



FUTURE EXTENSIONS TO THE NATIVE INTERFACE

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EXISTING NATIVE INTERFACE

> Drivers

- Provides native code via a set of callbacks
- Executed in context of a port
- Ports similar to processes

> NIF

- Native code
- Executed in context of a process

COMMUNICATING WITH PORTS

Message To Port	Message From Port
<code>{Pid, {command, Data}}</code>	<code>{Port, {data, Data}}</code>
<code>{Pid, close}</code>	<code>{Port, closed}</code>
<code>{Pid, {connect, NewPid}}</code>	<code>{Port, connected}</code>

Port BIFs
<code>port_command(Port, Data)</code>
<code>port_close(Port)</code>
<code>port_connect(Port, NewPid)</code>
<code>port_control(Port, Operation, Data)</code>
<code>erlang:port_call(Port, Operation, Data)</code>
<code>port_info(Port) / port_info(Port, Item)</code>

DRIVER CALLBACKS

```
typedef struct erl_drv_entry {
    int (*init)(void);
    ErlDrvData (*start)(ErlDrvPort port, char *command);
    void (*stop)(ErlDrvData drv_data);
    void (*output)(ErlDrvData drv_data, char *buf, int len);
    void (*ready_input)(ErlDrvData drv_data, ErlDrvEvent event);
    void (*ready_output)(ErlDrvData drv_data, ErlDrvEvent event);
    char *driver_name;
    void (*finish)(void);
    void *handle;
    int (*control)(ErlDrvData drv_data, unsigned int command, char *buf, int len, char **rbuf, int rlen);
    void (*timeout)(ErlDrvData drv_data);
    void (*outputv)(ErlDrvData drv_data, ErlIOVec *ev);
    void (*ready_async)(ErlDrvData drv_data, ErlDrvThreadData thread_data);
    void (*flush)(ErlDrvData drv_data);
    int (*call)(ErlDrvData drv_data, unsigned int command, char *buf, int len, char **rbuf, int rlen, unsigned int *flags);
    void (*event)(ErlDrvData drv_data, ErlDrvEvent event, ErlDrvEventData event_data);
    int extended_marker;
    int major_version;
    int minor_version;
    int driver_flags;
    void *handle2;
    void (*process_exit)(ErlDrvData drv_data, ErlDrvMonitor *monitor);
    void (*stop_select)(ErlDrvEvent event, void* reserved);
} ErlDrvEntry;
```

DRIVER CALLBACKS

Driver Callback	Operation	Call Reason
<code>output()</code> <code>outputv()</code>	<code>Port ! {Pid, {command, Data}}</code> <code>port_command()</code>	Process communication
<code>control()</code>	<code>port_control()</code>	
<code>call()</code>	<code>erlang:port_call()</code>	
<code>timeout()</code>	<code>driver_set_timer()</code>	Event subscription
<code>process_exit()</code>	<code>driver_monitor_process()</code>	
<code>ready_async()</code>	<code>driver_async()</code>	
<code>ready_input()</code> <code>ready_output()</code> <code>event()</code> <code>stop_select()</code>	<code>driver_select()</code>	

ISSUES WITH PORTS

- › Cannot handle arbitrary protocols
 - Only understand a fixed set of messages
 - › {Pid, {command, Data}}
 - For example, cannot handle the I/O protocol
 - Need to be paired with a process
 - › may imply unnecessary scheduling
- › Lacks functionality compared to processes
 - Cannot be monitored
 - No remote operations
- › Inefficient compared to NIFs
 - Argument processing
 - Heap management
- › Own implementation
 - Doesn't benefit from optimizations nor new functionality made for processes

NATIVE IMPLEMENTED FUNCTIONS

- › Easy to implement
- › Efficient
 - Direct call into native code
 - Reads and writes directly from/to process heap
- › Can ***not*** replace ports
 - Only code
 - No fixed execution context
 - Cannot subscribe for events

```
ERL_NIF_TERM hello_nif_world(ErlNifEnv *env, int argc, const ERL_NIF_TERM argv[])  
{  
    return enif_make_string(env, "Hello NIF world!", ERL_NIF_LATIN1);  
}
```

