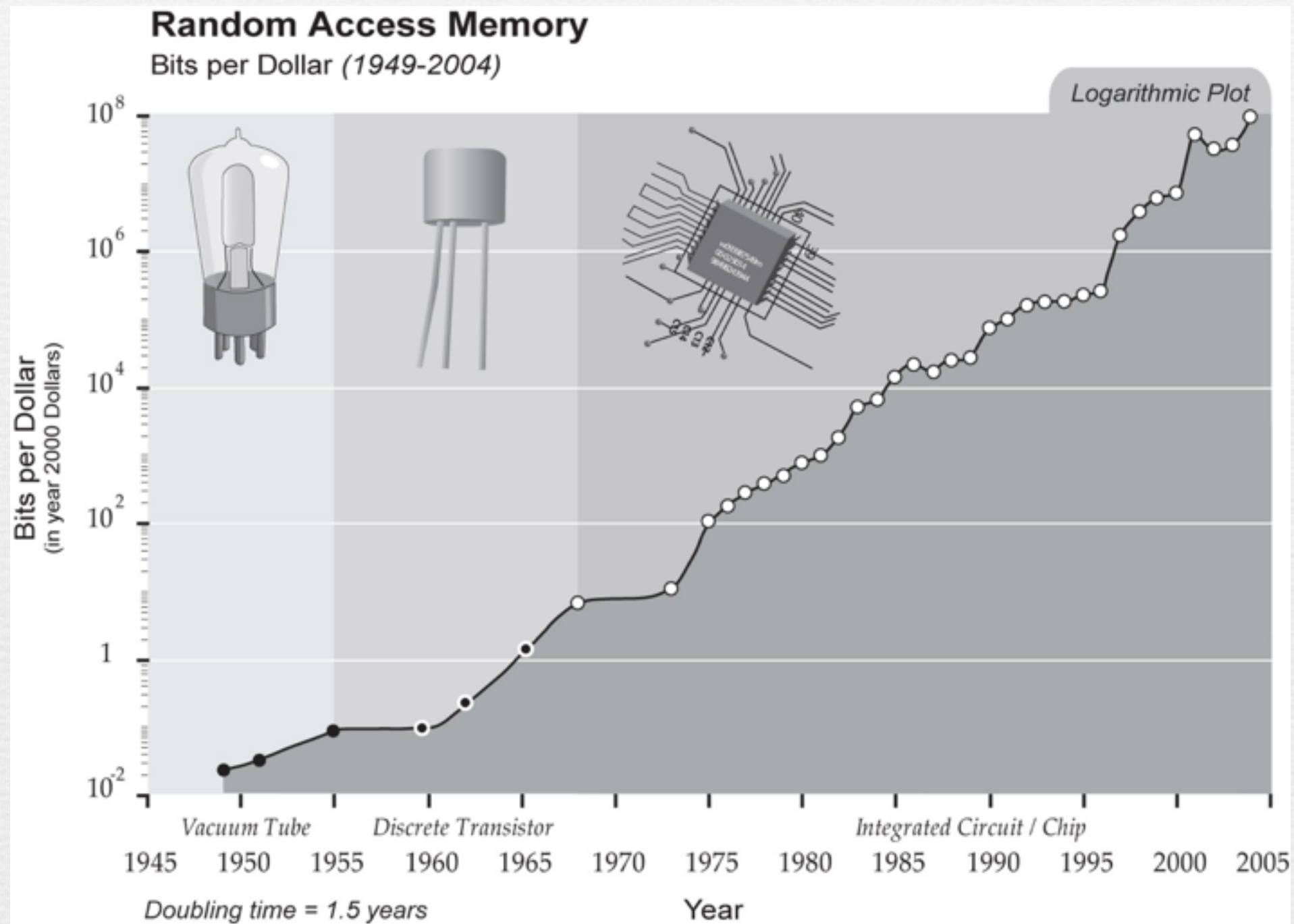


Figure 1.1 Growth in processor performance since the mid-1980s.



source: www.singularity.com

What happened?

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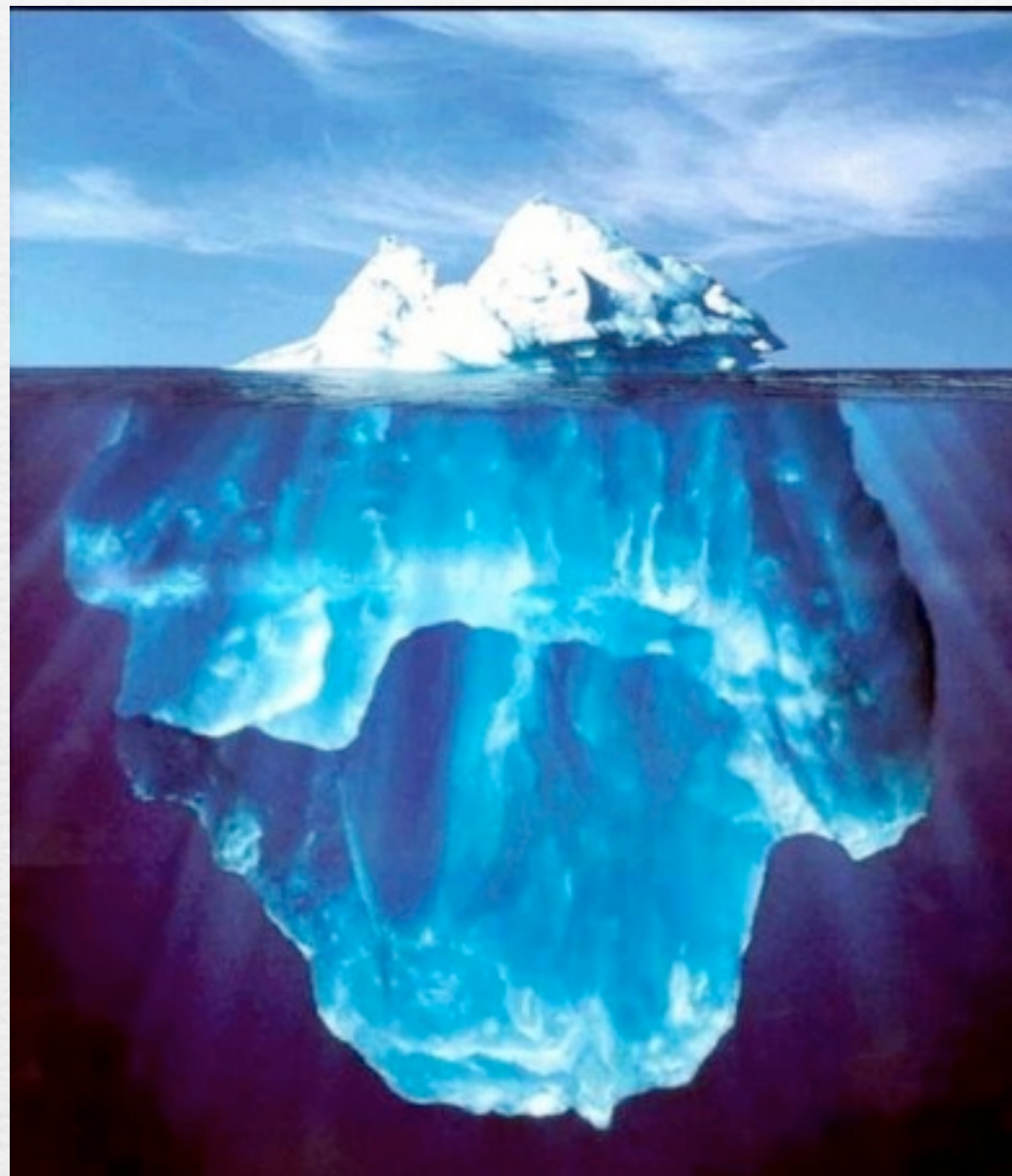
Software

- ▶ 1879 - Frege
- ▶ 1930 - Curry
- ▶ 1958 - LISP
- ▶ 1969 - agents/actors/smalltalk
- ▶ 1972 - Prolog
- ▶ 1978 - CSP
- ▶ 1983 - Occam (+hardware)
- ▶ 1986 - Parlog/Strand
- ▶ 1986 - Erlang
- ▶ 2011 - Elixir