NATIVE PROCESSES

- › Indistinguishable from ordinary processes
- › Minimal amount of native code
 - Execute both Erlang code and native code
- › Share implementation with ordinary processes
- › As efficient as NIFs
- › Share implementation with NIFs

NATIVE PROCESS CREATION

- › `spawn()`
- › Convert to native process in NIF
 - Call `enif_become_native_process()`
 - Process will enter a native event loop and will not return from NIF
 - Code is part of a NIF library

```
ERL_NIF_TERM enif_become_native_process(ErlNifEnv *env, ErlNProcSetup *setup);

typedef struct {
    ...
    void *state;
    void (*io_event)(ErlNifEnv *,
                    void *,
                    ErlNProcIoEventType,
                    ErlNProcIoDescr);
    void (*event)(ErlNifEnv *,
                 void *,
                 ErlNProcEventType,
                 ErlNifEnv *,
                 int,
                 ERL_NIF_TERM *);
} ErlNProcSetup;
```


IO-EVENT CALLBACK

```
void io_event_callback(ErlNifEnv *proc_env, ←  
                      void *state, ←  
                      ErlNProcIoEventType type, ←  
                      ErlNProcIoDescr io_descr); ←
```

```
typedef enum {  
    ERL_NPROC_IO_EVENT_INPUT_READY = 0,  
    ERL_NPROC_IO_EVENT_OUTPUT_READY = 1,  
    ERL_NPROC_IO_EVENT_STOP_SELECT = 2  
} ErlNativeProcIoEventType;
```

```
void enif_nproc_select(ErlNifEnv *env, ErlNProcIoDescr io_descr, int mode, int on);  
  
#define ERL_NPROC_IO_EVENT_READ 1  
#define ERL_NPROC_IO_EVENT_WRITE 2  
#define ERL_NPROC_IO_EVENT_USE 4
```

EXECUTING ERLANG CODE

```
void enif_nproc_schedule_ecall(ErlNifEnv *proc_env,
                              ERL_NIF_TERM module,
                              ERL_NIF_TERM function,
                              ERL_NIF_TERM arguments,
                              void (*result_func)(ErlNifEnv *proc_env,
                                                    void *state,
                                                    ErlNProcResultType result_type,
                                                    ERL_NIF_TERM result));
```

```
typedef enum {
    ERL_NPROC_RESULT_RETURN = 0,
    ERL_NPROC_RESULT_THROW = 1,
    ERL_NPROC_RESULT_ERROR = 2,
    ERL_NPROC_RESULT_EXIT = 3
} ErlNProcResultType;
```

- › May only be called by callbacks
- › Erlang code may call NIFs
- › Continuations of a callback call
- › No new callback calls until done

NATIVE PROCESS CODE CHANGE

- › Load new module
- › Call NIF in new module using `enif_nproc_schedule_ecall()`
- › Call `enif_native_process_code_change()` from NIF

```
ERL_NIF_TERM enif_native_process_code_change(ErlNifEnv *env, ErlNProcSetup *setup);

typedef struct {
    ...
    void *state;
    void (*io_event)(ErlNifEnv *,
                    void *,
                    ErlNProcIoEventType,
                    ErlNProcIoDescr);
    void (*event)(ErlNifEnv *,
                 void *,
                 ErlNProcEventType,
                 ErlNifEnv *,
                 int,
                 ERL_NIF_TERM *);
} ErlNProcSetup;
```

```
void *enif_get_nproc_state(ErlNifEnv *env);
```

REVERTING BACK TO AN ORDINARY PROCESS

```
void enif_nproc_become_ordinary_process(ErlNifEnv *env);
```

- › Call `enif_nproc_become_ordinary_process()`
- › May only be called in callback
- › Exits native event loop and return to Erlang code from NIF that converted to native process

REASONS TO WRITE NATIVE CODE

- › Optimizing
- › Access functionality not available in Erlang
 - Operating System services
 - › Sockets
 - › Files
 - › ...
 - Third party libraries
 - › ZLib
 - › OpenSSL
 - › wxWidgets
 - › ...