A heck of a lot of
hardware stuff
has happened in
the last 10 years

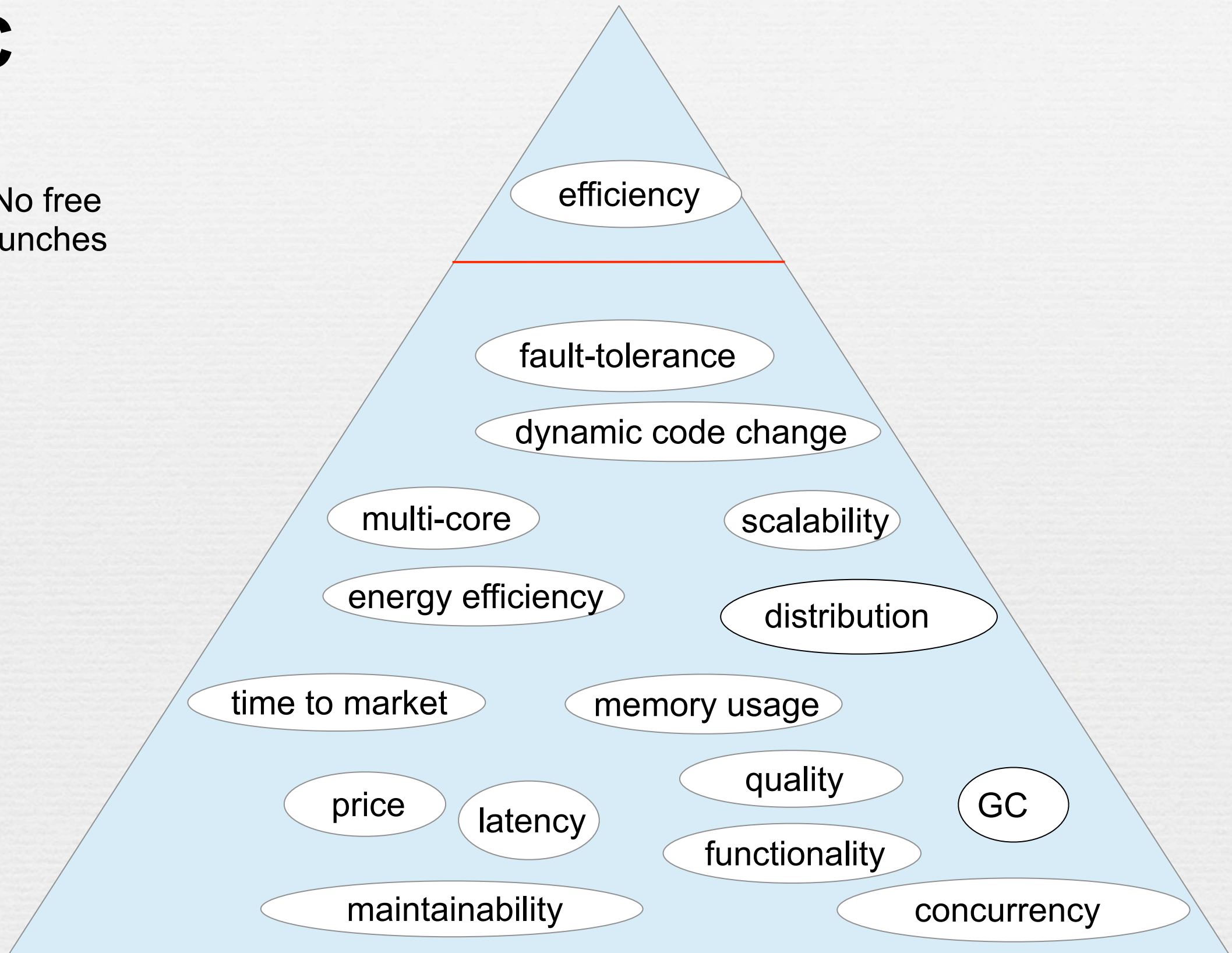
Hardware - sort

- ▶ 1879 - 0 Hz - 0 MB
- ▶ 1958 - ...
- ▶ 1980's TCP/Internet
- ▶ 1986 - 1.6 Mhz
- ▶ 2000 - 1GHz clocks
- ▶ 2000 - Always connected
- ▶ 2000 - Mobile revolution
- ▶ 2004 - GPRS/3G/WCDMA
- ▶ 2004 - multi cores /GB Ram/
- ▶ 2010 - 4G (LTE)/TB disk
- ▶ 2020 - Peta bytes?/K Cores