PROBLEMATIC NATIVE CODE

- › Unstable code
 - Runtime system crash
 - Internal runtime system inconsistency
- › Code that run for long periods of time
 - Destroy responsiveness of runtime system (also in the SMP runtime system)

UNSTABLE NATIVE CODE

- › Fix if possible
- › Runtime system executing native code is vulnerable
- › Isolate by running in separate runtime system
 - Unstable third party libraries
 - Debugging
- › Distribution transparency makes this easier
 - Ports aren't distribution transparent
 - Native processes are

LONG EXECUTION OF NATIVE CODE

- › Make preemptable by dividing a large workload into a number of smaller workloads
 - Preferred but not always possible
 - › Third party libraries
 - › Blocking wait
- › Threading
 - Async thread pool (drivers only)
 - Own driver/NIF threads

PREEMPTABLE NIF AND DRIVER

```

large_workload(Arg) ->
  case large_workload_nif(Arg) of
    {continue, Reductions, NewState} ->
      erlang:bump_reductions(Reductions),
      large_workload(NewState);
    {done, Reductions, Result} ->
      erlang:bump_reductions(Reductions),
      Result
  end.

```

```

static void large_workload_output_callback(ErlDrvData drv_data,
                                           char *buf, int len)
{
    int not_done;
    MyDrvData *data = (ErlDrvData *) drv_data;
    /* Do work... */
    if (not_done) driver_set_timer(data->port, 0);
}

static void timeout_callback(ErlDrvData drv_data)
{
    int not_done;
    MyDrvData *data = (ErlDrvData *) drv_data;
    /* Do more work... */
    if (not_done) driver_set_timer(data->port, 0);
}

```

PREEMPTABLE NATIVE PROCESS CALLBACK

```

void large_workload_callback(ErlNifEnv *proc_env, void *state, ErlNProcEventType type,
                             ErlNifEnv *arg_env, int argc, ERL_NIF_TERM *argv)
{
    ERL_NIF_TERM erl_state;
    /* Parse args, etc... */
    do_work(proc_env, state,
            ERL_NPROC_RESULT_RETURN, erl_state);
}

static void do_work(ErlNifEnv *env,
                    void *state,
                    ErlNProcResultType rtype,
                    ERL_NIF_TERM erl_state)
{
    int not_done, reductions;
    ERL_NIF_TERM new_erl_state;
    /* Do work... */
    if (not_done)
        enif_nproc_schedule_ecall(enif_make_atom(env, "my_module"),
                                enif_make_atom(env, "reschedule_point"),
                                enif_make_list(env, 2,
                                                enif_make_int(env, reductions),
                                                new_erl_state),
                                do_work);
}