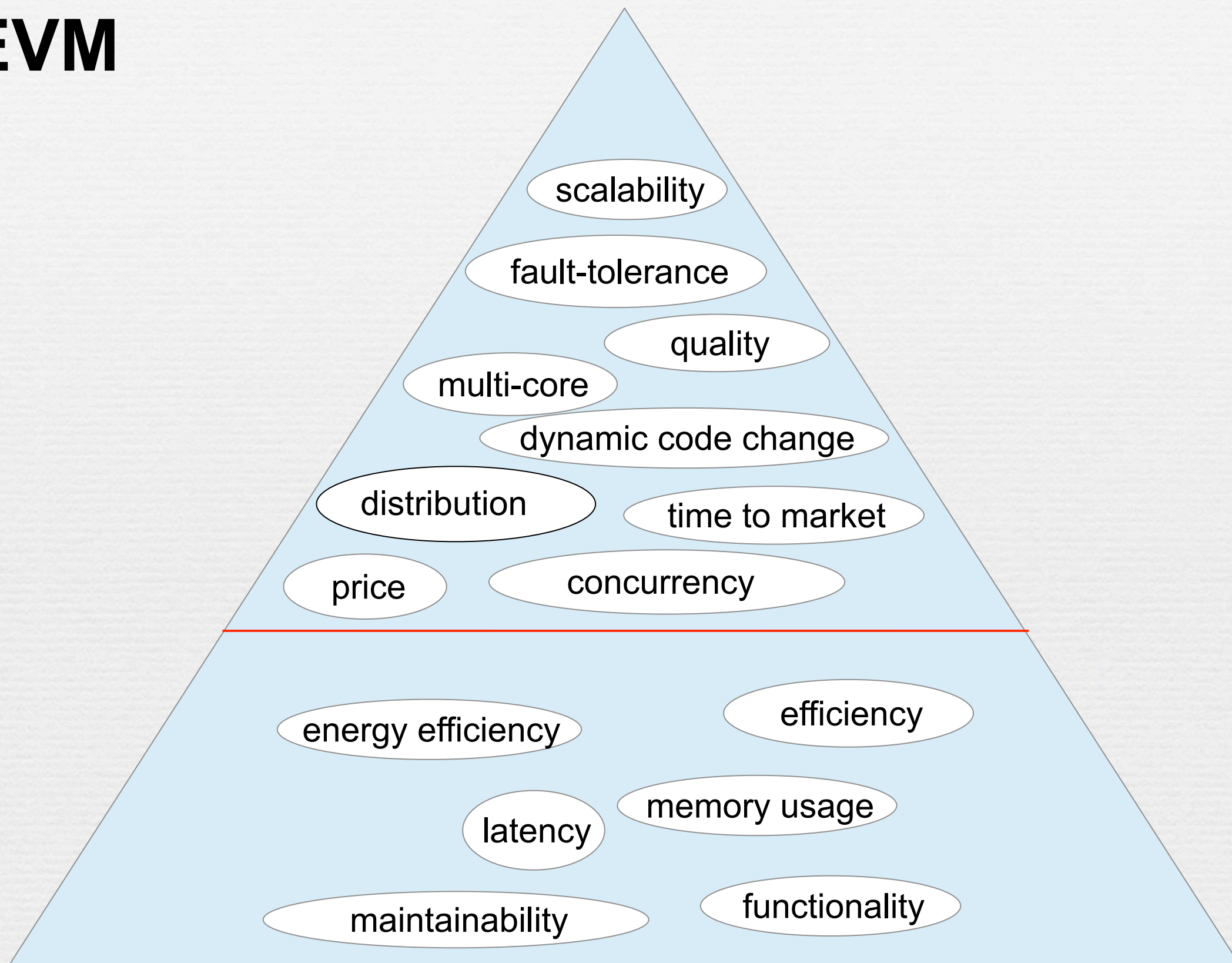


C

No free
lunches



EVM

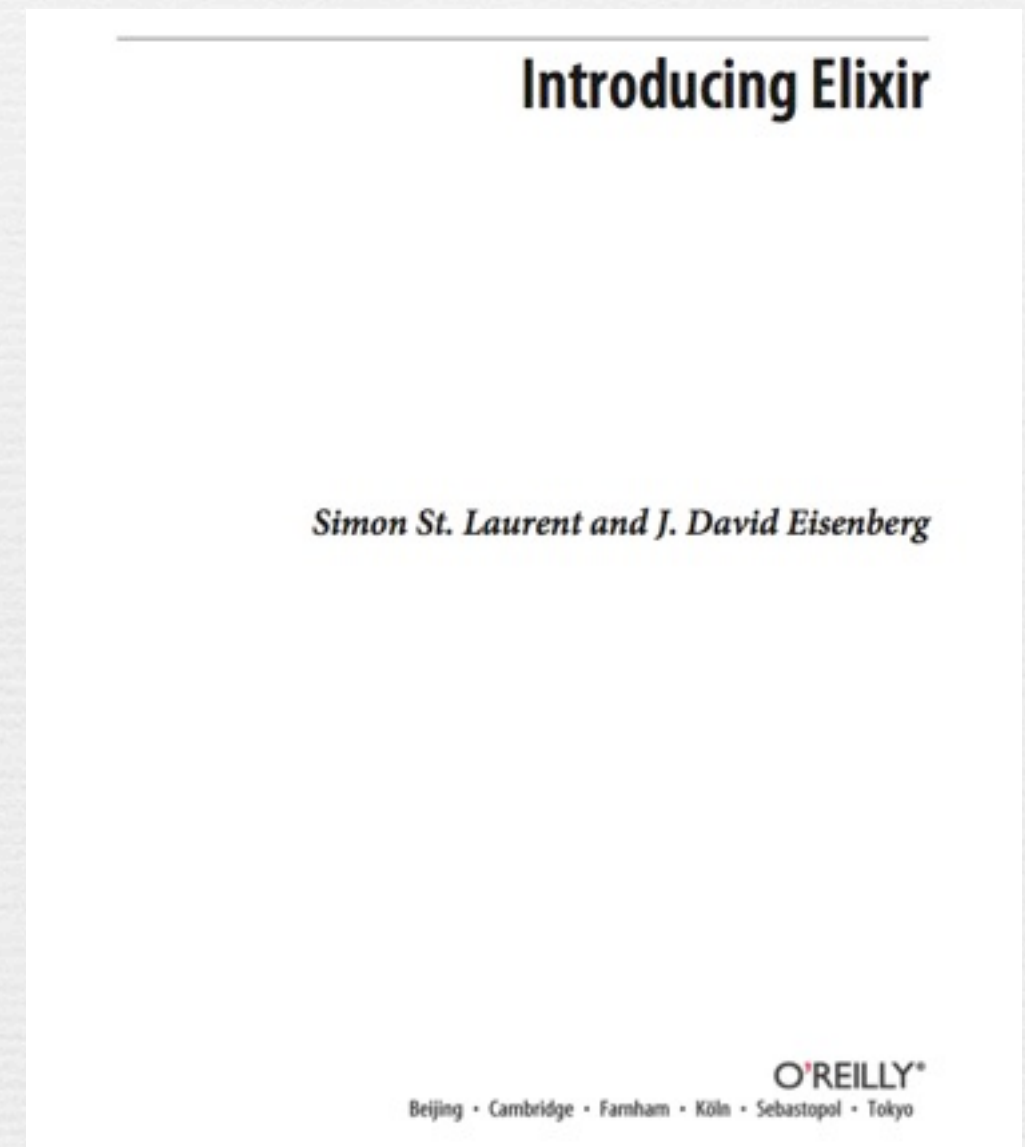


Languages on the EVM



Erlang
Prolog
LFE
LUA
Elixir
Joxa (lisp)
Reia

Elixir



Erlang

- 🌀 Language discussion limited by geography (pre WWW)
- 🌀 Closed source
- 🌀 “Funny syntax”
- 🌀 Started 1986
- 🌀 Book1 1993
- 🌀 Book2 2007 (21 years later)
- 🌀 Not marketed

Elixir

- 🌀 Language discussion on WWW
- 🌀 Open source
- 🌀 “Ruby syntax”
- 🌀 Started 2011
- 🌀 Books1+2 2013 (2 years later)
- 🌀 Marketed

Tomorrow

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- ❧ 2020 - 1 Million cores
- ❧ 2020 - 1 TB flash in mobile / 1 GB/sec mobile
PByte disks
- ❧ 2020 - 100 B connected devices / ubiquitous
networking

Today

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- ❧ MME (=Mobile Management Entity (LTE/4G))
- ❧ SGSN (=Serving GPRS Support Node (GPRS = General Packet Radio Service, 3G))
- ❧ WhatsApp



How will we program
all this new stuff?

Which X on the
EVM?