```

```

-module(my_module).
-export([reschedule_point/2]).

reschedule_point(Reds, State) ->
    erlang:bump_reductions(Reds),
    State.

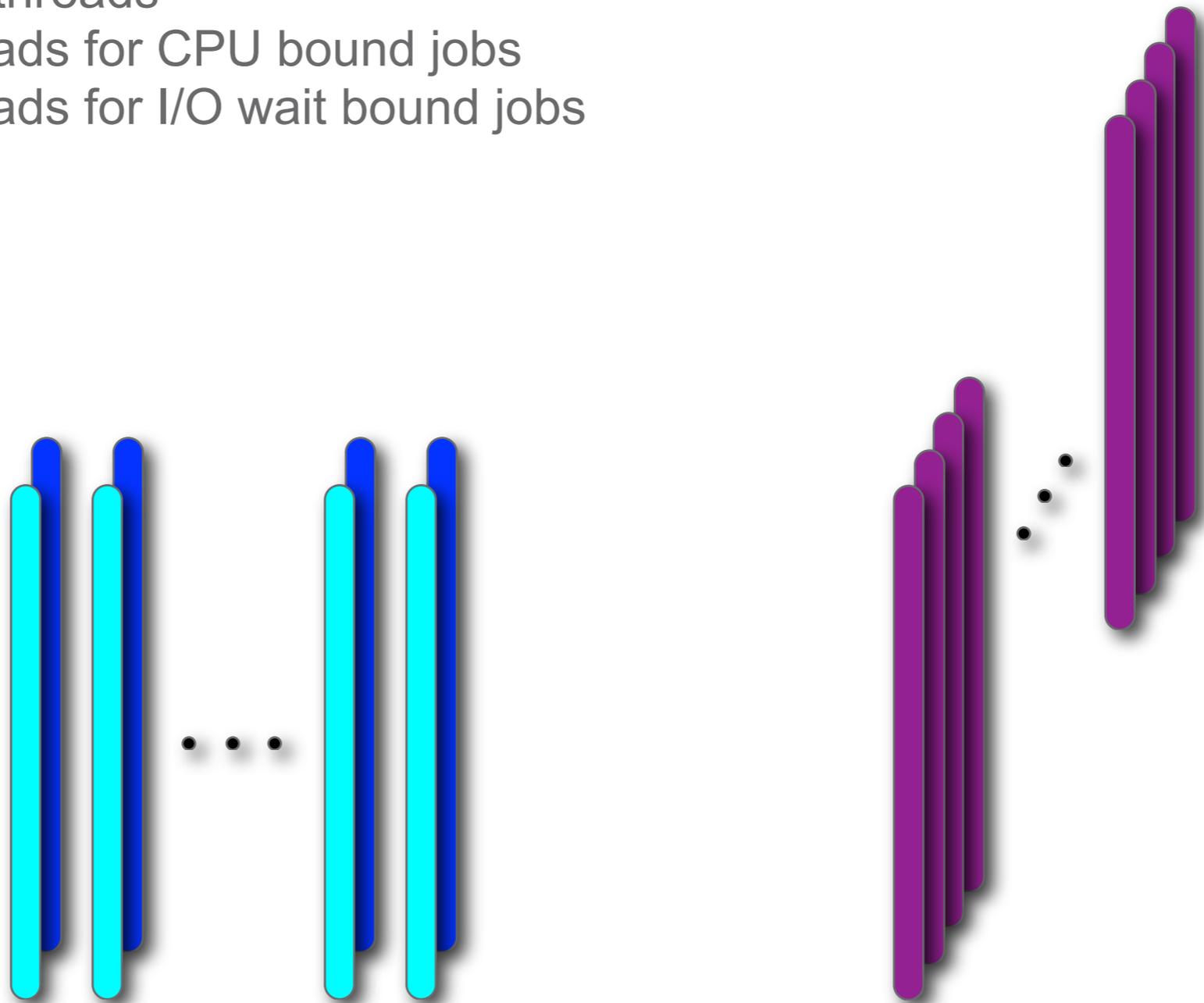
```

EASE MANAGEMENT OF LONG JOBS

- › Ease threading
- › Two sets of long jobs
 - CPU bound jobs
 - I/O wait bound jobs
- › When executed on other threads than schedulers threads
 - CPU bound have large impact on priority of other processes
 - I/O wait bound have a small or no impact on other processes
- › Two set of thread pools

DIRTY SCHEDULERS

- Ordinary scheduler threads
- Dirty scheduler threads for CPU bound jobs
- Dirty scheduler threads for I/O wait bound jobs



DIRTY JOBS

› Thread safety

- SMP runtime system; as on any scheduler thread
- Non-SMP runtime system
 - › Limited set of functions can be used
 - › Behavior slightly different

› Scheduling

- Borrowing reductions for CPU bound jobs
 - › May be automatically calculated on some platforms
 - › Native code need to bump reductions
 - › Punishment

› ERTS internal functionality

USING DIRTY SCHEDULERS

```

#define ERTS_DIRTY_JOB_FLG_CPU_BOUND 1
#define ERTS_DIRTY_JOB_FLG_IO_WAIT_BOUND 2

/* Dirty NIFs and Native Process Callbacks */
int enif_is_dirty_job(ErlNifEnv *env);
int enif_is_on_dirty_scheduler(ErlNifEnv *env);
int enif_have_dirty_schedulers(void);

ERL_NIF_TERM enif_schedule_dirty_nif(ErlNifEnv *env, int flags,
                                     ERL_NIF_TERM (*dirty_nif)(ErlNifEnv *, int
                                                                ERL_NIF_TERM *));
ERL_NIF_TERM enif_schedule_dirty_nif_finalizer(ErlNifEnv *env,
                                               ERL_NIF_TERM (*dirty_nif_finalizer)(
                                                                ErlNifEnv *,
                                                                ERL_NIF_TERM));

void enif_nproc_schedule_dirty_event(...);
void enif_schedule_dirty_event_finalizer(...);

void enif_nproc_schedule_dirty_io_event(...);
void enif_schedule_dirty_io_event_finalizer(...);

/* Dirty Driver Callbacks */

...

```

EXAMPLE OF DIRTY SCHEDULER USE

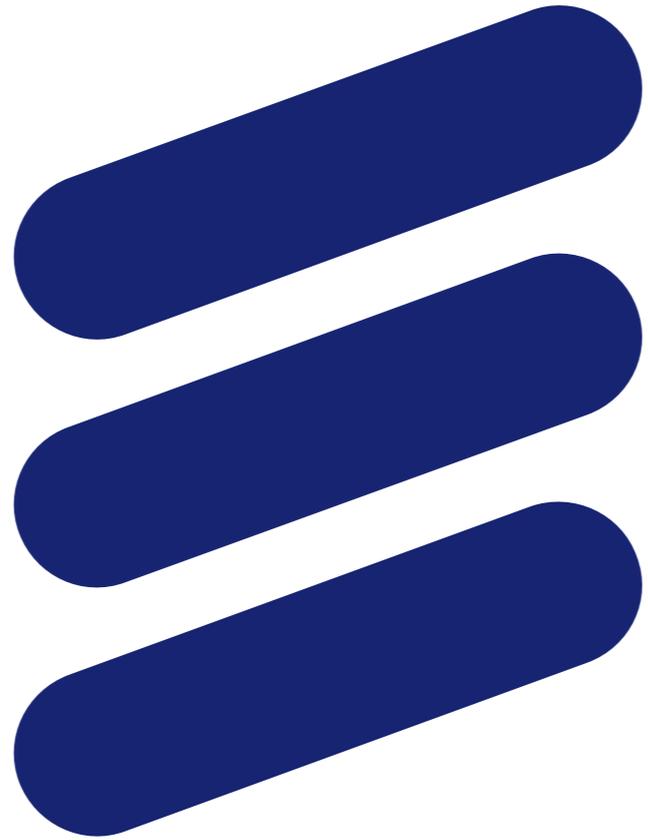
```
static ERL_NIF_TERM perhaps_long_job(ErlNifEnv *env, int argc, const ERL_NIF_TERM argv[])
{
→ if (will_be_short_job(env, argc, argv))
    return do_job(env, argc, argv);
    else
→ return enif_schedule_dirty_nif(env, ENIF_DIRTY_JOB_CPU_BOUND, do_job);
}

static ERL_NIF_TERM do_job(ErlNifEnv *env, int argc, const ERL_NIF_TERM argv[])
{
    ERL_NIF_TERM result;
    /* Do work... */
→ return enif_schedule_dirty_nif_finalizer(env, result, do_job_finalizer);
}

static ERL_NIF_TERM do_job_finalizer(ErlNifEnv *env, ERL_NIF_TERM result)
{
    ERL_NIF_TERM modified_result;
    /* Finalize... */
    return modified_result;
}
```

SUMMARY

- › Native processes make drivers obsolete
 - Will be able to handle any protocol
 - Possible to minimize the need for native code
 - Will have access to the same functionality as ordinary processes
 - More efficient
 - Distribution transparent
 - Share implementations
- › Dirty schedulers
 - Now easy to handle any code that misbehaves by monopolizing a scheduler thread
 - Ordinary schedulers should never be harassed by misbehaving jobs
- › Designed by the OTP-team; based on ideas hatched by Tony Rogvall
- › When will these features show up?
 - Hopefully R15, but this is *no* promise



ERICSSON